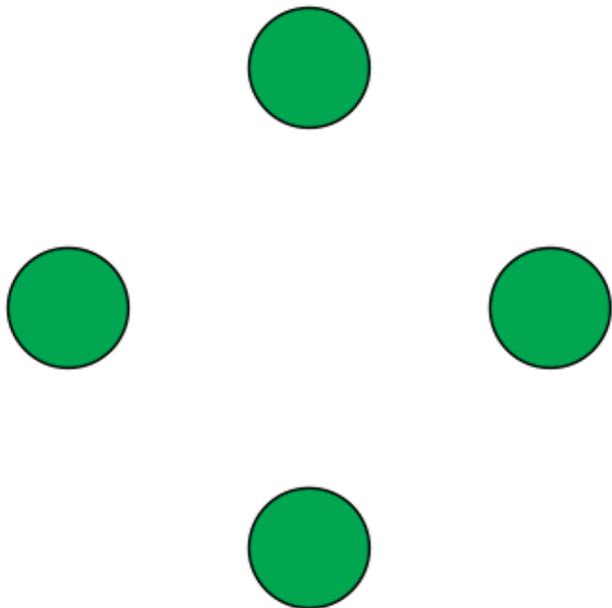


# FnF-BFT: Exploring Performance Limits of BFT Protocols

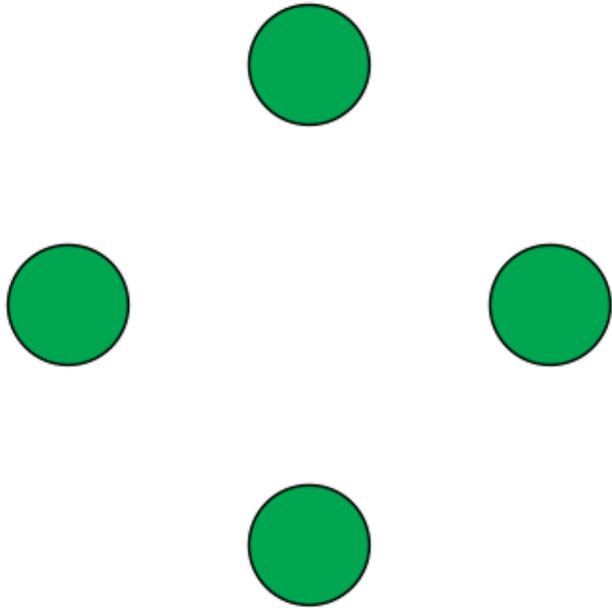


Zeta Avarikioti, Lioba Heimbach, Roland Schmid, Laurent Vanbever, Roger Wattenhofer, Patrick Wintermeyer  
ETH Zurich – Distributed Computing – [www.disco.ethz.ch](http://www.disco.ethz.ch)

# Byzantine fault tolerance

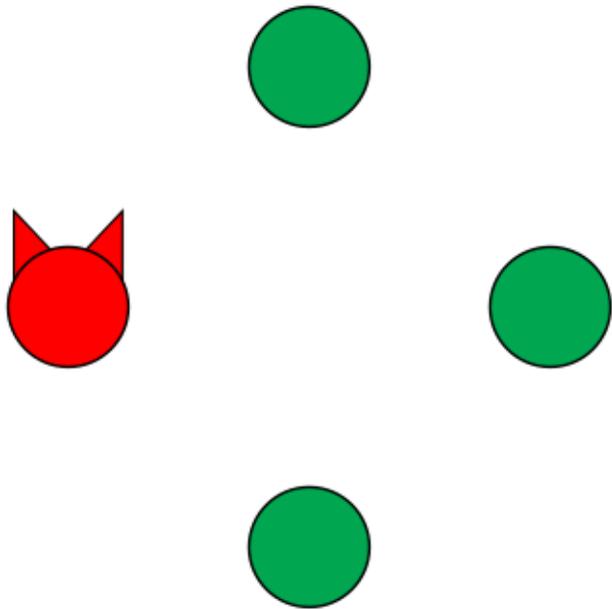


# Byzantine fault tolerance



$$n = 3f + 1 \text{ replicas}$$

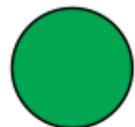
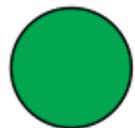
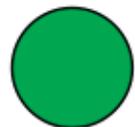
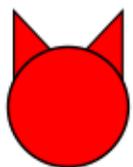
# Byzantine fault tolerance



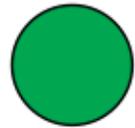
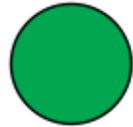
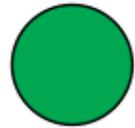
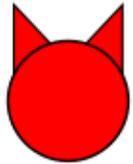
$n = 3f + 1$  replicas

$f$  replicas may byzantine

# Properties

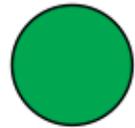
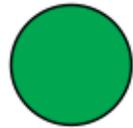
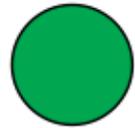
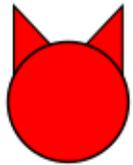


# Properties



safety

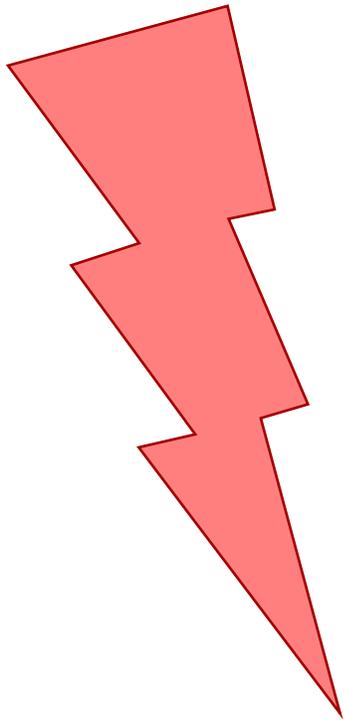
# Properties



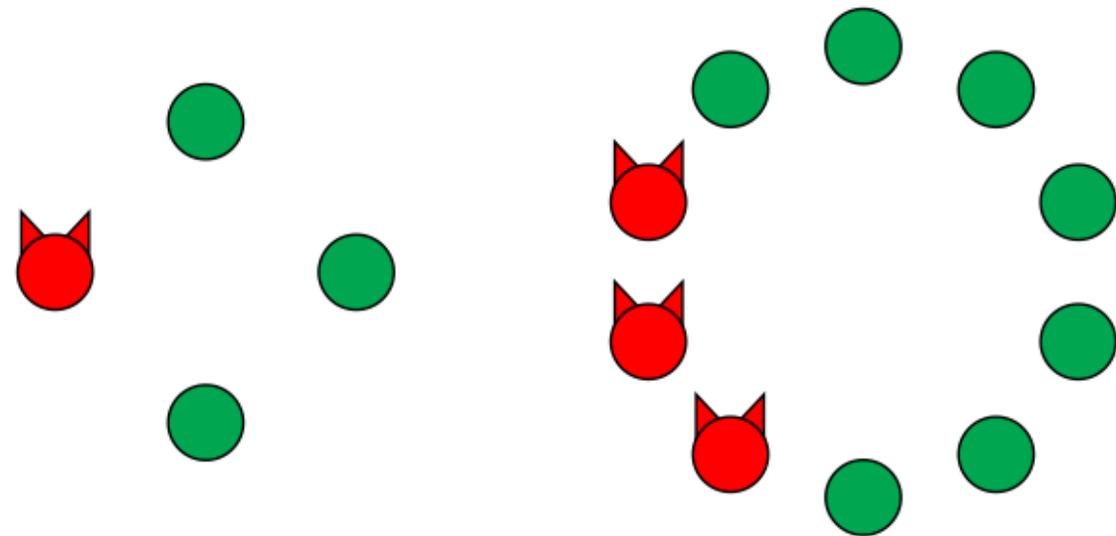
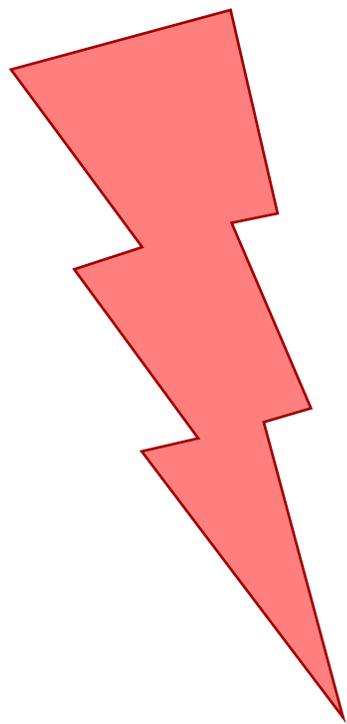
safety

liveness

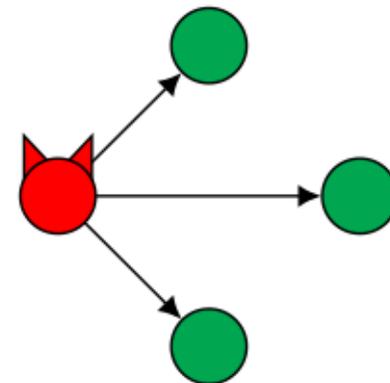
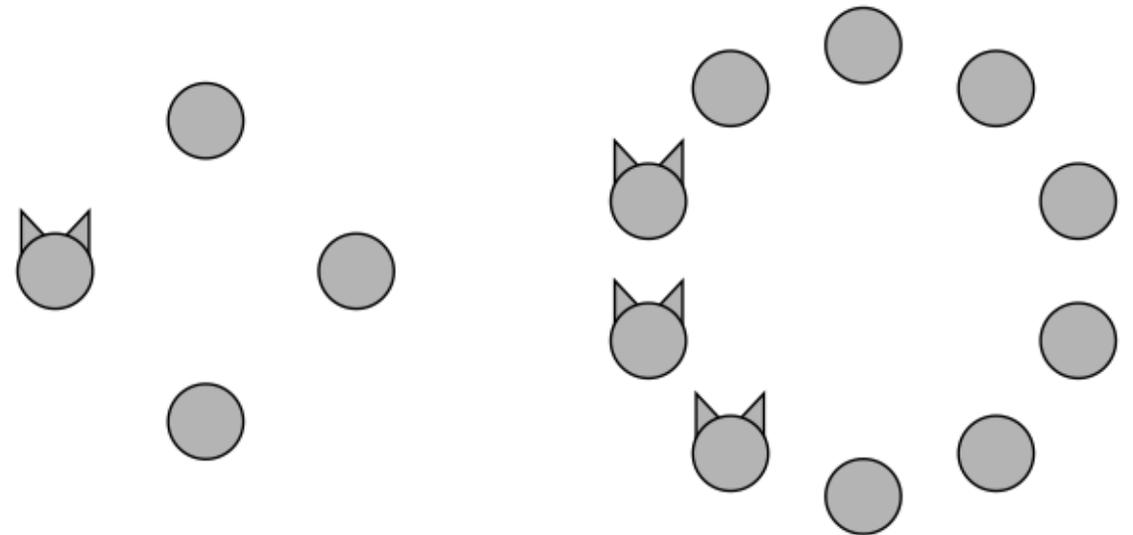
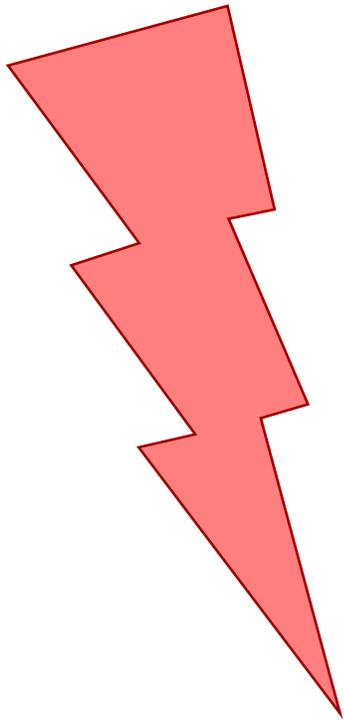
# BFT protocols

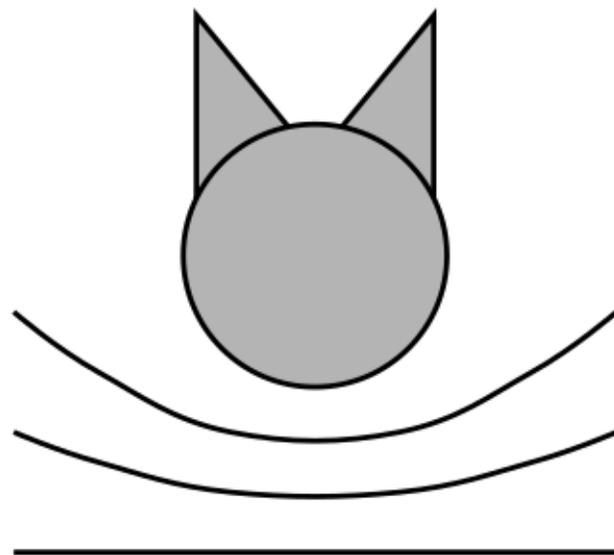
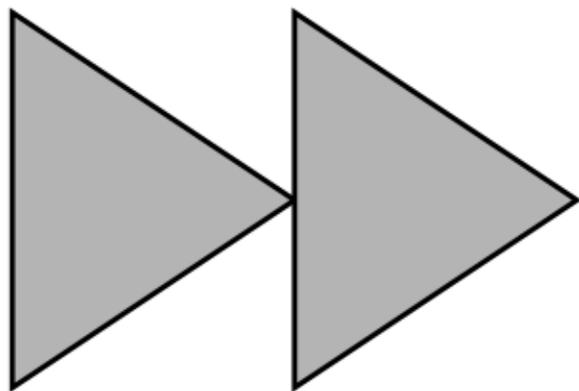
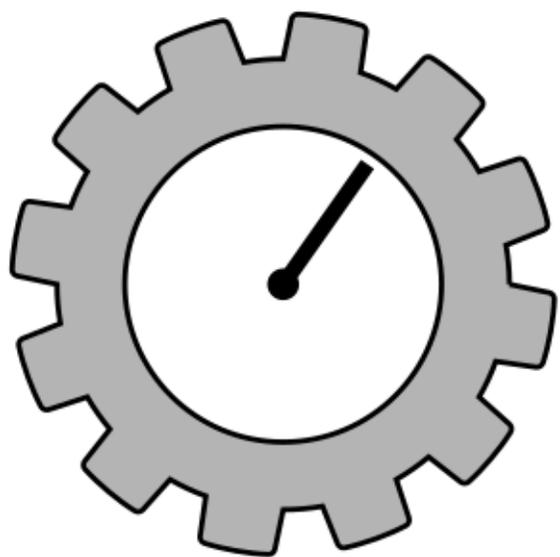


