

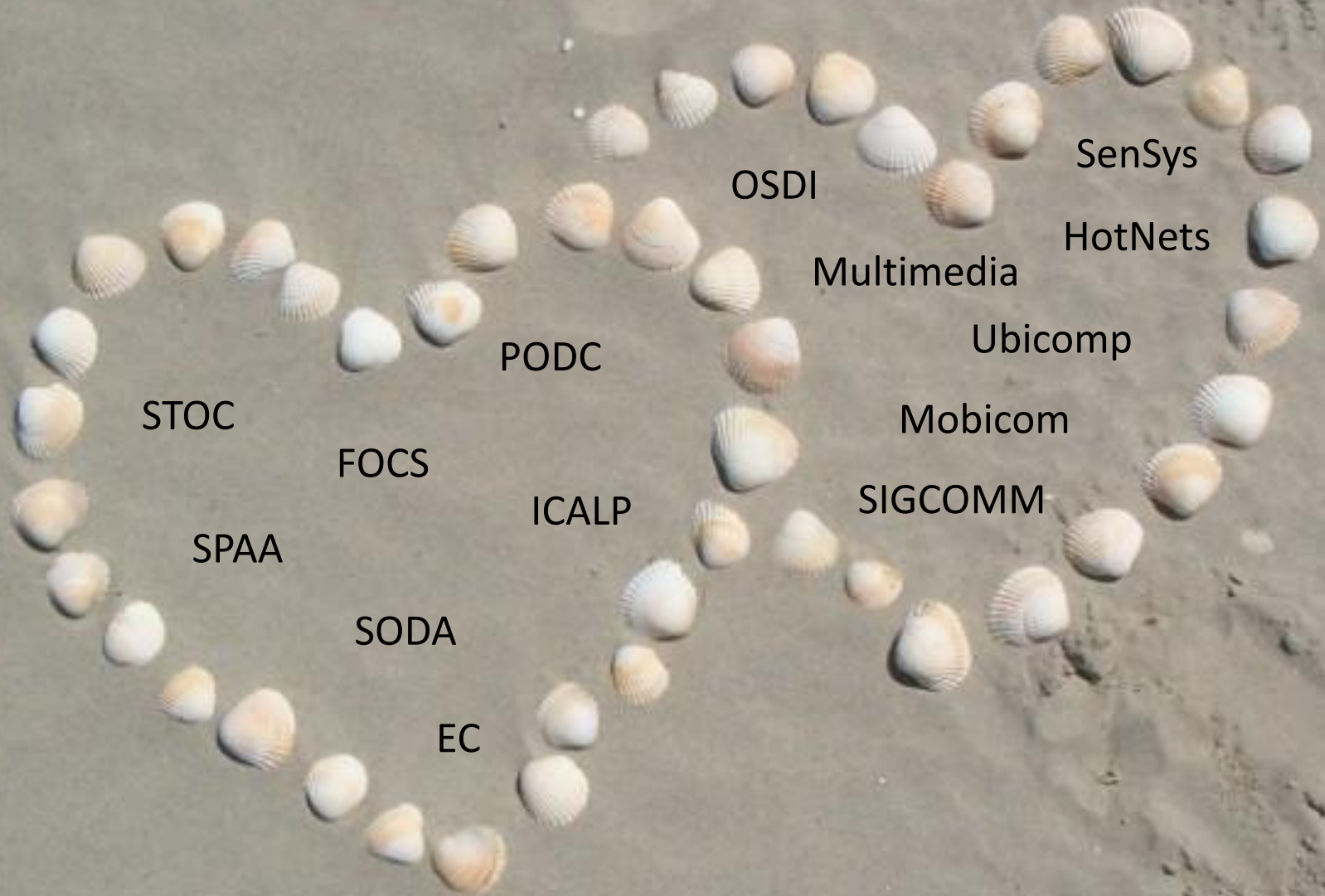
Sensor Networks

Where Theory Meets Practice



Roger Wattenhofer

Theory Meets Practice



STOC

SPAA

SODA

