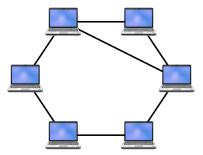
Randomness vs. Time in Anonymous Networks

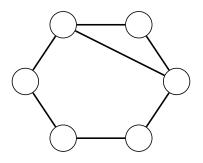
Jochen Seidel Jara Uitto

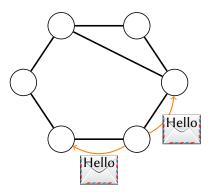
Roger Wattenhofer

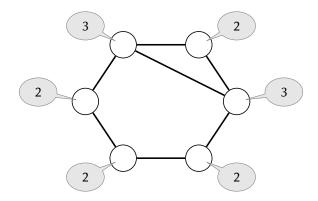
ETH Zurich – Distributed Computing Group – www.disco.ethz.ch



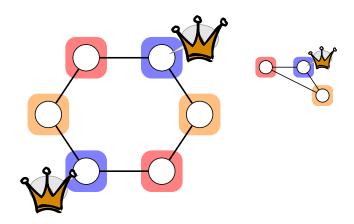




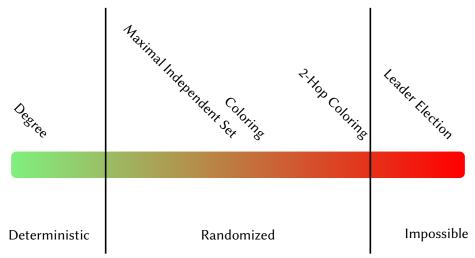


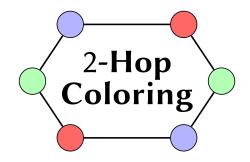


Leader Election

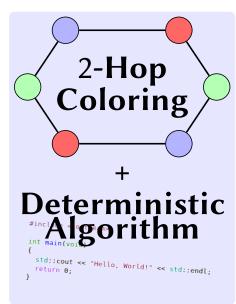


Computability





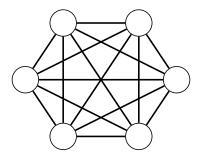
mersenne twister engine«UIntType,w,n,m,r,a,u,d,s,b,t,c,l,f>::twist() const UIntType upper mask = (-static cast<UIntType>(0)) << r;</pre> const UIntType lower mask = -upper_mask; const std::size t unroll factor = 6; const std::size t unroll extral = (n-m) % unroll factor; const std::size t unroll extra2 = (m-1) % unroll factor; // split loop to avoid costly modulo operations { // extra scope for MSVC brokenness w.r.t. for scope for(std::size t j = 0; j < n-m-unroll extral; j++) {</pre> UIntType y = (x[j] & upper mask) | (x[j+1] & lower mask); $x[j] = x[j+m] \land (y >> 1) \land ((x[j+1]61) * a);$ Raindom unter set i set ask; gorithm for(std::size t j = n-1-unroll extra2; j < n-1; j++) { UIntType y = (x[j] & upper mask) | (x[j+1] & lower mask); 1 $x[j] = x[j - (n-m)] \land (y >> 1) \land ((x[j+1]\delta 1) + a);$ UIntType $y = (x[n \cdot 1] & upper mask) | (x[0] & lower mask);$ $x[n-1] = x[n-1] \land (y >> 1) \land ((x[0]\delta 1) * a);$ i = 0;std::size t w, std::size t m, std::size t m, std::size t r, template<class UIntType, UIntType a, std::size t u, UIntType d, std::size t S,



2**-Hop** // extra scope for MSVC brokenness w.r.t. for scope = θ; j < n-m-unroll extral; j++) { & upper mask) | (x[j+1] & lower mask); Coloring iny ⁼ random bitisistic 1 std::cout << "Hello, World!" << std::endl;</pre>

2-Hop // extra scope for MSVC brokenness w.r.t. for scope = θ; j < n-m-unroll extral; j++) { & upper mask) | (x[j+1] & lower mask); Coloring iny randor 1 ministic std::cout << "Hello, World!" << std::endl;</pre>

2**-Hop** // extra scope for MSVC brokenness w.r.t. for scope = 0; j < n-m-unroll extral; j++) { Coloring any randoi 1 n b ISTIC How fast?cout << "Hello, World!" << std::endl; urn 0;



log n





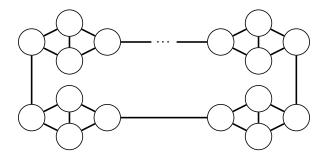
. . .