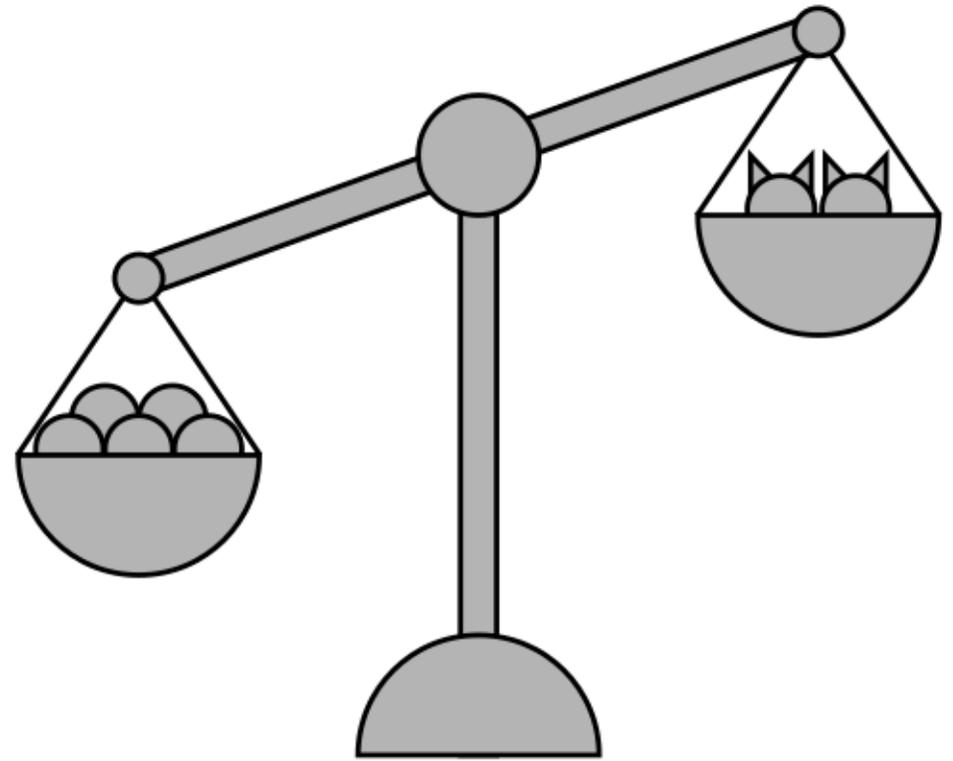
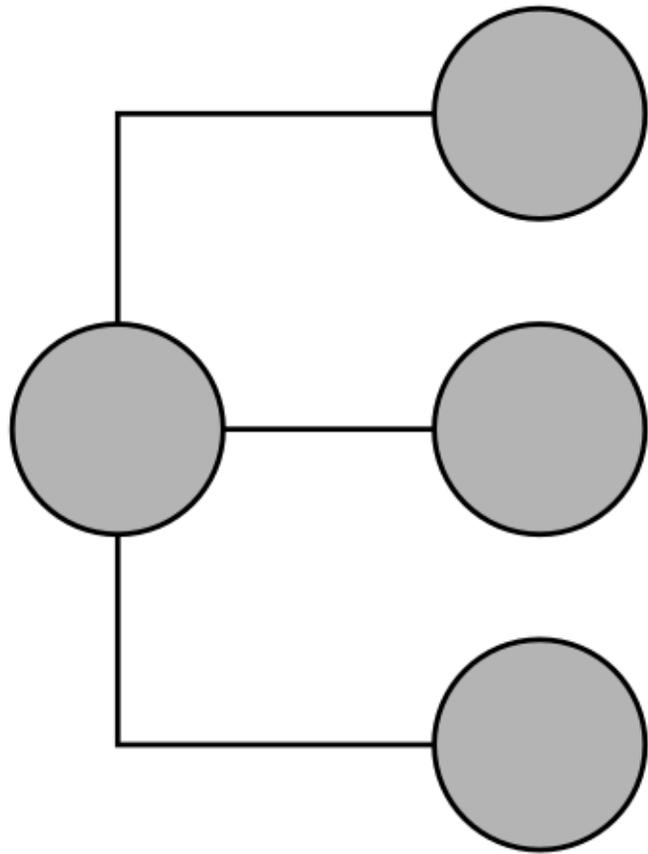
# BFT protocols



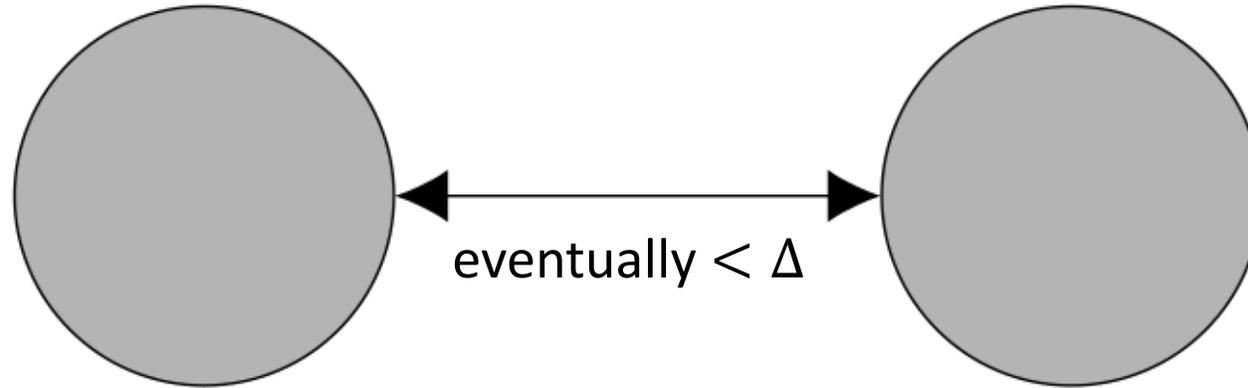
# BFT protocols



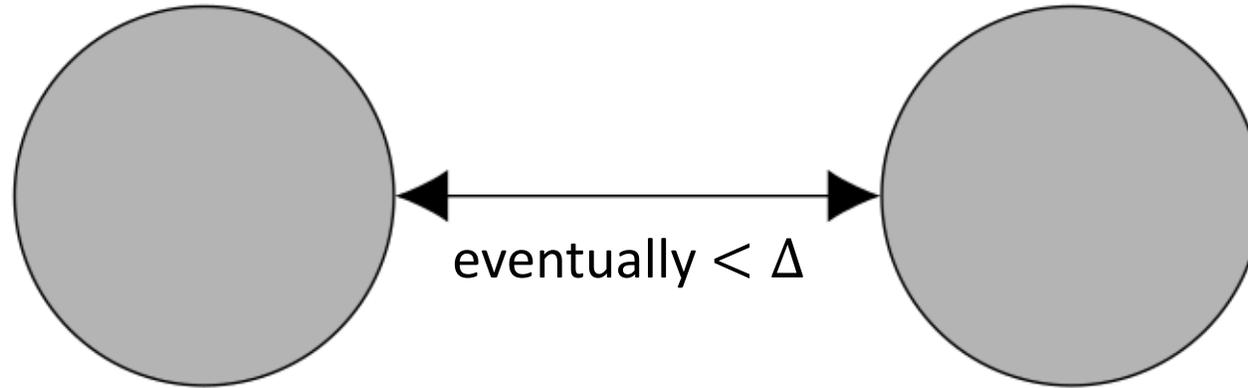




# Communication model

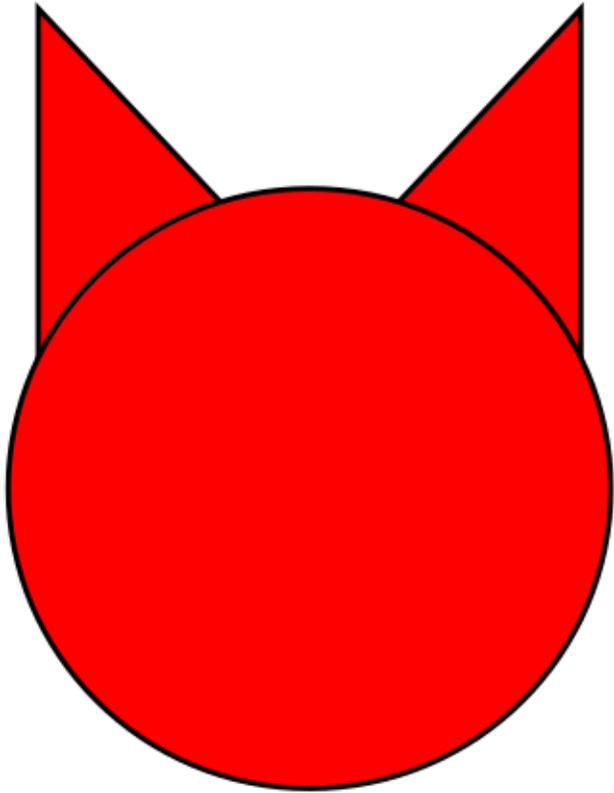


# Communication model



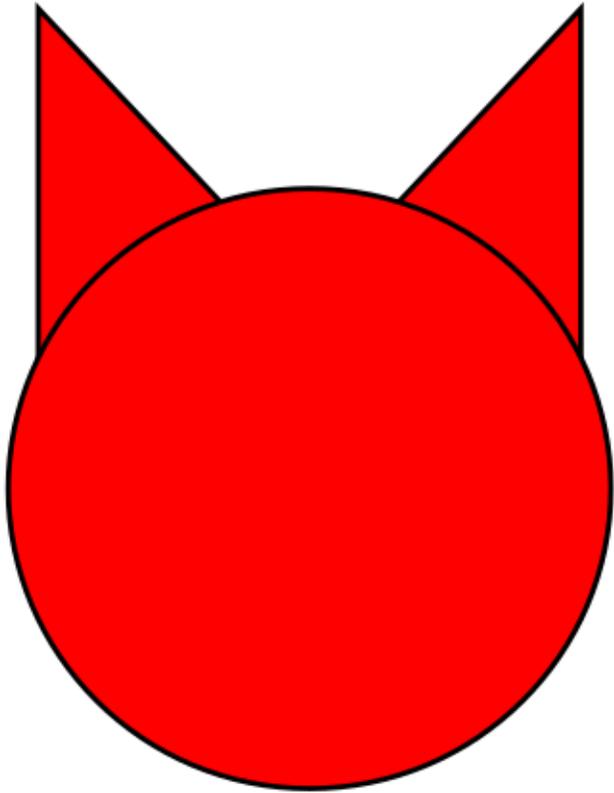
$n = 3f + 1$  replicas,  $f$  replicas may byzantine

# Model



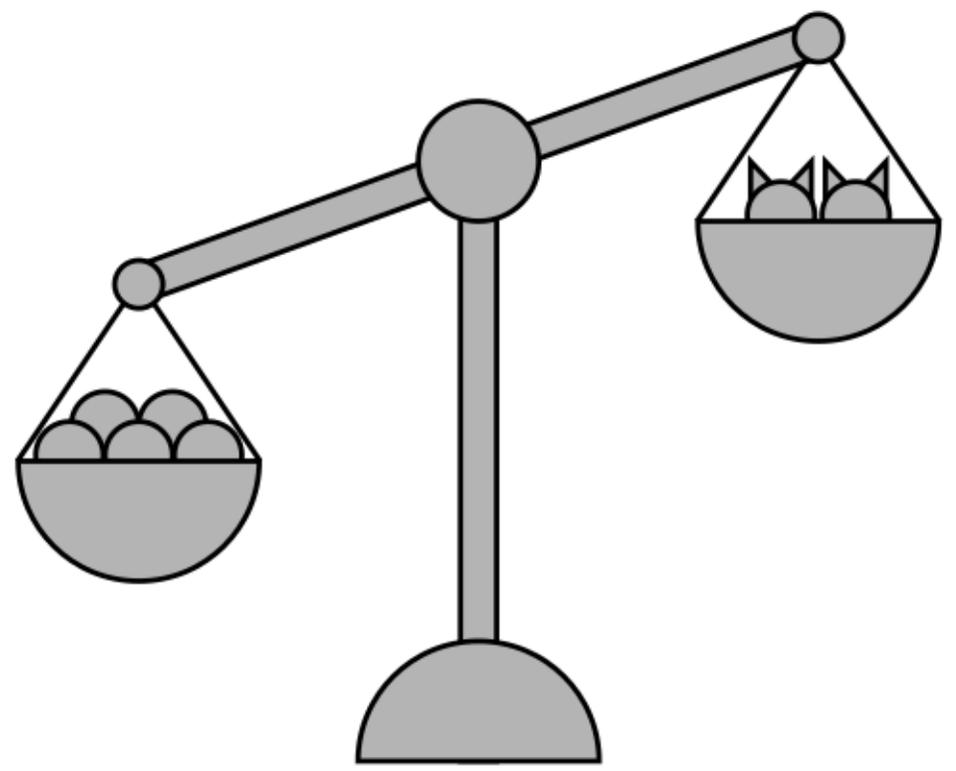
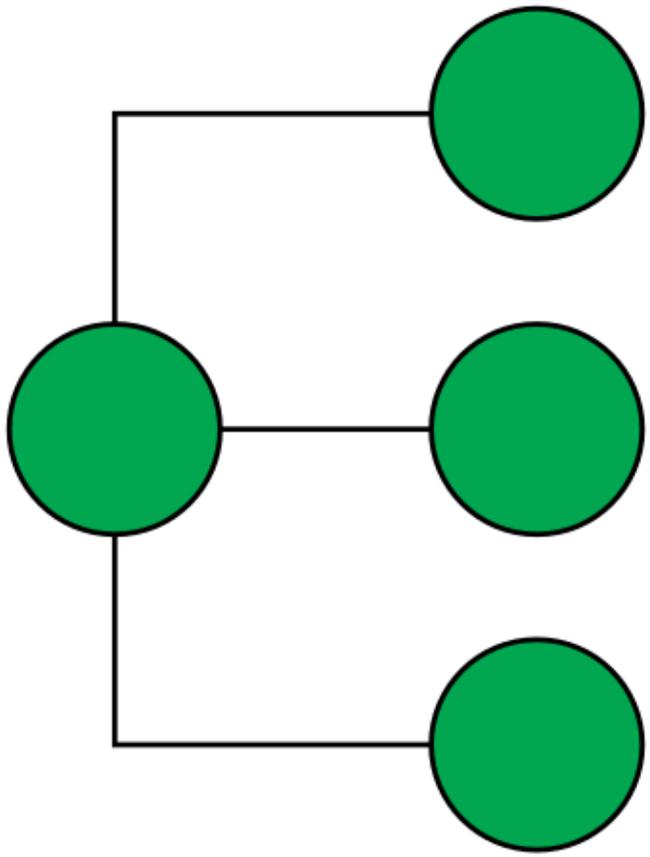
$f$  byzantine replicas

# Model

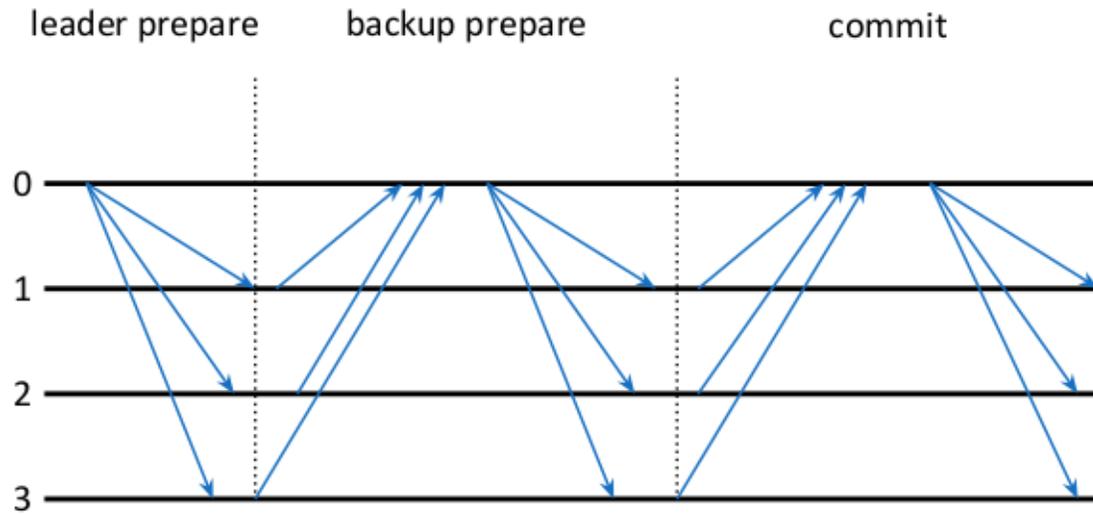


$f$  byzantine replicas

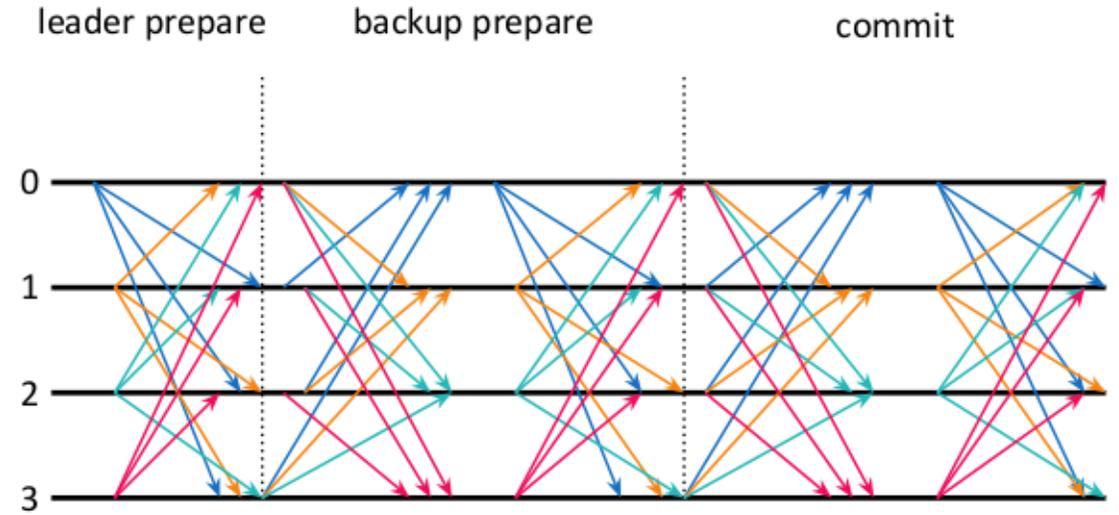
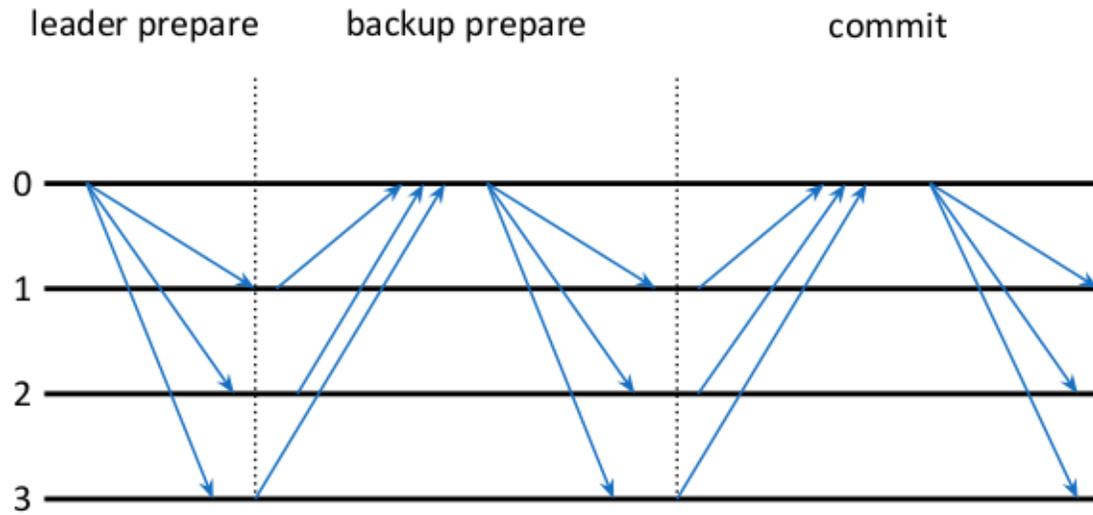
any number byzantine clients