EC

FOCUS

SICALP

PODC

OSDI

Multimedia

Mobicom

SIGCOMM

Ubicomp

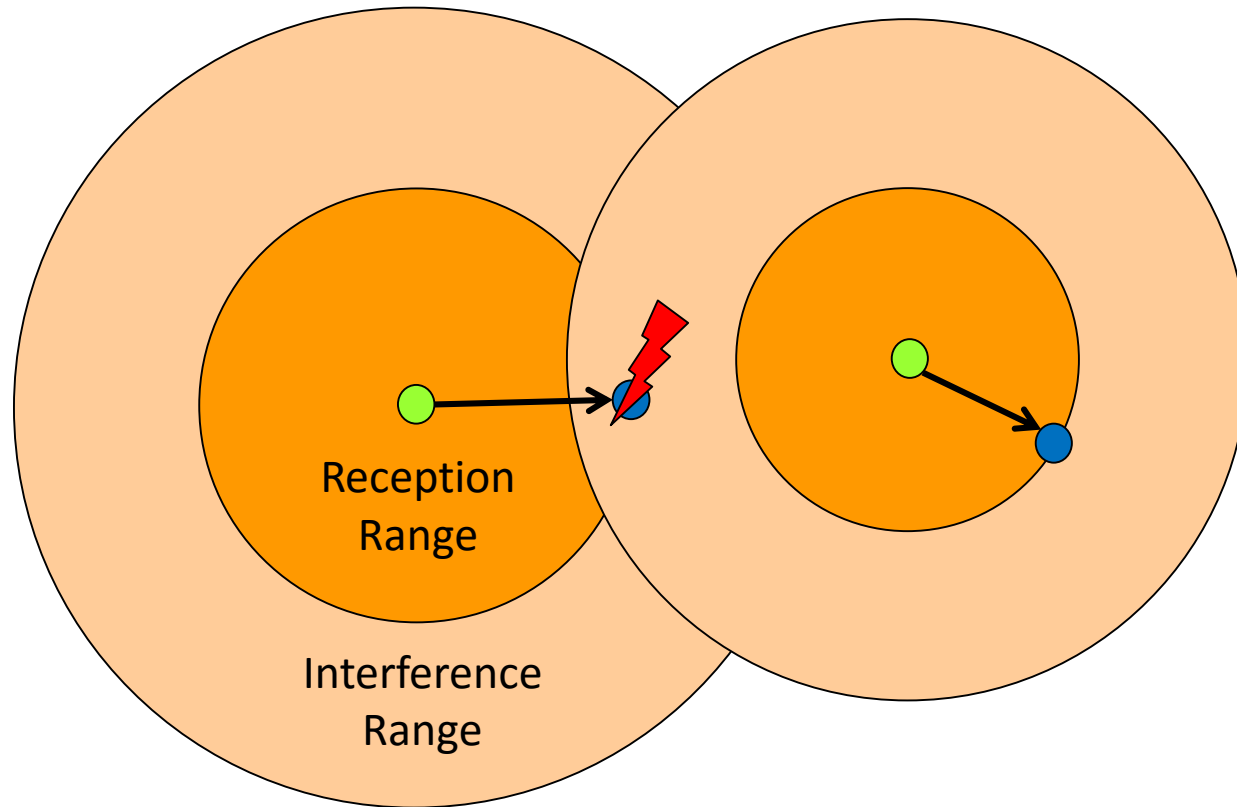
HotNets

SenSys

Wireless Communication?

