## Clairvoyant Mechanisms for Online Auctions

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## Online Auctions




Valuation




Preemption Price



## What if we knew the future?

| $\$ 1$ | $\$ 10$ | $\$ 5$ | $\$ 200$ | $\$ 1500$ | $\$ 2500$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $\$ 10$ | $\$ 2000$ | $\$ 200$ | $\$ 750$ | $\$ 5000$ | $\$ 3000$ |

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## Clairvoyant Model



Calculate in every round $r: \frac{\operatorname{opt}(r)}{\operatorname{gain}(r)}=\frac{1}{1}$
Score for a set $S$ of accepted bidders: $\max _{r} \frac{\operatorname{opt}(r)}{\operatorname{gain}(r)}$

## Clairvoyant Model



Calculate in every round $r: \frac{\operatorname{opt}(r)}{\operatorname{gain}(r)}=\frac{10}{1}$
Score for a set $S$ of accepted bidders: $\max _{r} \frac{\operatorname{opt}(r)}{\operatorname{gain}(r)}$

## Clairvoyant Model



Calculate in every round $r: \frac{\operatorname{opt}(r)}{\operatorname{gain}(r)}=\frac{10}{-5}$
Score for a set $S$ of accepted bidders: $\max _{r} \frac{\operatorname{opt}(r)}{\operatorname{gain}(r)}$

## Clairvoyant Model



Calculate in every round $r: \frac{\operatorname{opt}(r)}{\operatorname{gain}(r)}=\frac{200}{1}$
Score for a set $S$ of accepted bidders: $\max _{r} \frac{\operatorname{opt}(r)}{\operatorname{gain}(r)}$

## Clairvoyant Model



Calculate in every round $r: \frac{\operatorname{opt}(r)}{\operatorname{gain}(r)}=\frac{1500}{1500-10}$
Score for a set $S$ of accepted bidders: $\max _{r} \frac{\operatorname{opt}(r)}{\operatorname{gain}(r)}$

## Clairvoyant Model



Calculate in every round $r: \frac{\operatorname{opt}(r)}{\operatorname{gain}(r)}=\frac{2500}{1500-10}$
Score for a set $S$ of accepted bidders: $\max _{r} \frac{\operatorname{opt}(r)}{\operatorname{gain}(r)}$

## Clairvoyant Model



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## Clairvoyant Model

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Difficulty $\Delta=\min _{S} \max _{r} \frac{\operatorname{opt}(r)}{\operatorname{gain}(S, r)}$

## $\Delta$ Online Mechanisms



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|  |  |  |  |  |

Difficulty $\Delta=\min _{S} \max _{r} \frac{\operatorname{opt}(r)}{\operatorname{gain}(S, r)}$

## $\Delta$ Online Mechanisms

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Difficulty $\Delta=\min _{S} \max _{r} \frac{\operatorname{opt}(r)}{\operatorname{gain}(S, r)}$

## $\Delta$ Online Mechanisms



Difficulty $\Delta=\min _{S} \max _{r} \frac{\operatorname{opt}(r)}{\operatorname{gain}(S, r)}$

## Theorem

An online mechanism that knows $\Delta$ can be $\Delta^{5}$ competitive

Lower Bound

| 1 | $\Delta^{0}$ | $\Delta^{1}$ | $\Delta^{2}$ | $\Delta^{3}$ | $\Delta^{4}$ | $\Delta^{5}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | $\Delta^{0+2}$ | $\Delta^{1+2}$ | $\Delta^{2+2}$ | $\Delta^{3+2}$ | $\Delta^{4+2}$ | $\Delta^{5+2}$ |

Lower Bound


## Lower Bound


$v_{j}$

## Lower Bound


$v_{j} \Delta$
$v_{j} \Delta^{2}$

$v_{j} \Delta^{2}$
$v_{j} \Delta^{2}$
$v_{j} \Delta^{1000}$

$v_{j} \Delta^{1337}$ 1
$\mathrm{v}_{j} \Delta^{1336}$

$\Delta^{4+2}$

## Lower Bound




## Lower Bound



Theorem
An online mechanism that knows $\Delta$ can be $\Delta^{5}$ competitive; this is optimal.

| $\$ 10$ | $\$ 100$ | $\$ 5$ | $\$ 200$ | $\$ 1500$ | $\$ 2500$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $\$ 10$ | $\$ 2000$ | $\$ 200$ | $\$ 750$ | $\$ 5000$ | $\$ 3000$ |

Difficulty $\Delta=\min _{S} \max _{r} \frac{\operatorname{opt}(r)}{\operatorname{gain}(S, r)}$

## Theorem

An online mechanism that knows $\Delta$ can be $\Delta^{5}$ competitive; this is optimal.