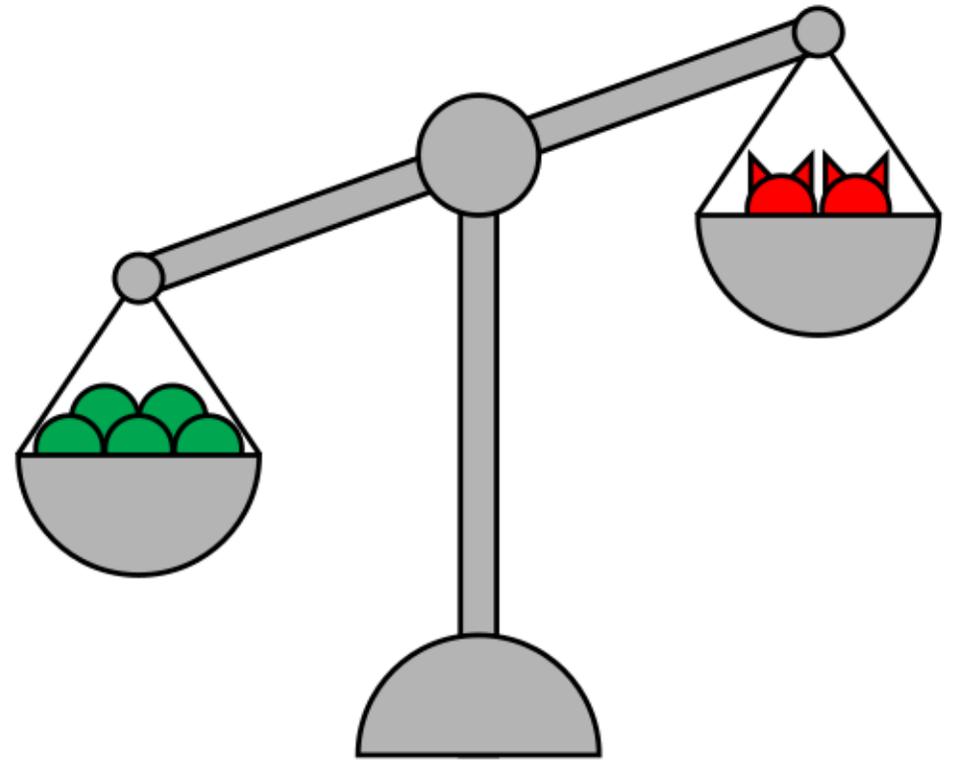
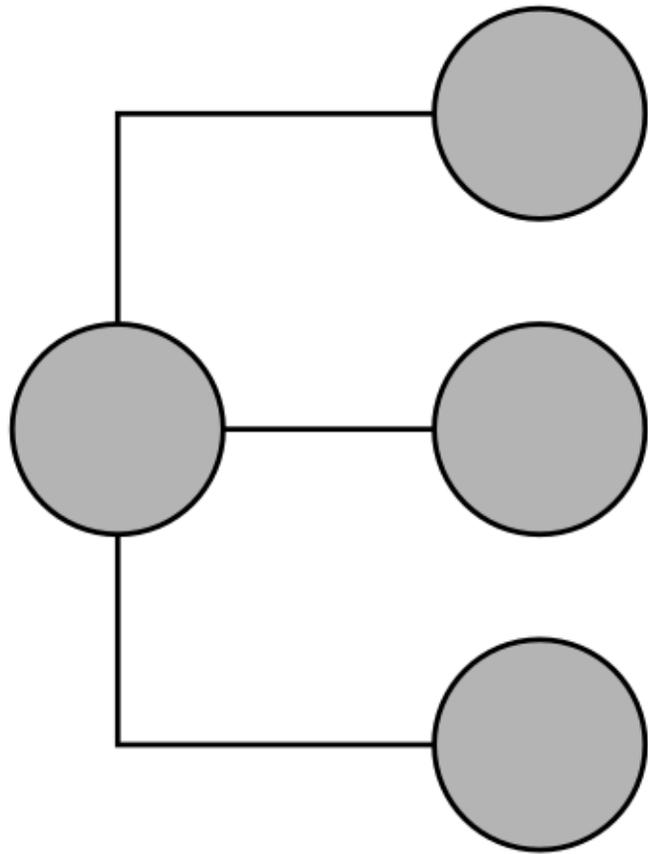