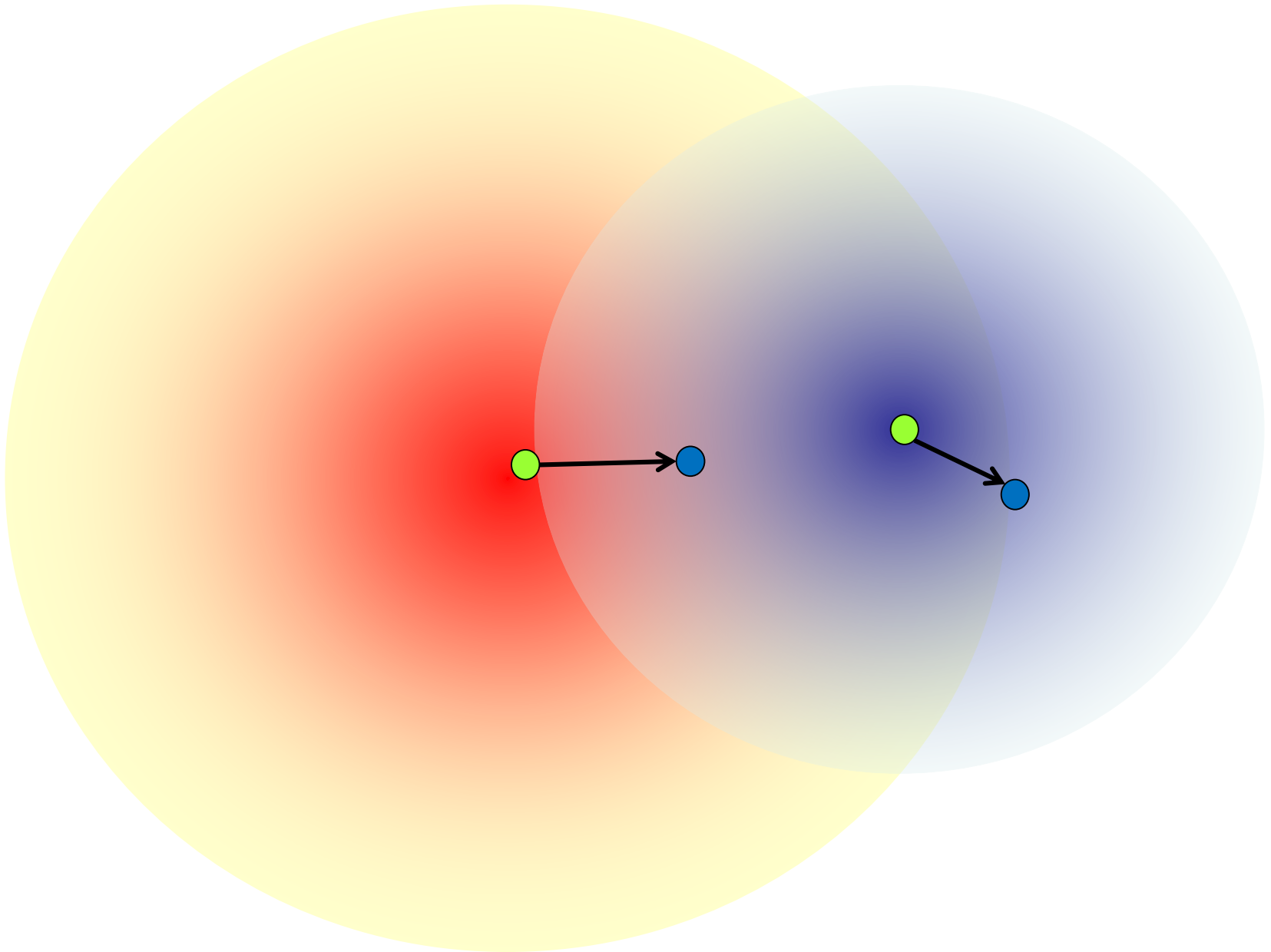
Capacity!