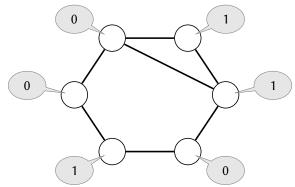


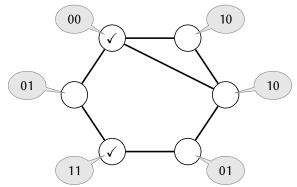
log n

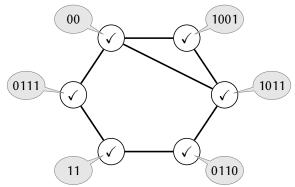
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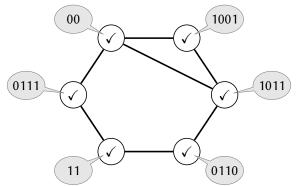


log n

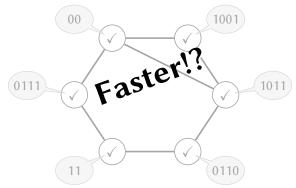


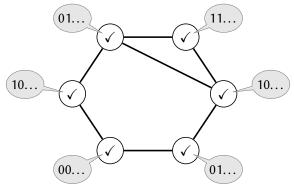




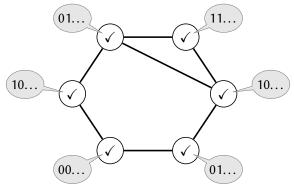


log log *n* rounds, log *n* bits (w.h.p. & in expectation)



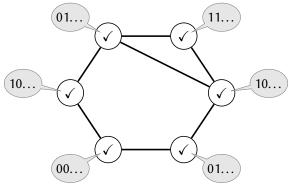


∞ random bits \Rightarrow 1 round

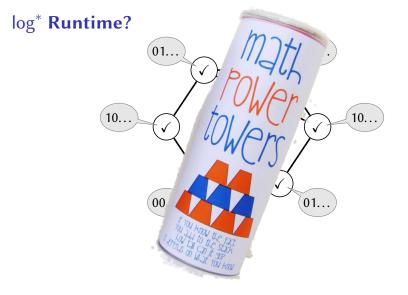






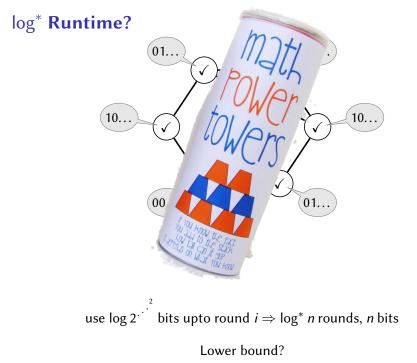


Idea: Use the inverse of log*

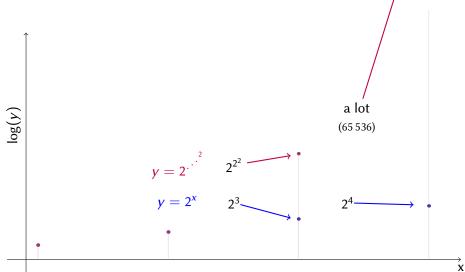


Idea: Use the inverse of log*

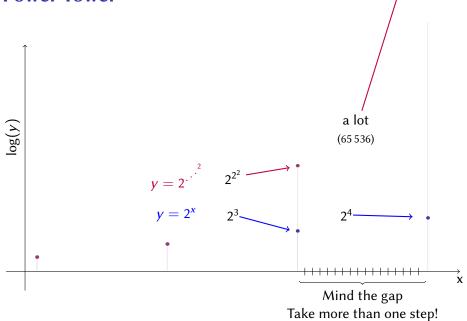




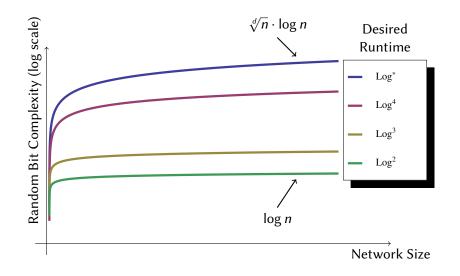
Power Tower







Time vs. Random Bit Complexity



Time vs. Random Bit Complexity

Theorem

Anonymous networks can be 2-hop colored in $O(d \cdot f(n))$ rounds.¹

The achieved random bit complexity is

$$O\left(\sqrt[d]{\frac{\lceil \log f^{-1}(f(n)+1)\rceil}{\lceil \log f^{-1}(f(n))\rceil}} \cdot \log n\right)$$

¹With high probability and in expectation. For reasonable desired runtime f, i.e., between log^{*} and log log.

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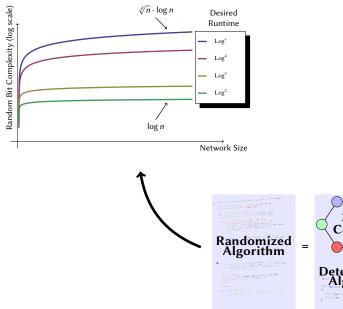
$$O\left(\sqrt[d]{\frac{\lceil \log f^{-1}(f(n)+1)\rceil}{\lceil \log f^{-1}(f(n))\rceil}} \cdot \log n\right)$$

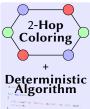
Theorem

This is asymptotically optimal.

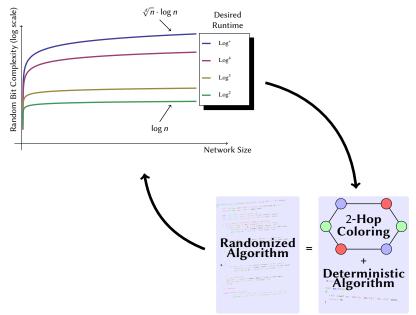
¹With high probability and in expectation. For reasonable desired runtime f, i.e., between \log^* and $\log \log$.

Recap





Recap



Thank You