# Parallelize

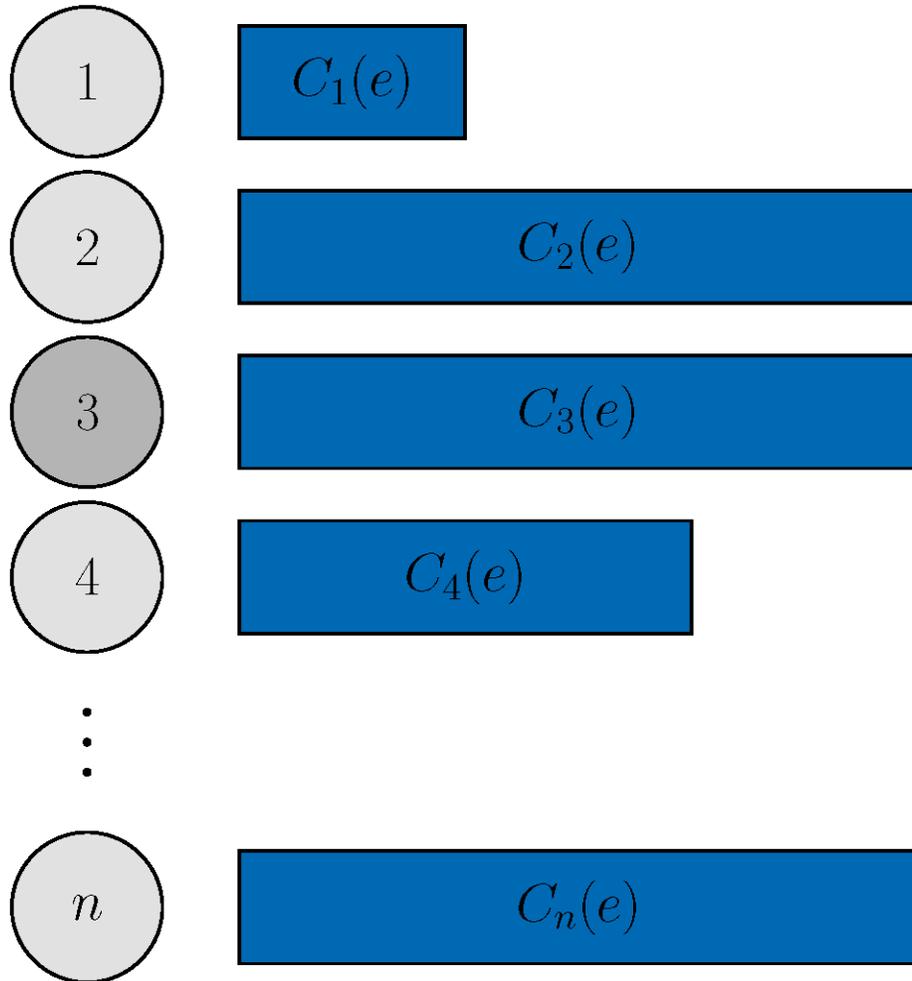


# Parallelize

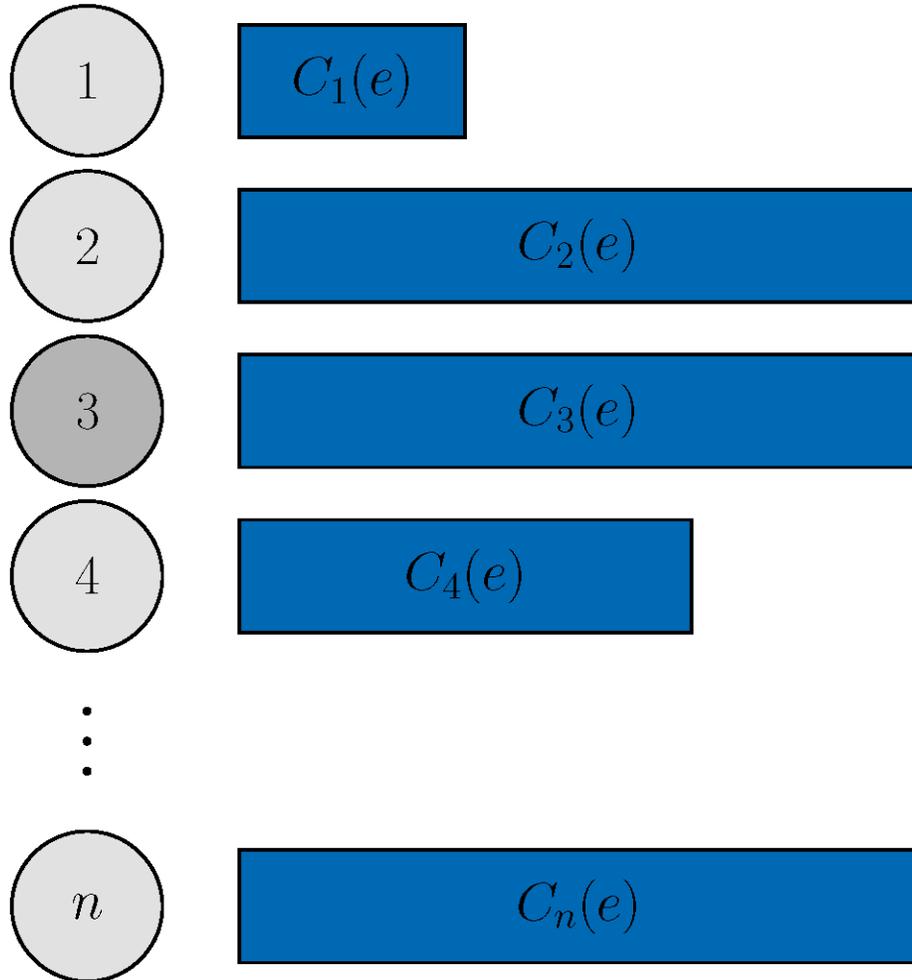




# Parallel leaders

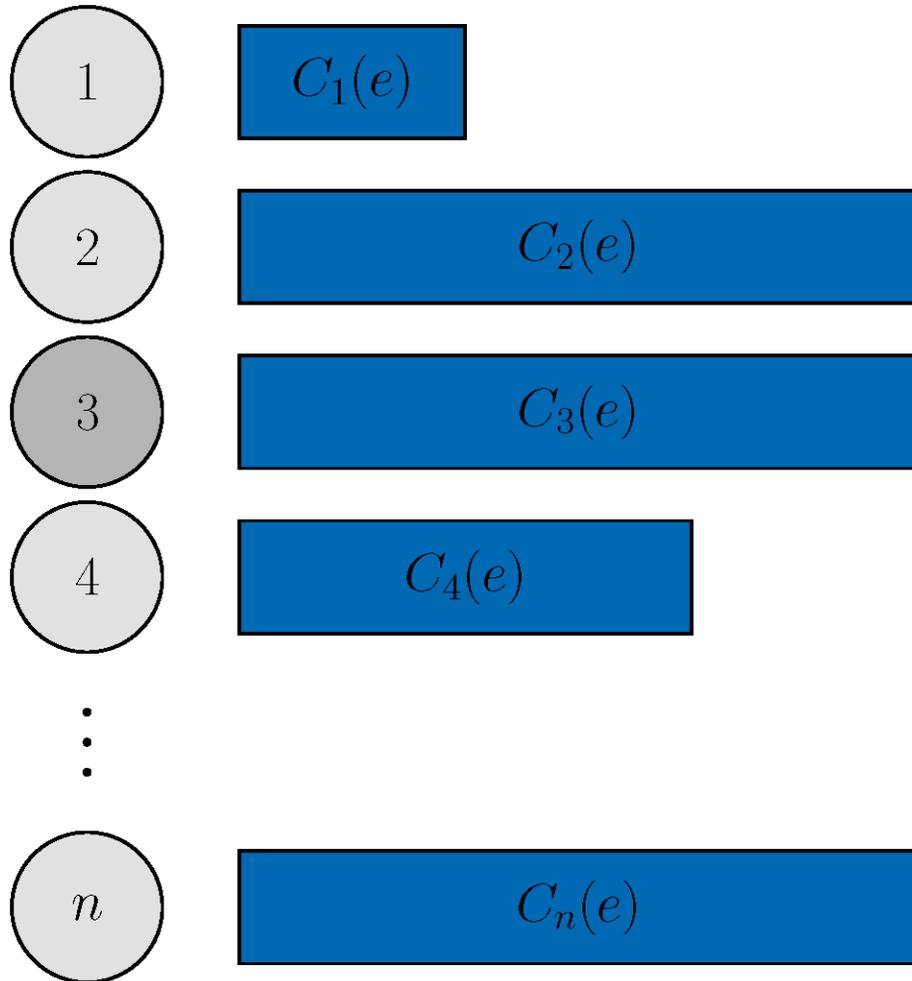


# Configuration



$C_v(e)$  number of requests  
assigned to leader  $v$

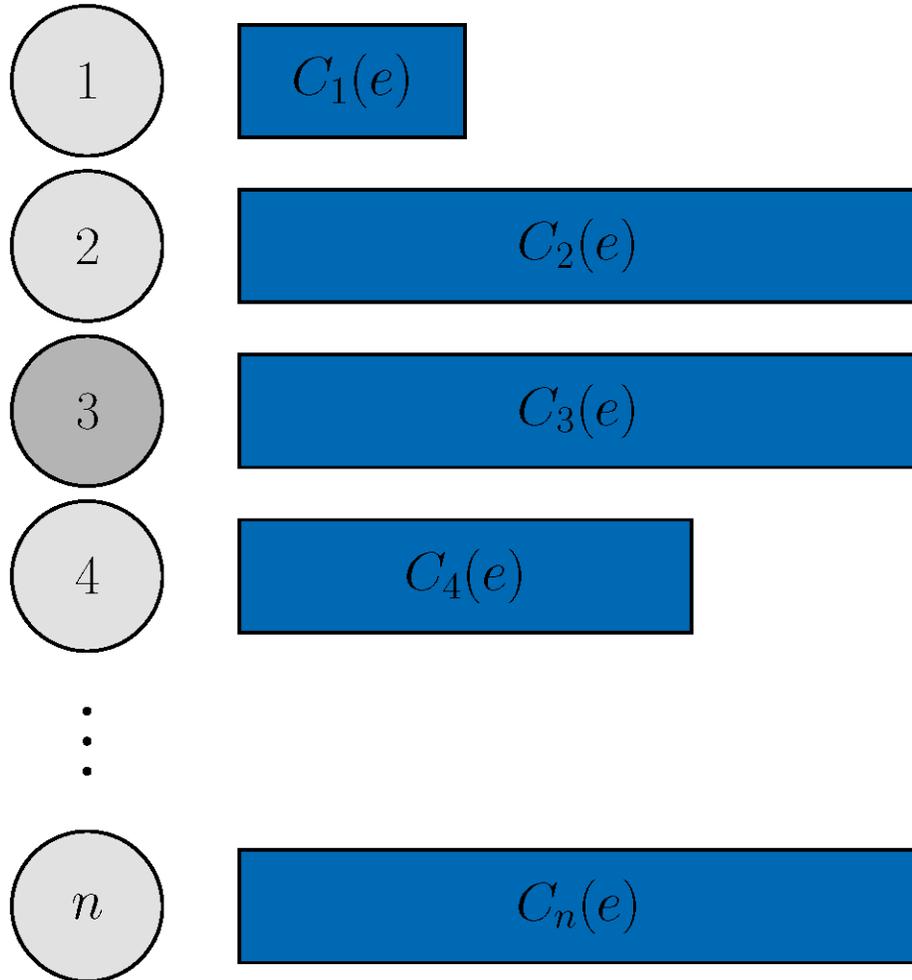
# Configuration



$C_v(e)$  number of requests  
assigned to leader  $v$

$C_v(e)$  initially  $C_{\min} \in \Omega(n^2)$   
for all  $v \in [n]$

# Optimization

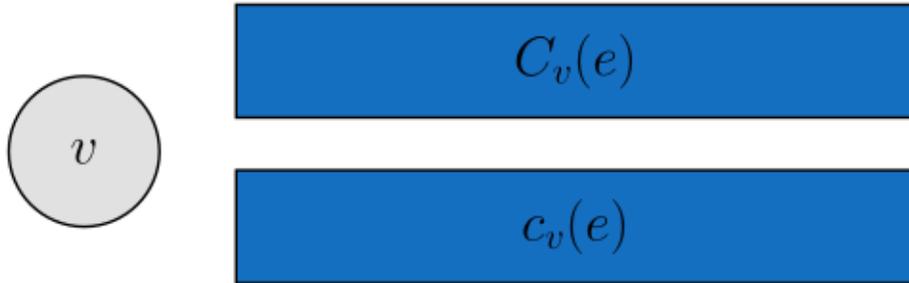


$C_v(e)$  number of requests  
assigned to leader  $v$

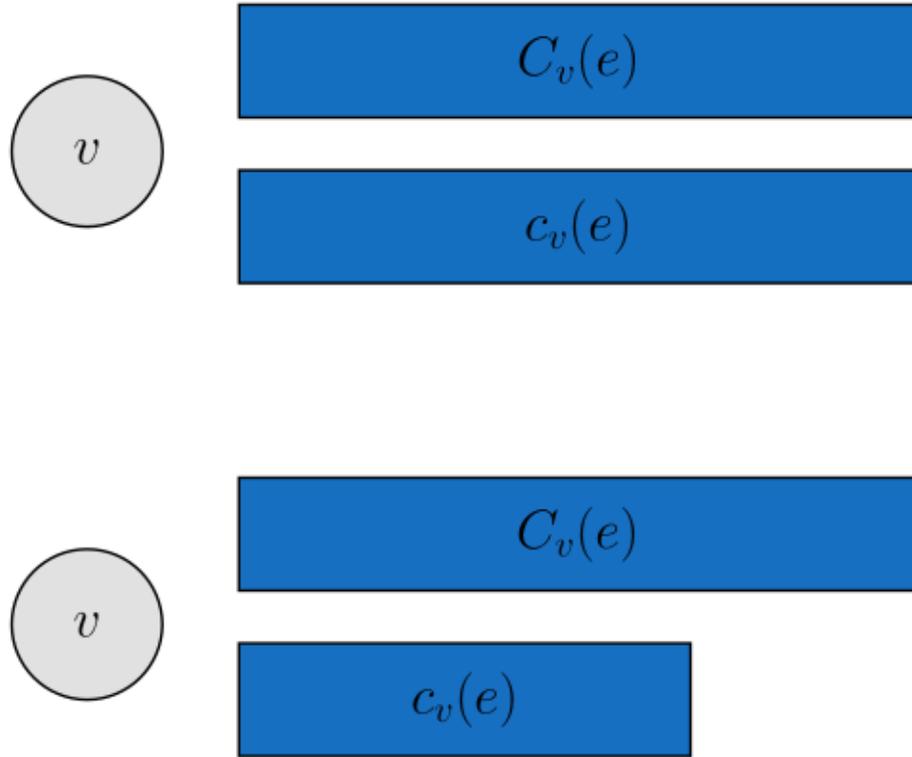
$C_v(e)$  initially  $C_{\min} \in \Omega(n^2)$   
for all  $v \in [n]$

$C_v(e)$  updated with  
leader  $v$ 's performance

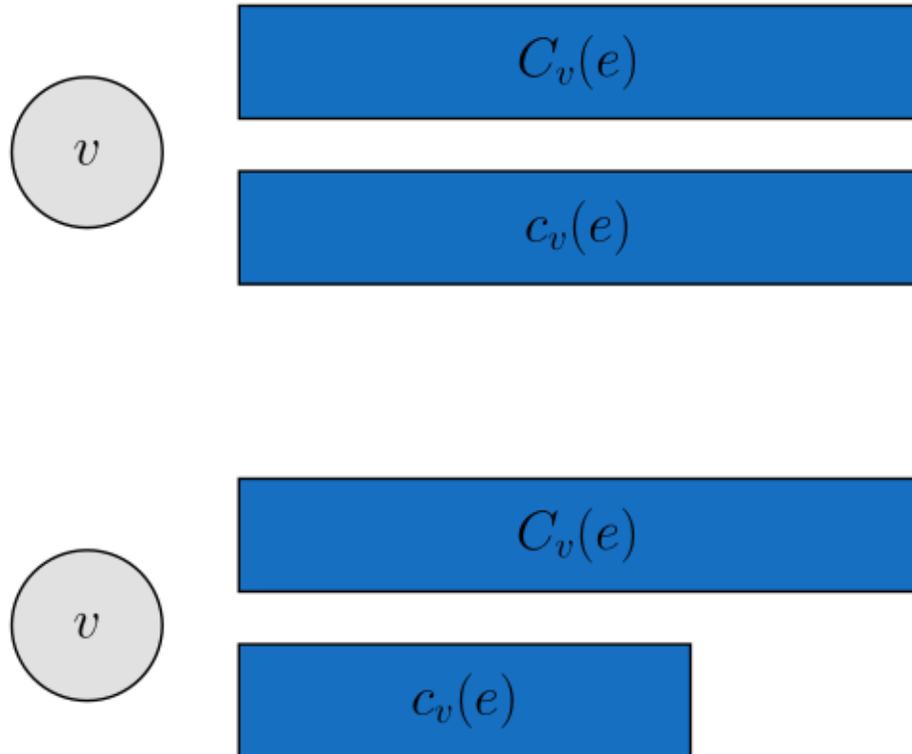
# Optimization



# Optimization

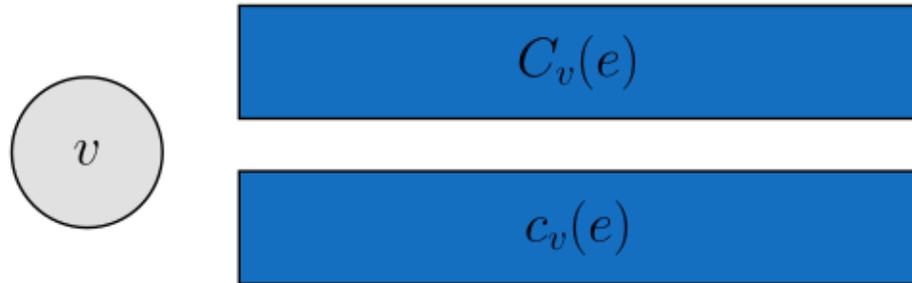


# Optimization

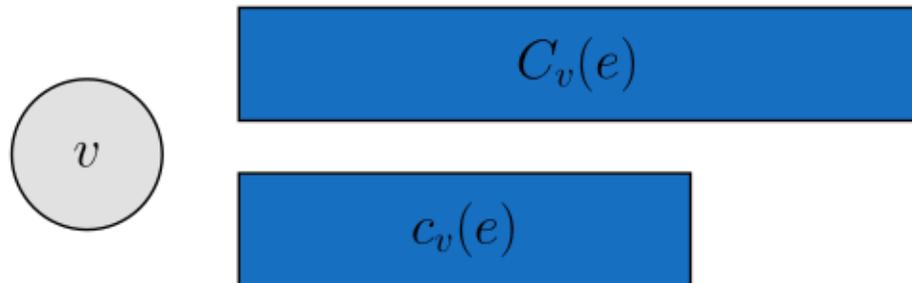


$$C_v(e + 1) = 2 \cdot c_v(e)$$

# Optimization

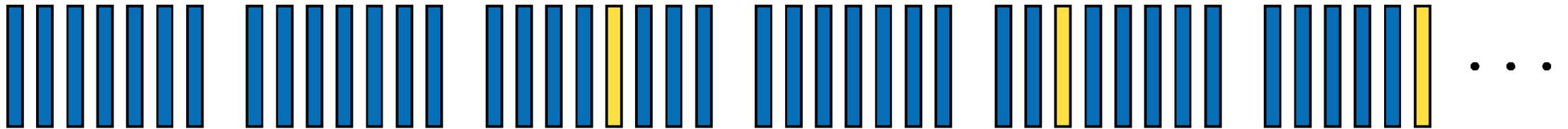


$$C_v(e + 1) = 2 \cdot c_v(e)$$

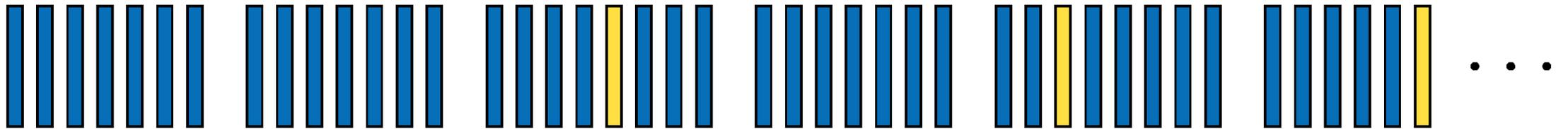


$$C_v(e + 1) = \max(C_{\min}, \max_{i \in \{0, \dots, f\}} c_v(e - i))$$

# Primary rotation

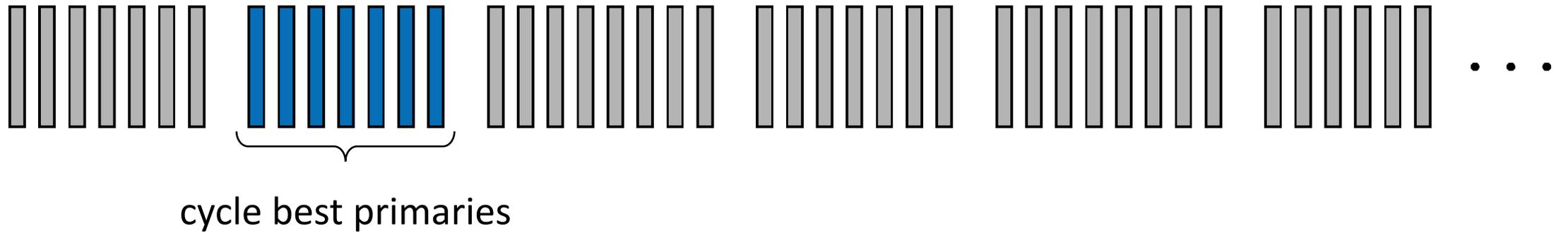


# Primary rotation

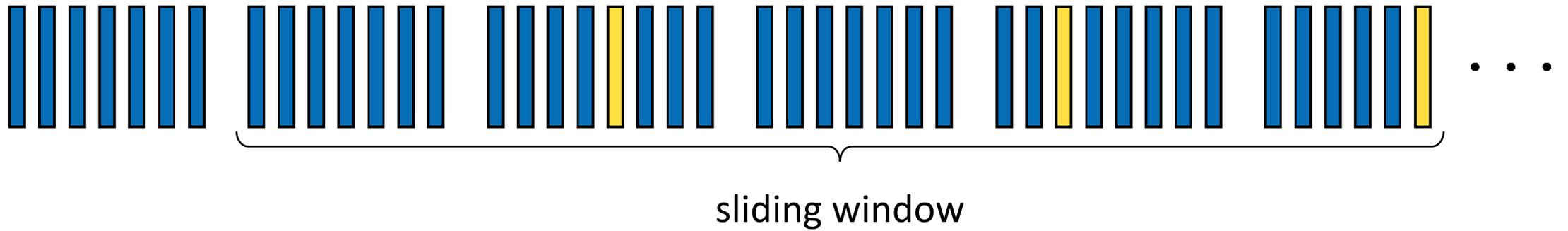


explore and exploit

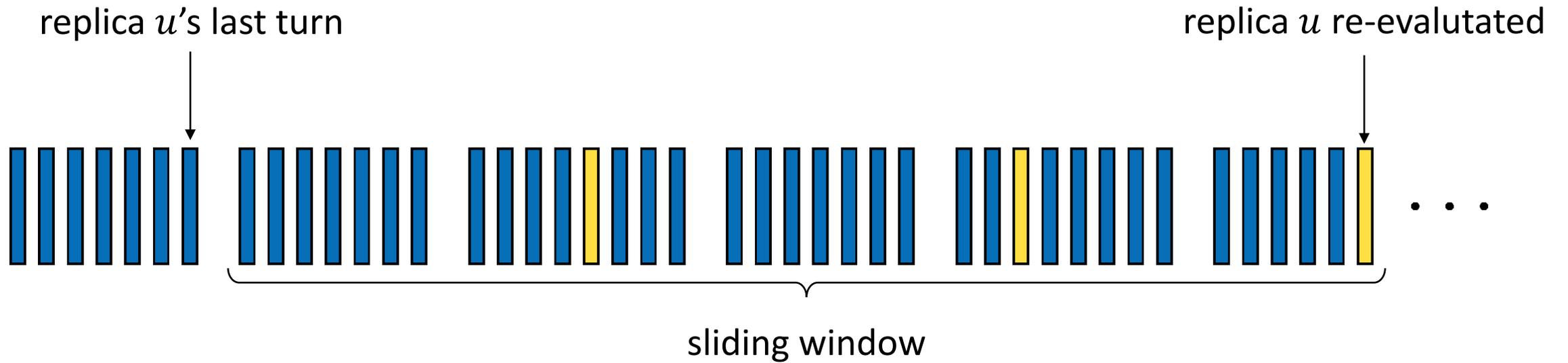
# Exploit best primaries

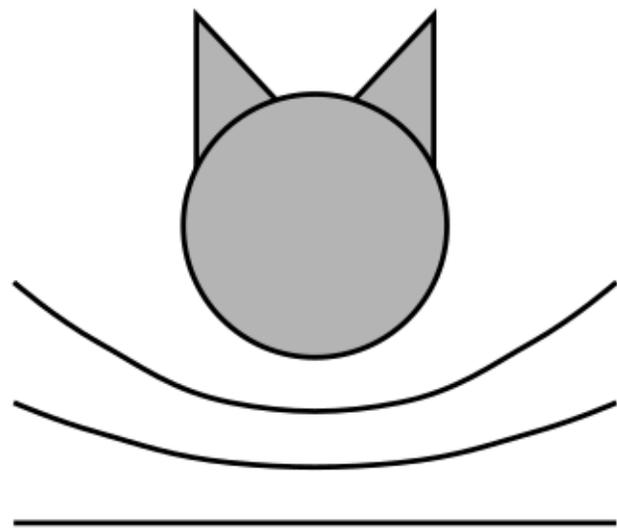
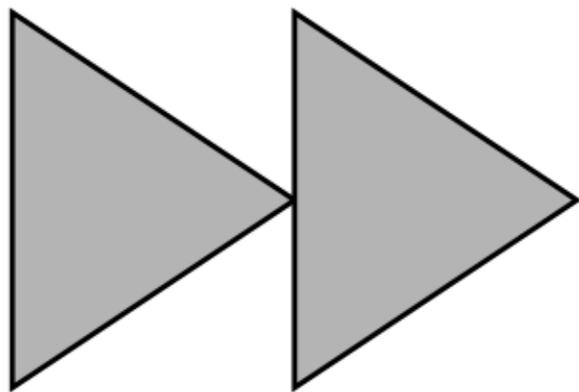


