Wireless Network Algorithms Looking Back & Moving Forward



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Theory & Practice

PODC STOC FOCS ICALP SPAA ESA SODA

EC

OSDI SenSys ICML CHI IPSN NeurIPS Mobicom SIGCOMM

SPECTRE : Spectral Conditioning Helps to Overcome the Expressivity Limits of One-shot Graph Generators

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Automating Rigid Origami Design

Jeremia Geiger, Karolis Martinkus, Oliver Richter, Roger Wattenhofer



International Symposium on Algorithmics of Wireless Networks

Ad Hoc Networks Autonomous Mobile Robots Communication Protocols Complexity and Computability **Computational Models** Data Aggregation and Fusion Dynamic Networks, Temporal Graphs Energy Management, Power Saving Fault Tolerance and Dependability Game Theoretic Aspects Infrastructure Discovery Internet of Things Localization Medium Access Control

Mobility and Dynamics Obstacle Avoidance Pattern Formation, Experimental Analysis Population Protocols, Swarm Computing **Resource Efficiency RFID Algorithms Routing and Data Propagation** Self-stabilization, Self-* Properties Sensor Networks Systems and Testbeds Time Synchronization **Topology Control** Tracking Virtual Infrastructures

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• Synchronous – all switchable nodes

 \rightarrow possibly infinite



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• Synchronous – all switchable nodes

 \rightarrow possibly infinite

• Sequential – one node at a time

 \rightarrow always stabilizes!

Weighted Minority Process



Weighted Minority Process



Weighted Minority Process



• Basic – for any improvement



• Basic – for any improvement



- Basic for any improvement
- Proportional conflict with *e* portion of (weighted) neighborhood



- Basic for any improvement
- **Proportional** conflict with *e* portion of (weighted) neighborhood



Upper bound:

• Sequential model: $O(2^n)$

Upper bound:

• Sequential model: $O(2^n)$


Stabilization time

Upper bound:

• Sequential model: $O(2^n)$



Stabilization time

Upper bound:

• Sequential model: $O(2^n)$

Lower bound?

Majority vs. Minority

Majority process (take most frequent color)



Minority process

(take least frequent color)



steps: $2^{\Theta(n)}$

steps: ?

Majority vs. Minority

Majority process (take most frequent color)



Minority process

(take least frequent color)



steps: $2^{\Theta(n)}$

steps: 2⁽ⁿ⁾





Lower weightsMore switches











 $Q=\frac{5}{6}$



 $Q=\frac{5}{6}$



e =





e =









e =

























• Adversarial: a specific sequence

• Adversarial: a specific sequence



• Adversarial: a specific sequence

Benevolent



• Adversarial: a specific sequence

• **Benevolent:** every sequence

• Adversarial: a specific sequence

• Benevolent: every sequence

 \rightarrow allow only one possible sequence!

Benevolent case


Benevolent case – logical gates



Benevolent case – logical gates



Benevolent case – logical gates



More Variants by [Papp, W]

Unweighted $\rightarrow \Theta(n^2)$

Random initialization $\rightarrow \Theta(n^2)$

Random init & proportional \rightarrow depends on ϱ

More general model
$$\rightarrow f(\lambda) := \max_{\varphi \in (0, \frac{1-\lambda}{2}]} \frac{\log\left(\frac{1-\varphi}{\lambda+\varphi}\right)}{\log\left(\frac{1-\varphi}{\varphi}\right)}$$

Sensor Networks

Algorithms for Sensor NetworksWhat Is It Good For?!



[W, Algosensors 2008!]

[PermaSense]

Efficiency and Reliability



Energy Efficiency



[Burri, von Rickenbach, W]



Energy Efficiency



Energy Efficiency



duty cycling, wake up e.g. every 10 seconds parent synchronizes children no network wide synchronization mean energy consumption: 0.066mW, 10y battery

Reliability



Reliability



nodes send beacons to reconnect orphans collisions are explicitly accepted availability & reliability: 99% to 99.999%

Wireless vehicle detection systems for outdoor parking lots

[tinynode]

Where's the Algorithmic Theory?

"no network wide synchronization"

Time Synchronization

Network Synchronization is Tough

Network Synchronization



Tree Based Protocols



Synchronization Error	FTSP	PulseSync
Average (t > 2000s)	23.96 µs	4.44 μs
Maximum (t > 2000s)	249 μs	38 µs

Error with Distance



Neighbor Synchronization



Neighbor Synchronization?





Tree-based Algorithms e.g. FTSP Neighborhood Algorithms e.g. GTSP

Theorem: Neighbor Sync is Somewhat Tough

Model: Drift & Jitter



Reasonable Time Must Behave!



no stopping

no jumping




















Example: Neighbor Sync is Hard

sync to fastest neighbor message delay = 10



Sync To Fastest Neighbor: Local Skew Can Be Diameter Average of Neighbors: Local Skew Can Be Diameter Squared

Better Protocol?

Reminder: Drift & Jitter



Theorem: Neighbor Sync is <mark>Somewhat</mark> Tough

neighbor sync error = log diameter lower bound: difficult proof matching upper bound: not trivial as well

[Lenzen, Locher, W, JACM]

"The Future"





Starlink

-

Global Positioning System









Atomic Clock Inside Transmit Time + "Position" + ...



Tracking

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Energy Problem





Delay Problem



Position and time only 30s after turning camera on

Delay Problem



Better: Record 1ms of raw GPS data when taking picture

Coarse Time Navigation



Coarse Time Results



15 minute duty cycle18 μW mean power2 years with coin cell

[Eichelberger, von Hagen, W]

GPS Spoofing

[Eichelberger, von Hagen, W, 2019]

Marca









Indoor «GPS»







[Eichelberger, Luchsinger, Tanner, W]

-light Awar

Does it work?



Still many challenges...




Questions? Comments?

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