# Re-evaluate primaries



# Re-evaluate primaries



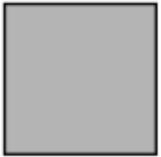


# Efficiency

complexity measure: authenticator complexity

# Efficiency

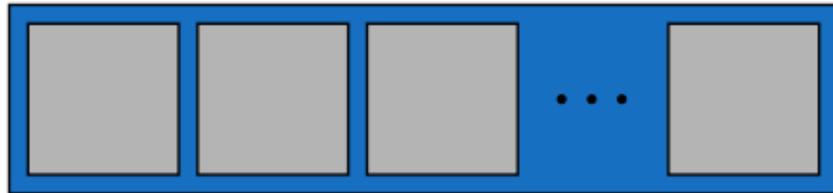
complexity measure: authenticator complexity



request creation cost:  $O(n)$

# Efficiency

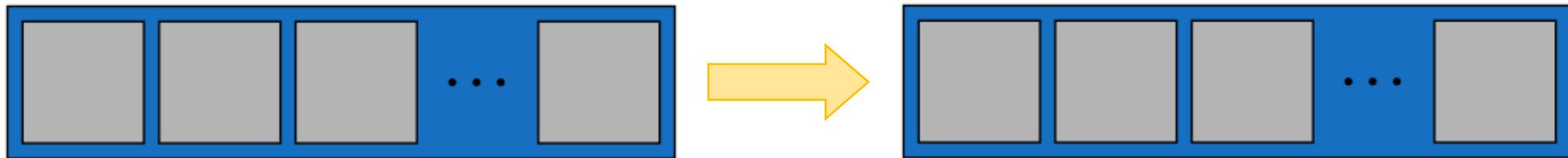
complexity measure: authenticator complexity



average number of requests:  $\Omega(n^3)$

# Efficiency

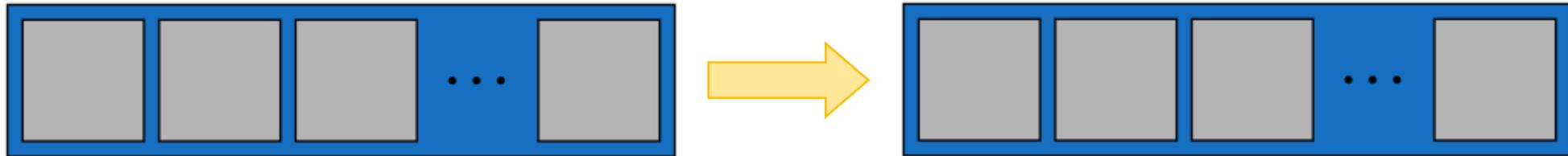
complexity measure: authenticator complexity



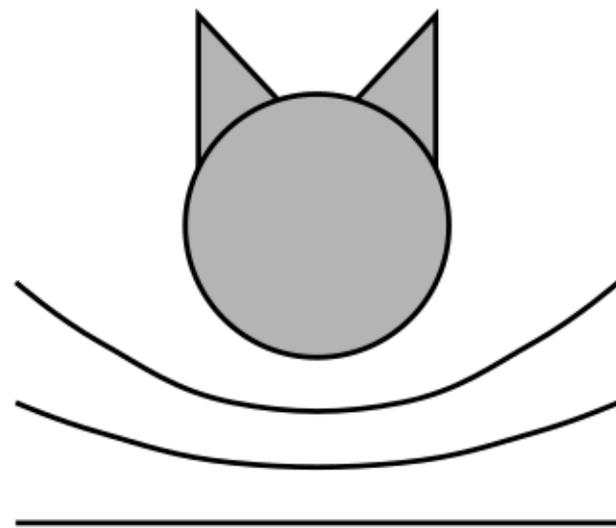
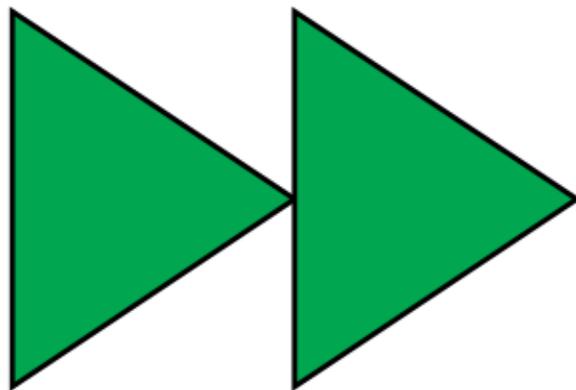
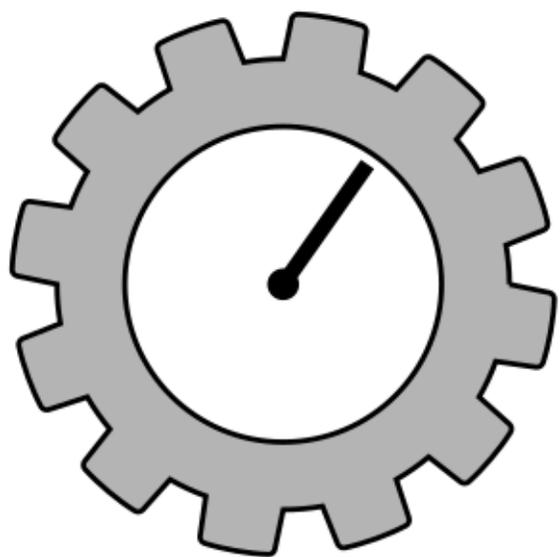
epoch-change cost:  $O(n^3)$

# Efficiency

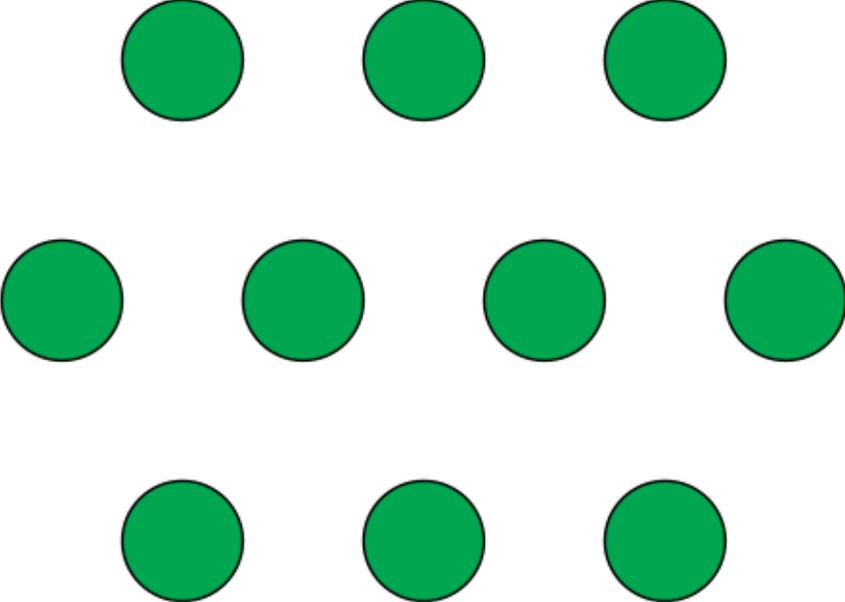
complexity measure: authenticator complexity



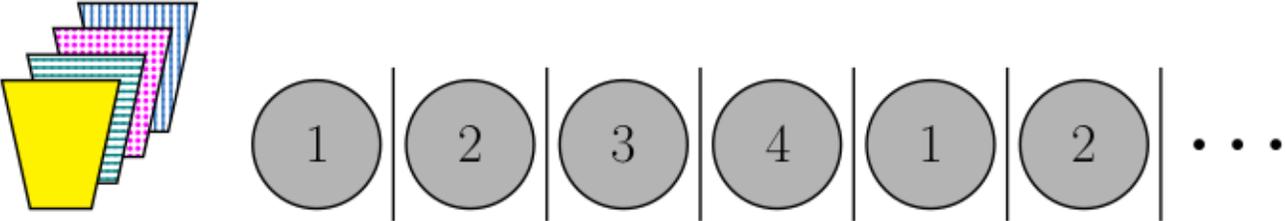
amortized request cost:  $O(n)$



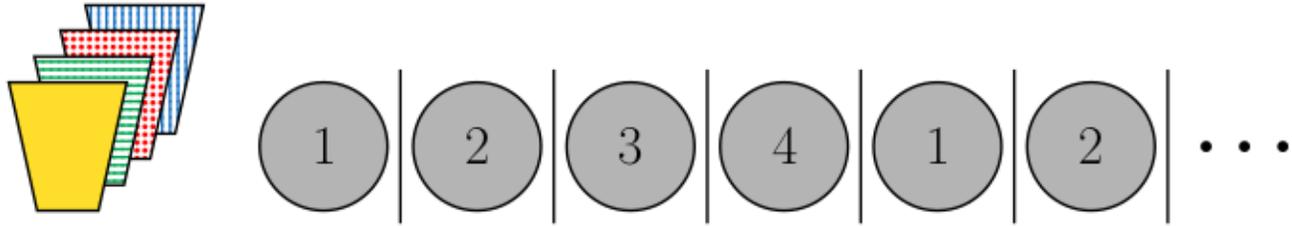
# Assumptions



# Speedup

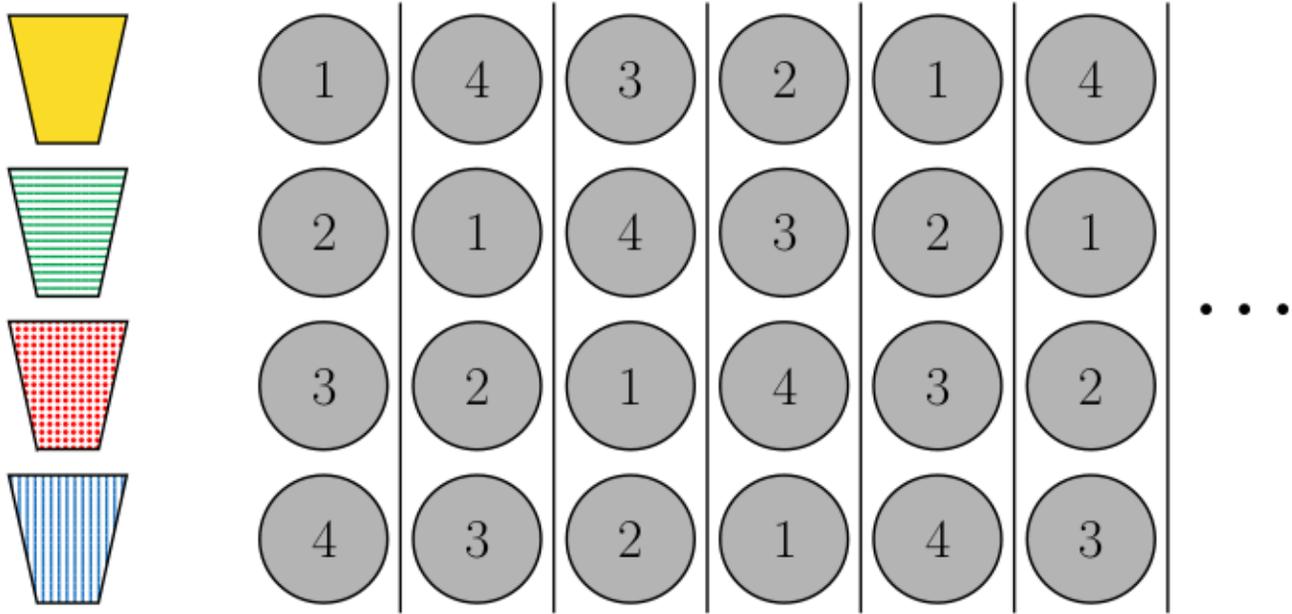


# Speedup

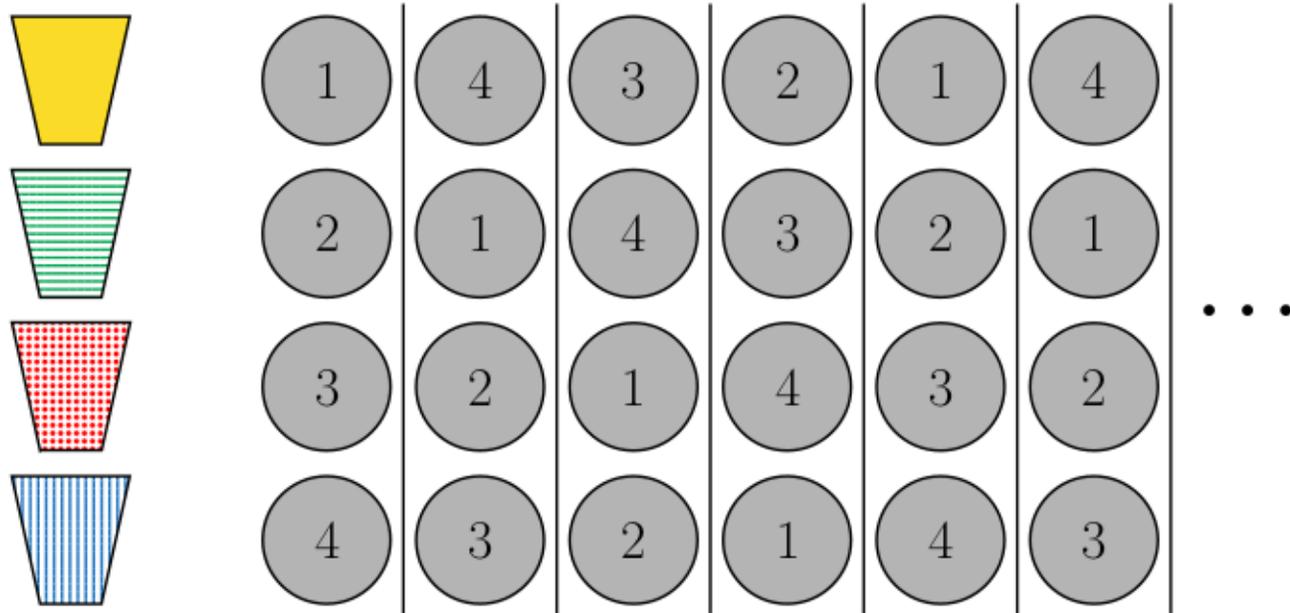


time units per request:  
 $\Theta(n)$

# Speedup

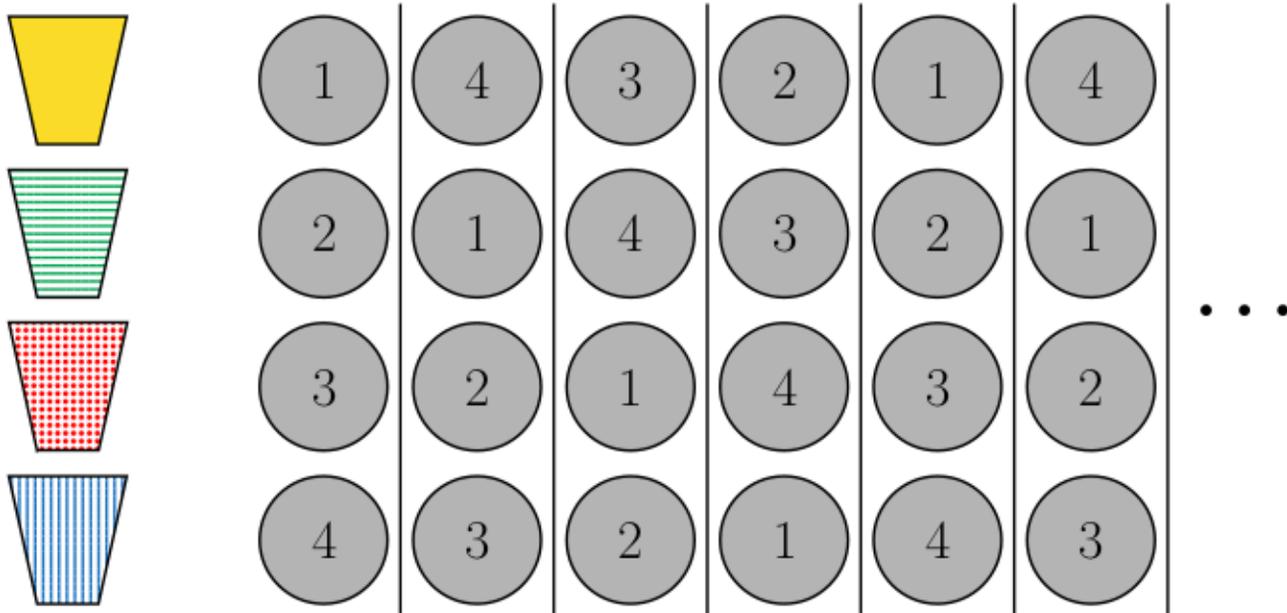


# Speedup



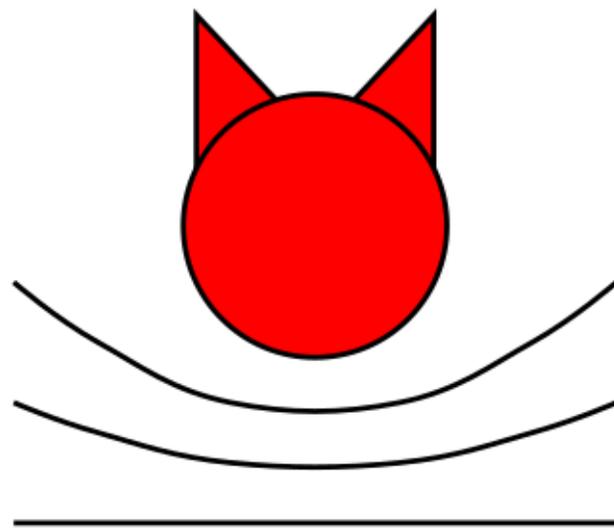
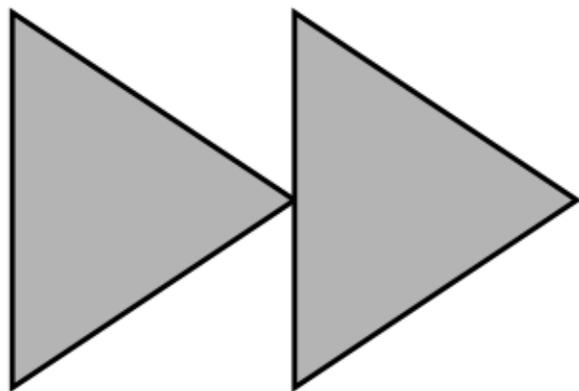
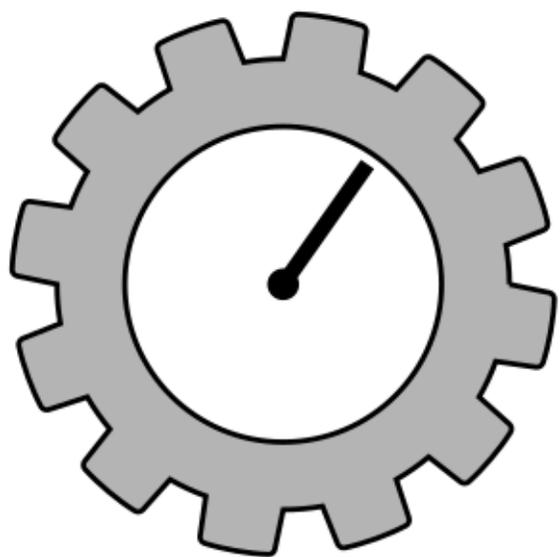
time units per request:  
 $\Theta(1)$

# Speedup

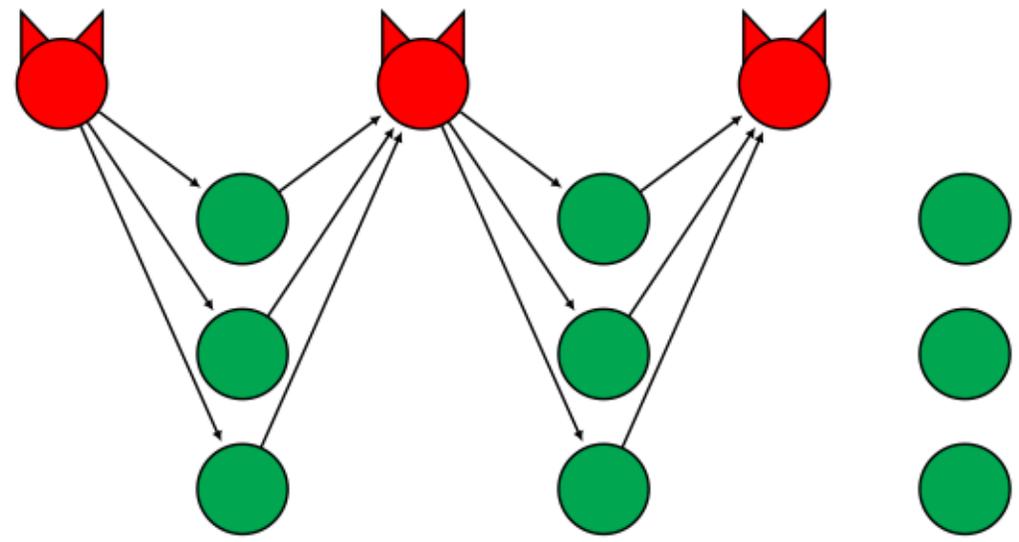


time units per request:  
 $\Theta(1)$

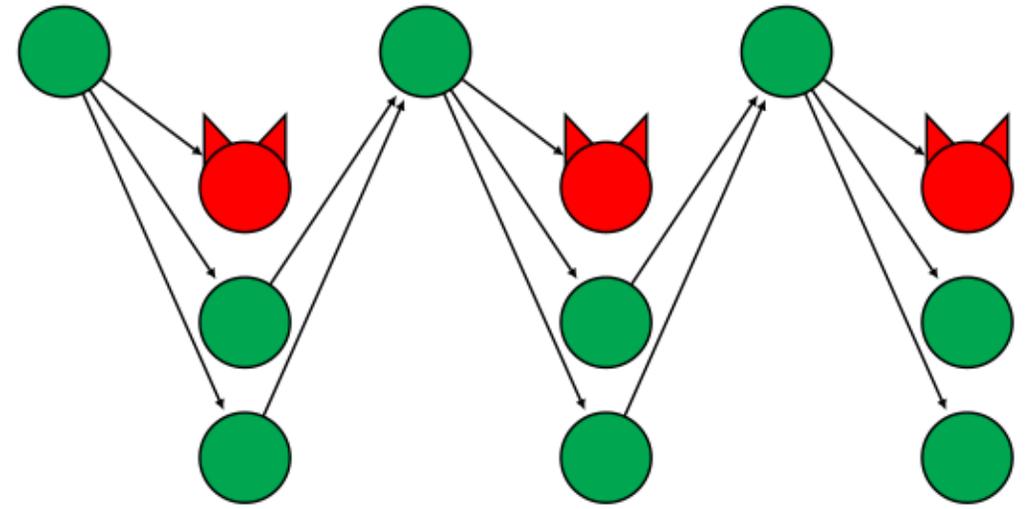
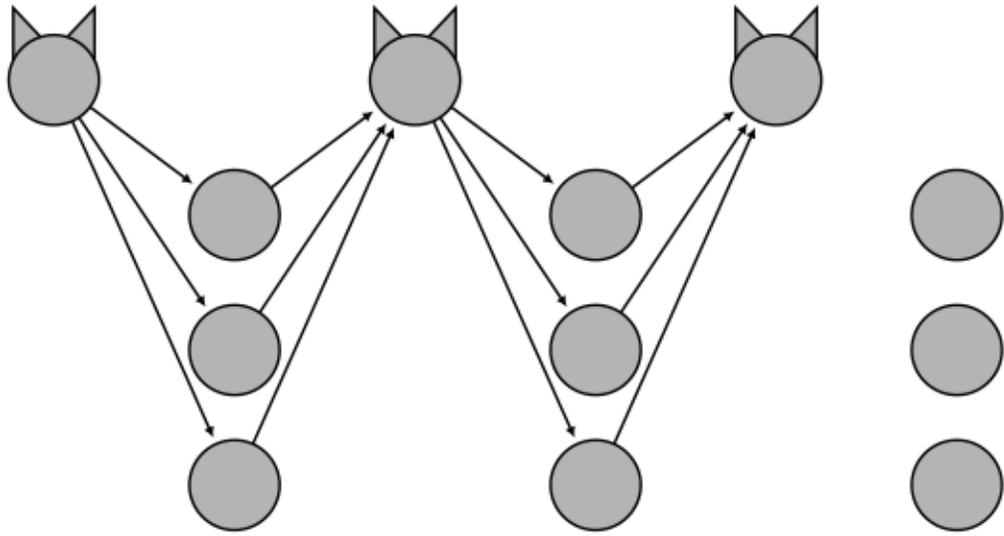
speedup:  
 $\Theta(n)$



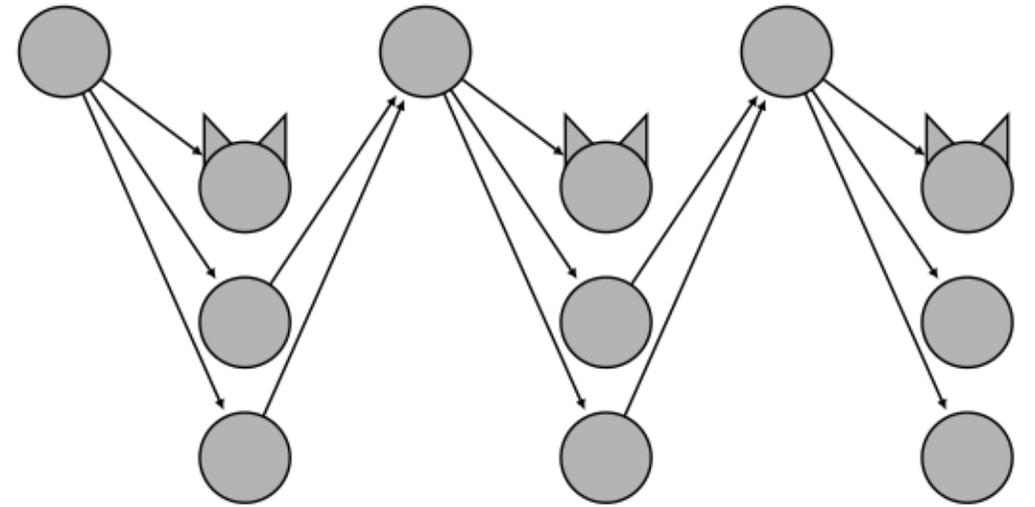
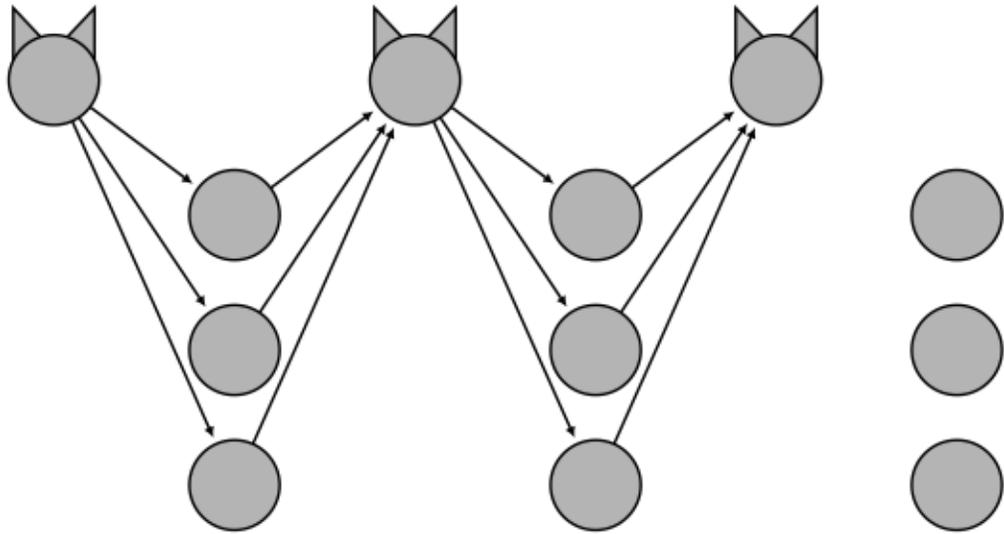
# Correct primary epochs



# Correct primary epochs

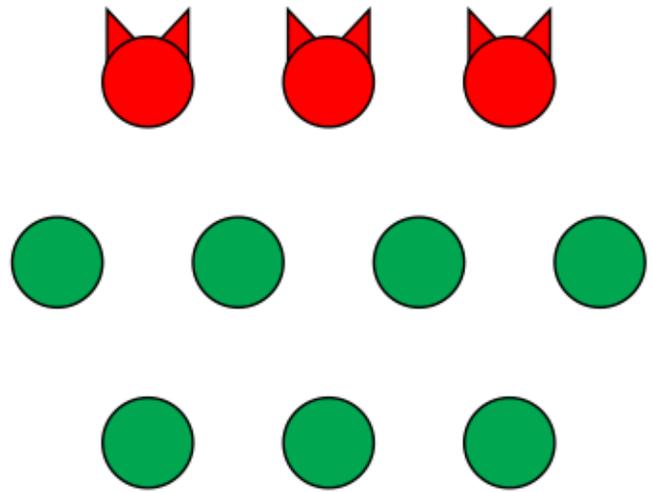


# Correct primary epochs

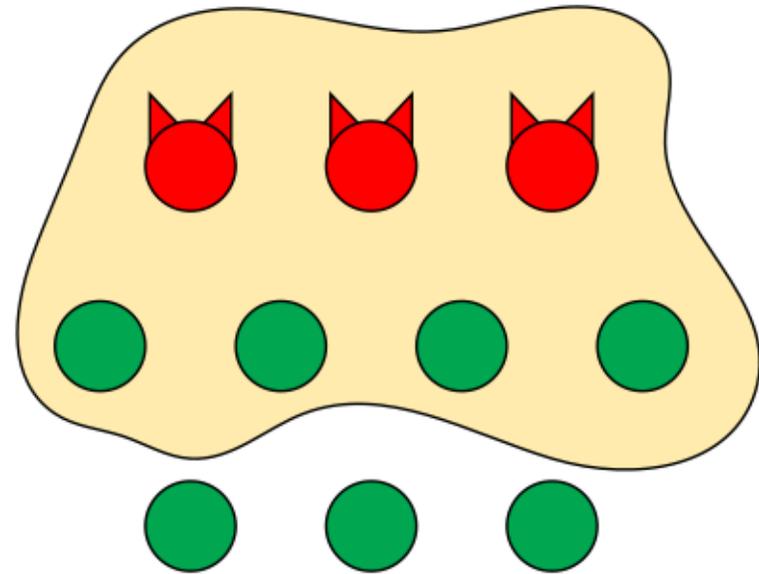
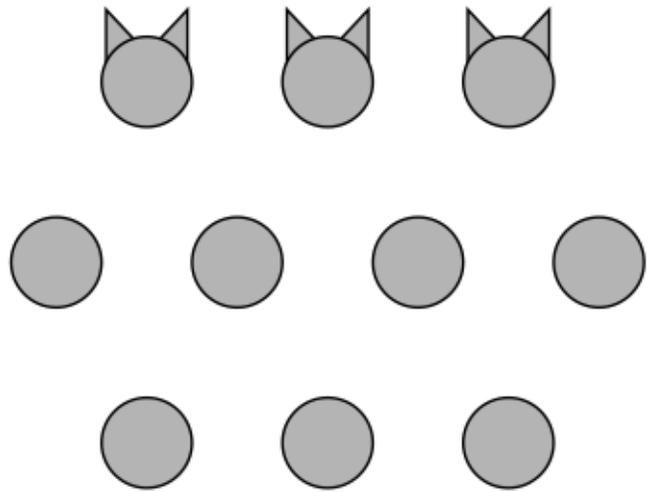


effective utilization:  $\frac{8}{9}$

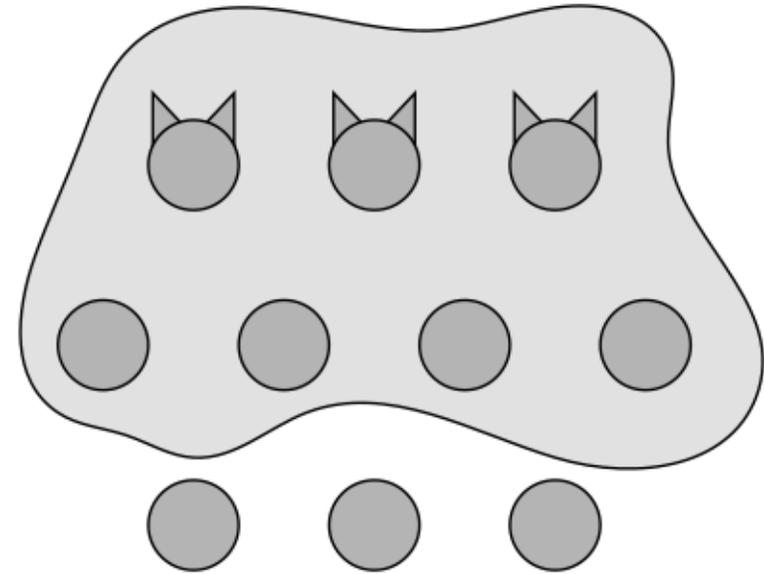
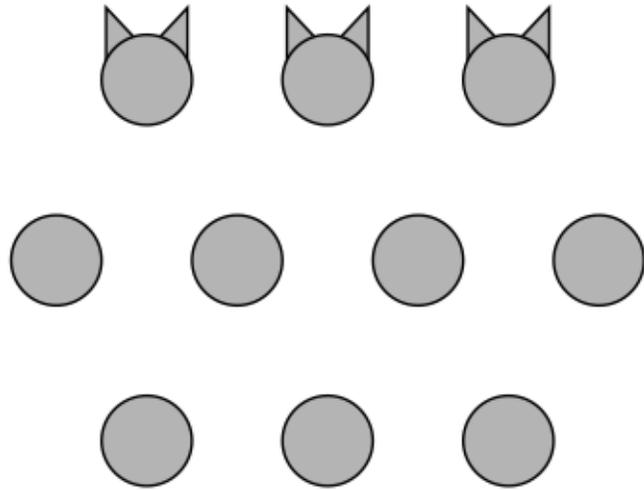
# Byzantine primary epochs



# Byzantine primary epochs

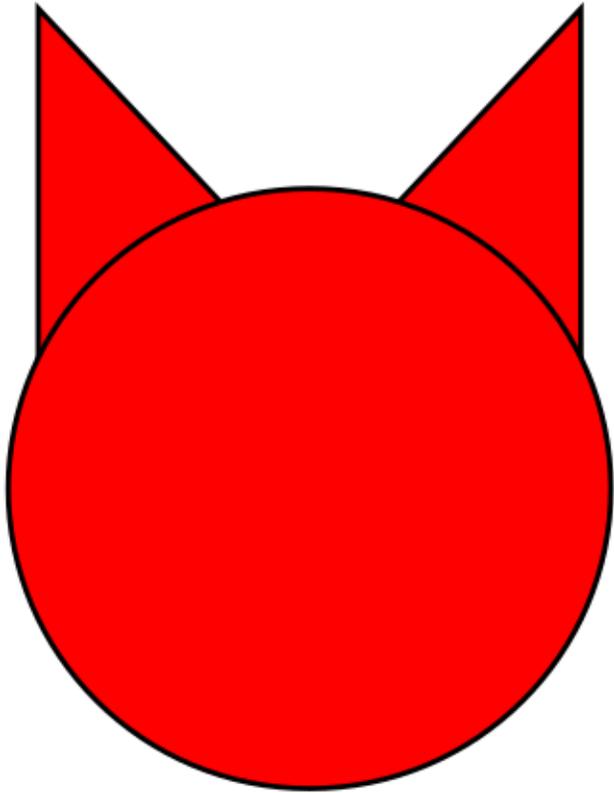


# Byzantine primary epochs



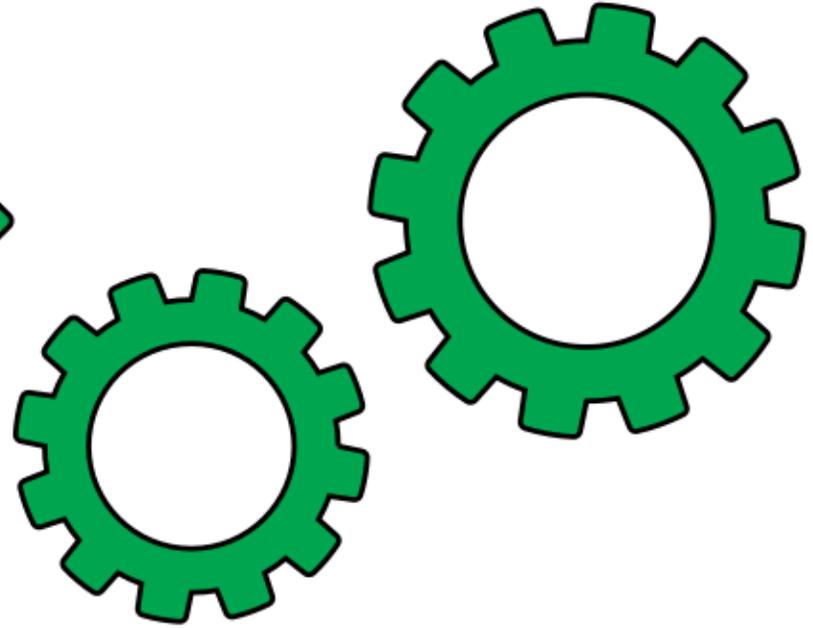
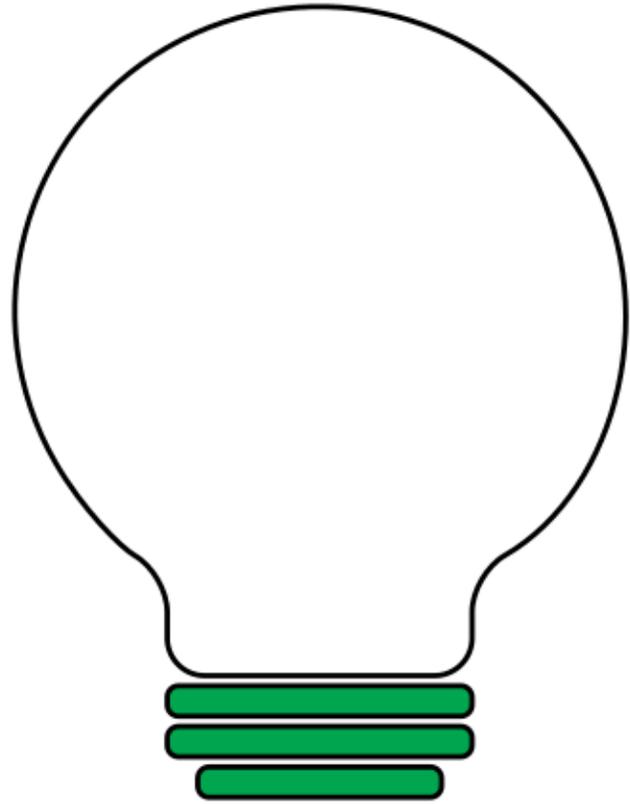
ratio of byzantine primaries:  $\frac{f}{g}$

# Effective utilization

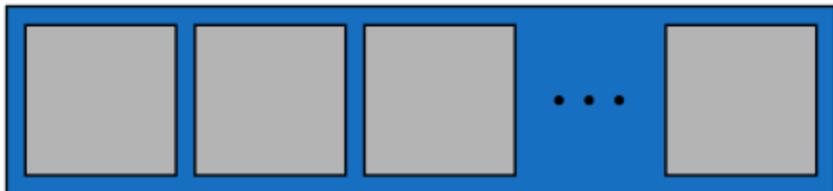


effective utilization:

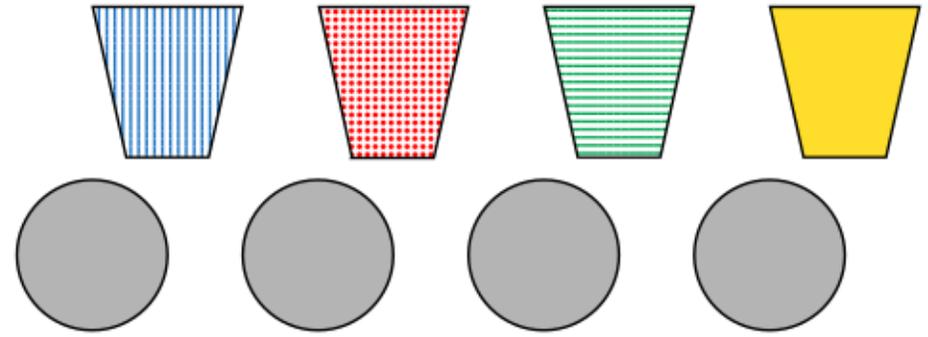
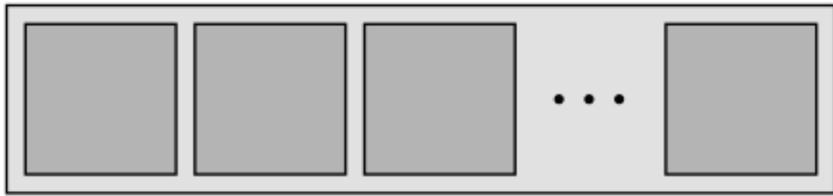
$$\frac{8}{9} \cdot \frac{g-f}{g} \geq \frac{16}{27}$$



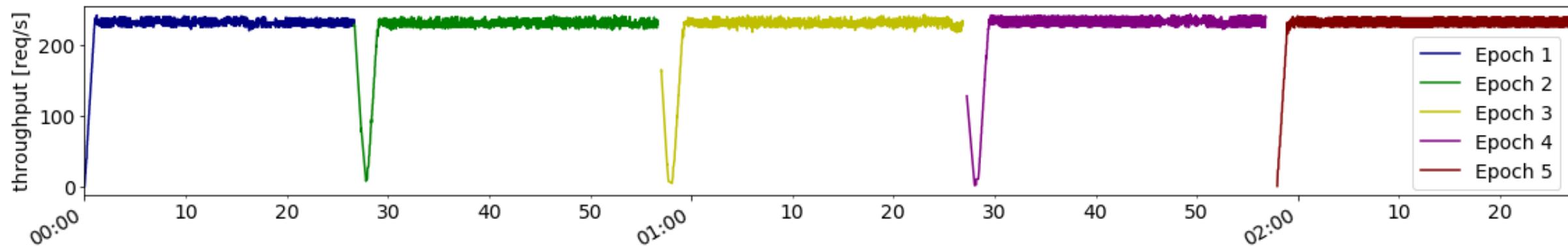
# Setup



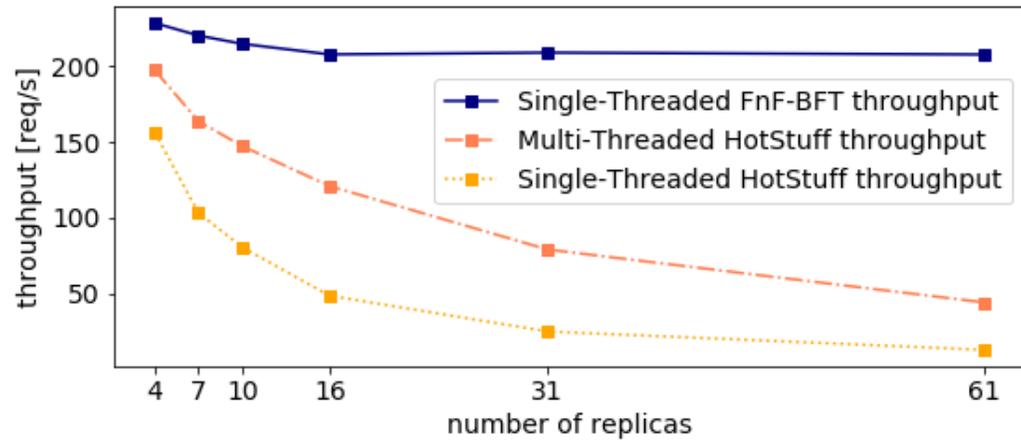
# Setup



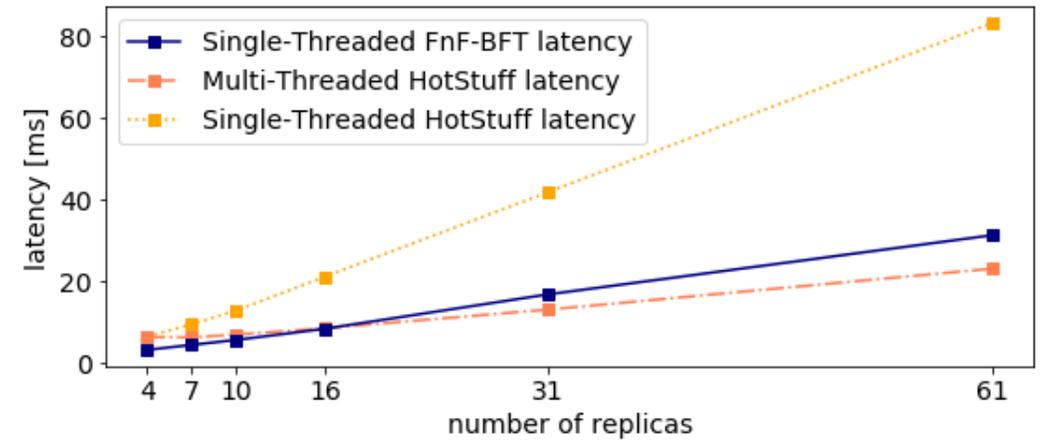
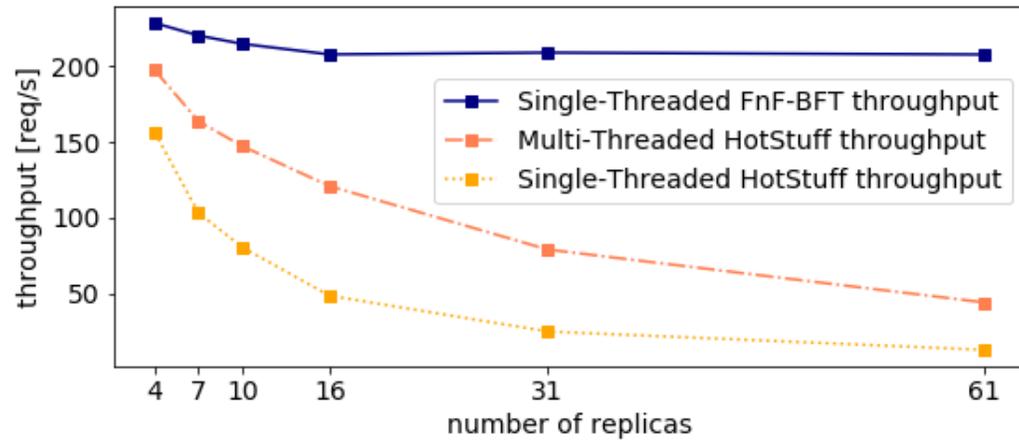
# Performance

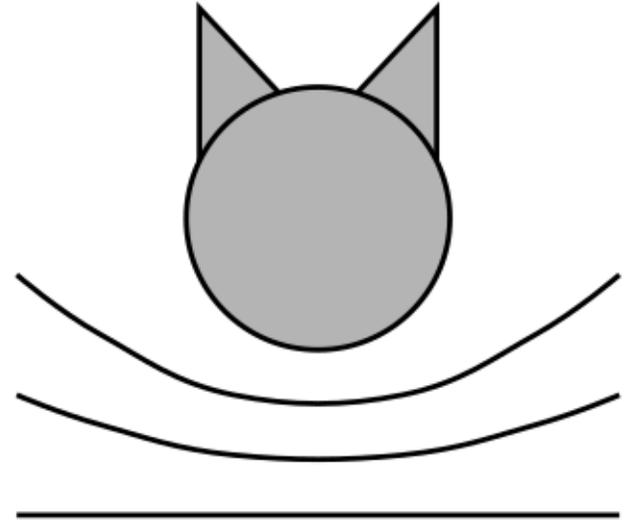
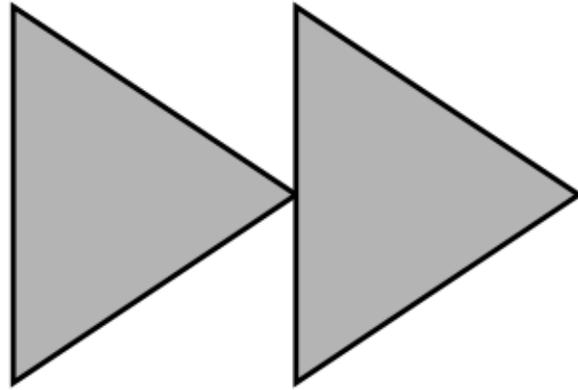
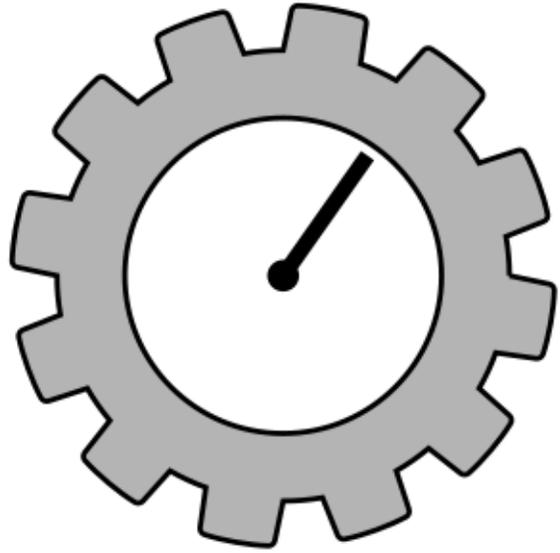


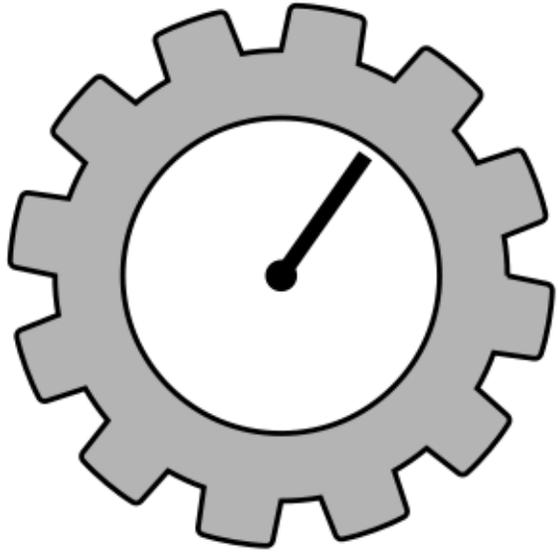
# Performance



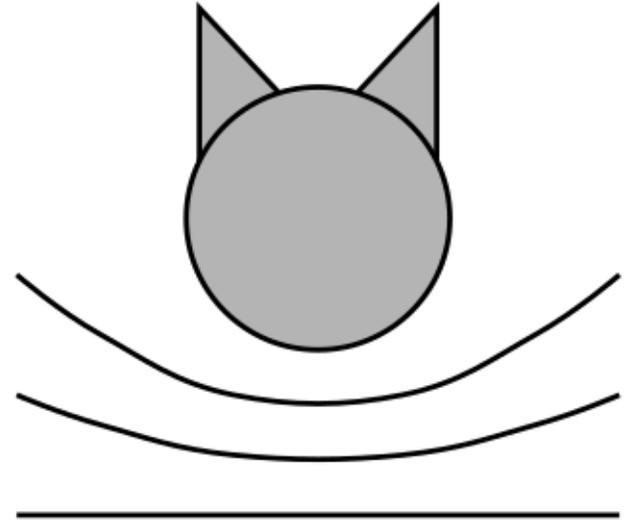
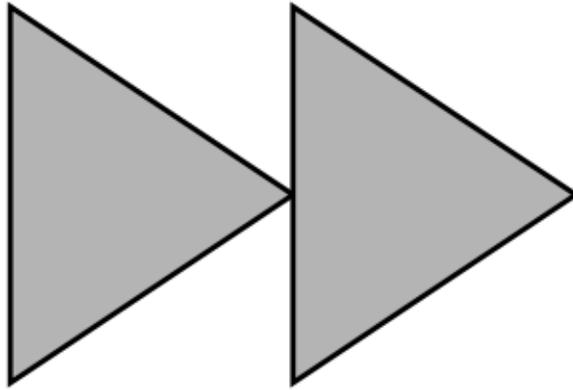
# Performance

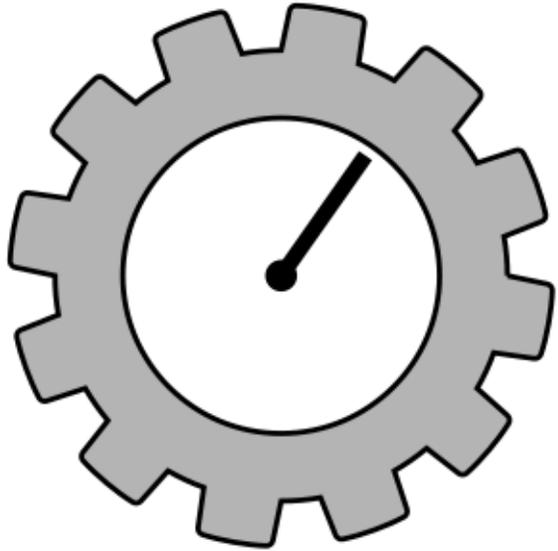




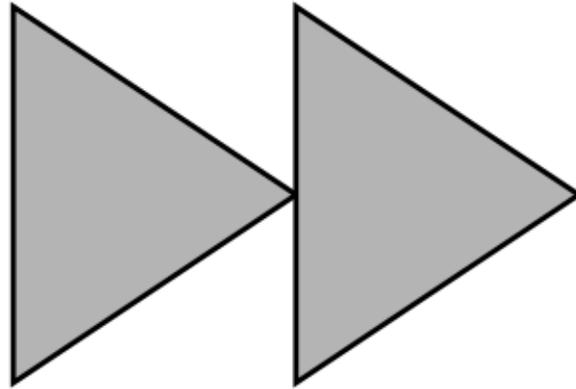


amortized request cost:  
 $O(n)$

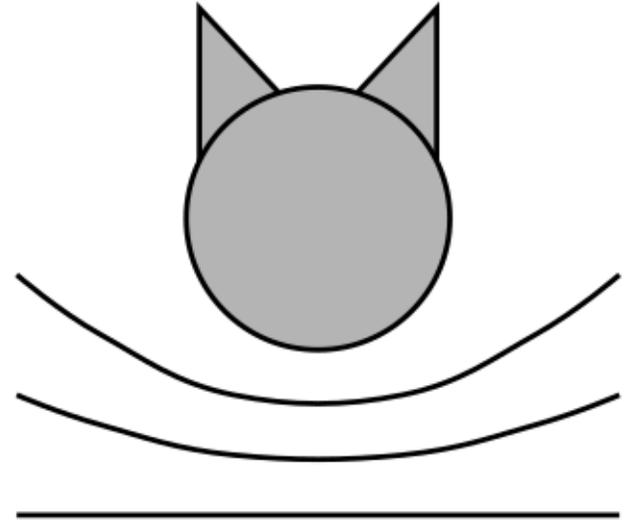


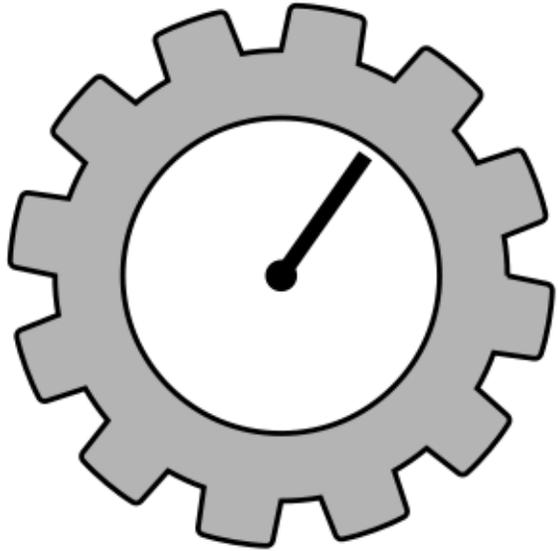


amortized request cost:  
 $O(n)$

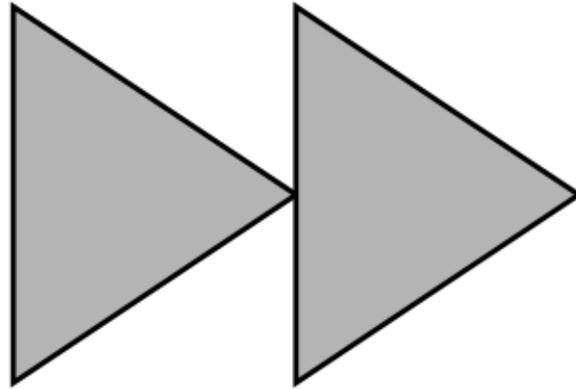


speedup:  
 $\Theta(n)$

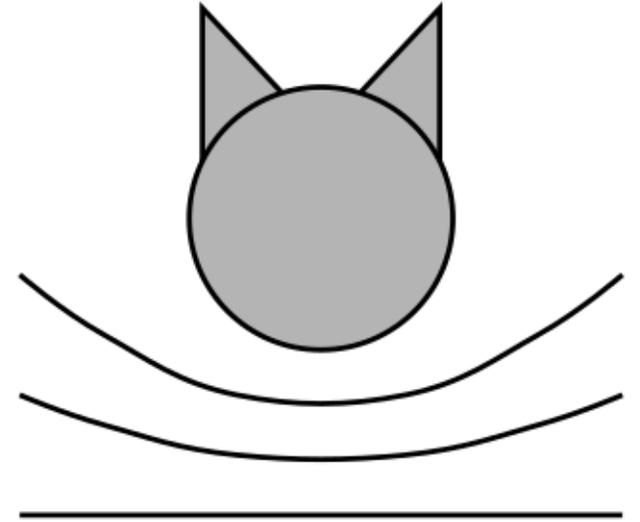




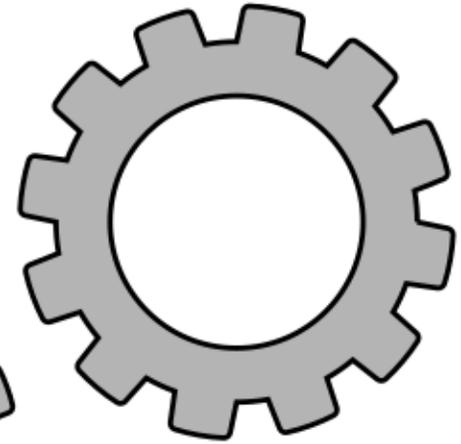
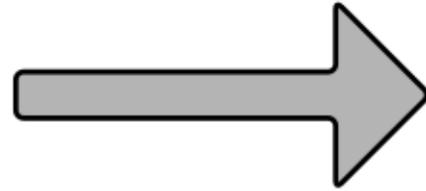
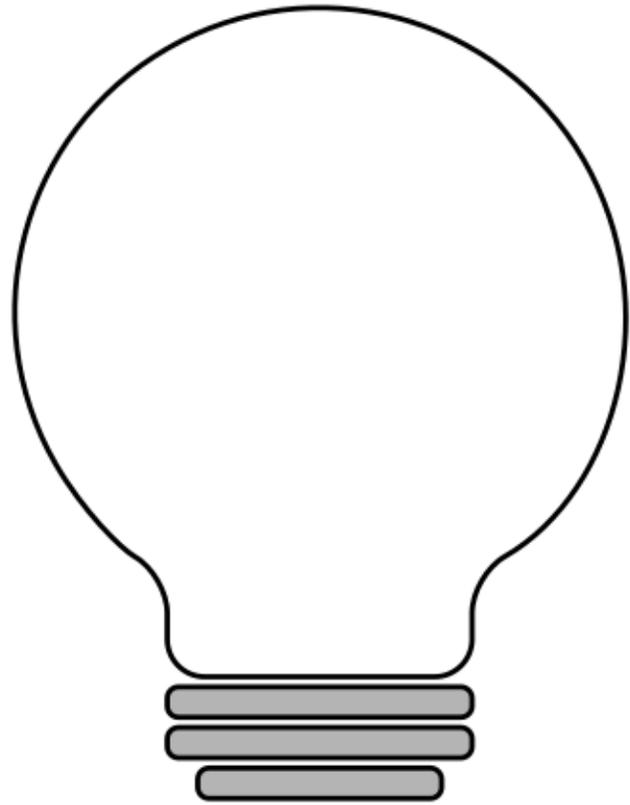
amortized request cost:  
 $O(n)$

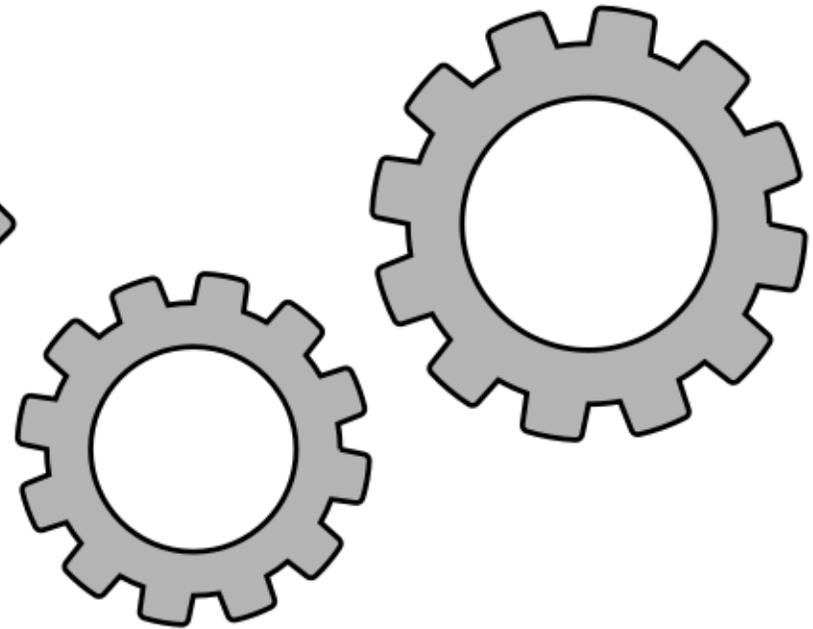
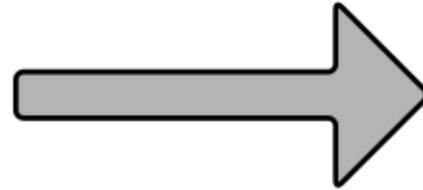
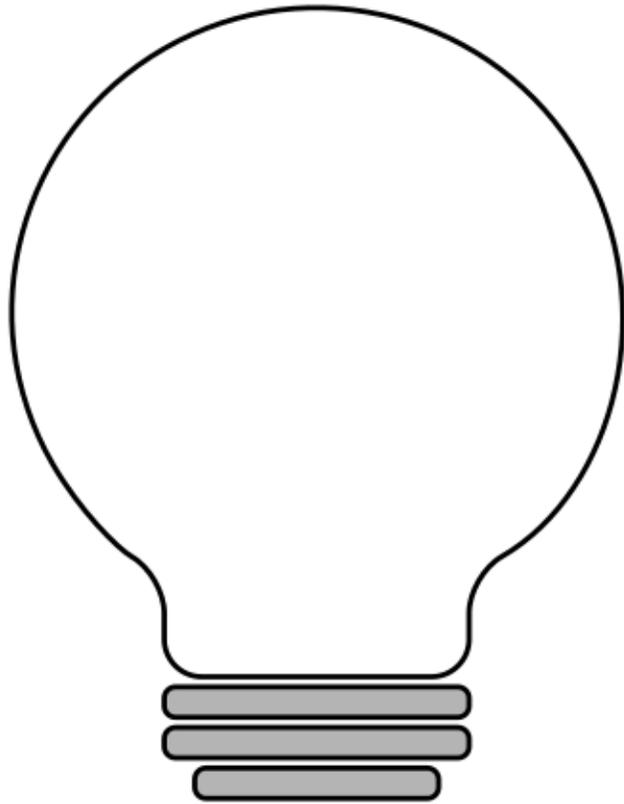


speedup:  
 $\Theta(n)$



effective utilization:  
 $\geq \frac{16}{27}$





significantly improved  
scaling capabilities

Thank You!  
Questions & Comments?



Zeta Avarikioti, Lioba Heimbach, Roland Schmid, Laurent Vanbever, Roger Wattenhofer, Patrick Wintermeyer  
ETH Zurich – Distributed Computing – [www.disco.ethz.ch](http://www.disco.ethz.ch)

# Configuration parameters

Configuration Parameter	Setting
Requests per block	1
Threads per replica	1
Threads per client	4
Epoch timeout	30s
No progress timeout	2s
Blocks per checkpoint ( $K$ )	50
Watermark window size ( $2 * K$ )	100
Initial epoch watermark bounds	10000

Table 1. FNF-BFT configuration parameters used across all experiments.