Protocol Model



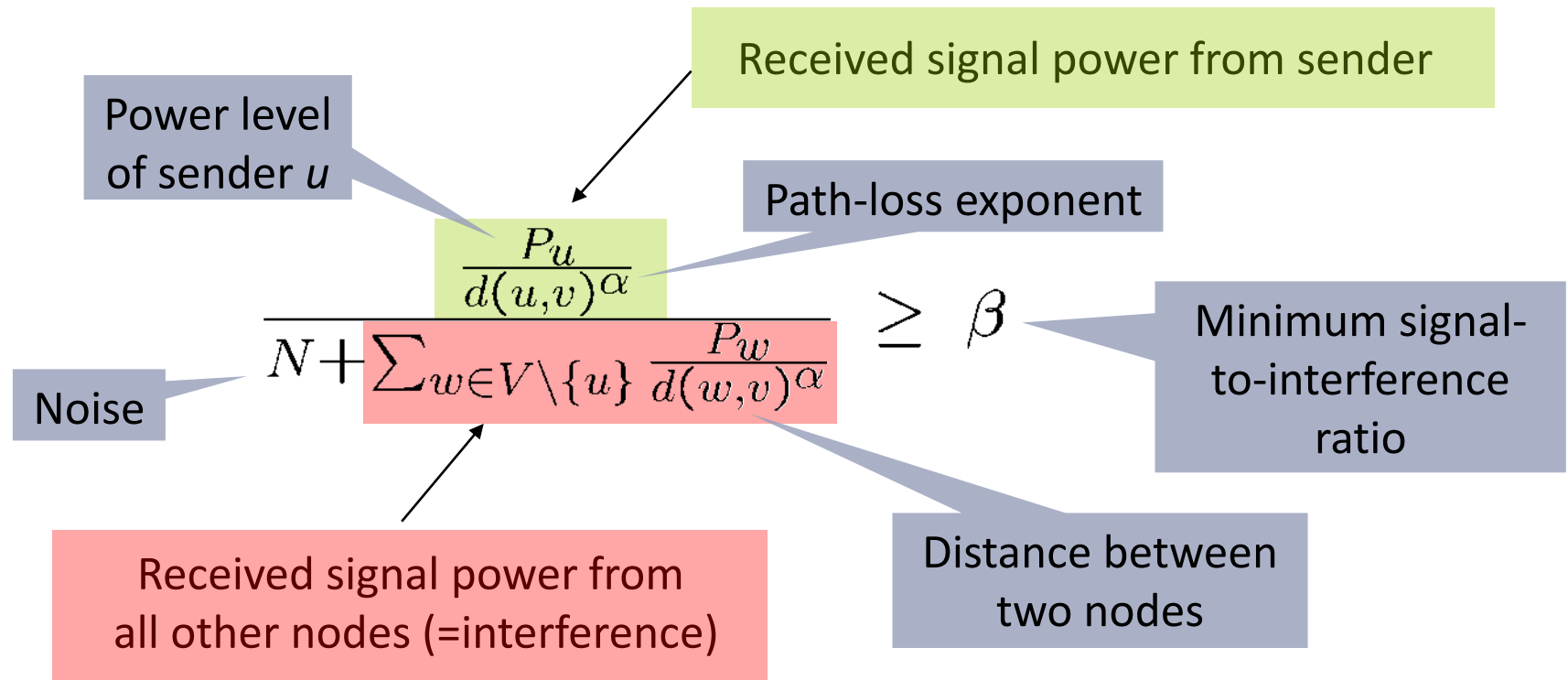


Physical (SINR) Model

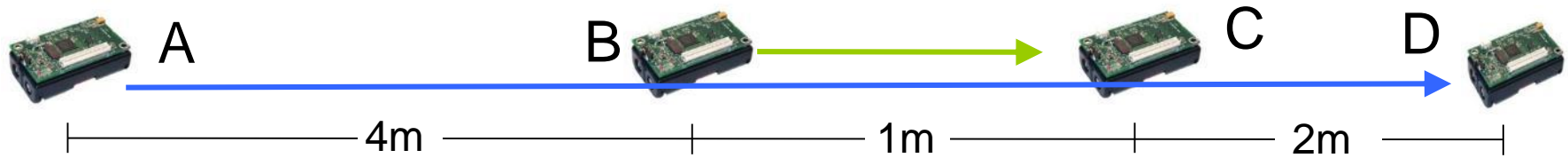




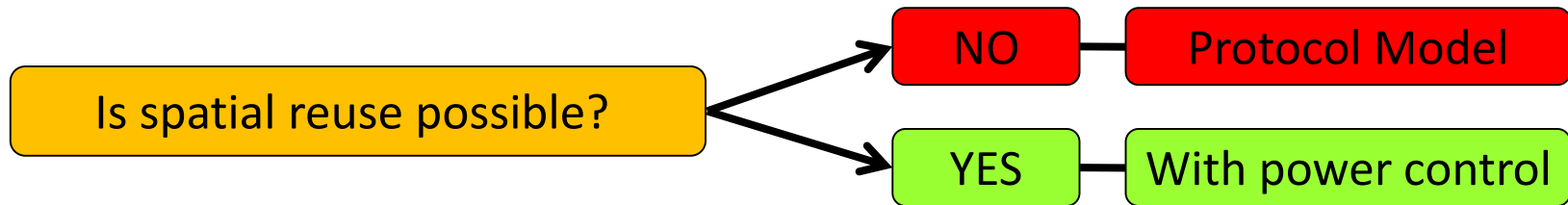
Signal-To-Interference-Plus-Noise Ratio (SINR) Formula



Example: Protocol vs. Physical Model





Assume a **single frequency** (and no fancy decoding techniques!)



Let $\alpha=3$, $\beta=3$, and $N=10\text{nW}$

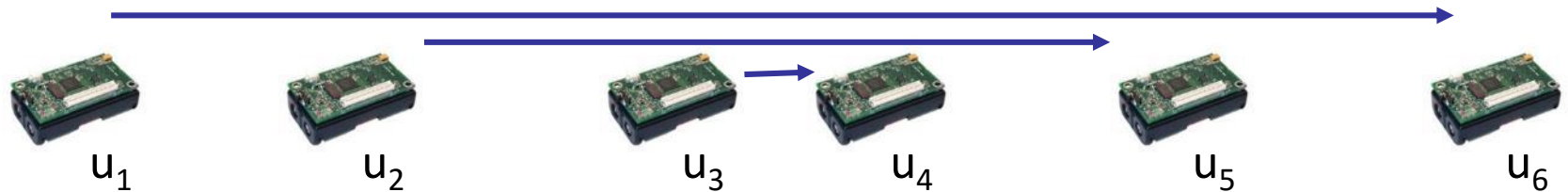
Transmission powers: $P_B = -15\text{ dBm}$ and $P_A = 1\text{ dBm}$

SINR of A at D:
$$\frac{1.26\text{mW}/(7\text{m})^3}{0.01\mu\text{W} + 31.6\mu\text{W}/(3\text{m})^3} \approx 3.11 \geq \beta$$
 

SINR of B at C:
$$\frac{31.6\mu\text{W}/(1\text{m})^3}{0.01\mu\text{W} + 1.26\text{mW}/(5\text{m})^3} \approx 3.13 \geq \beta$$
 

This works in practice

... even with very simple hardware



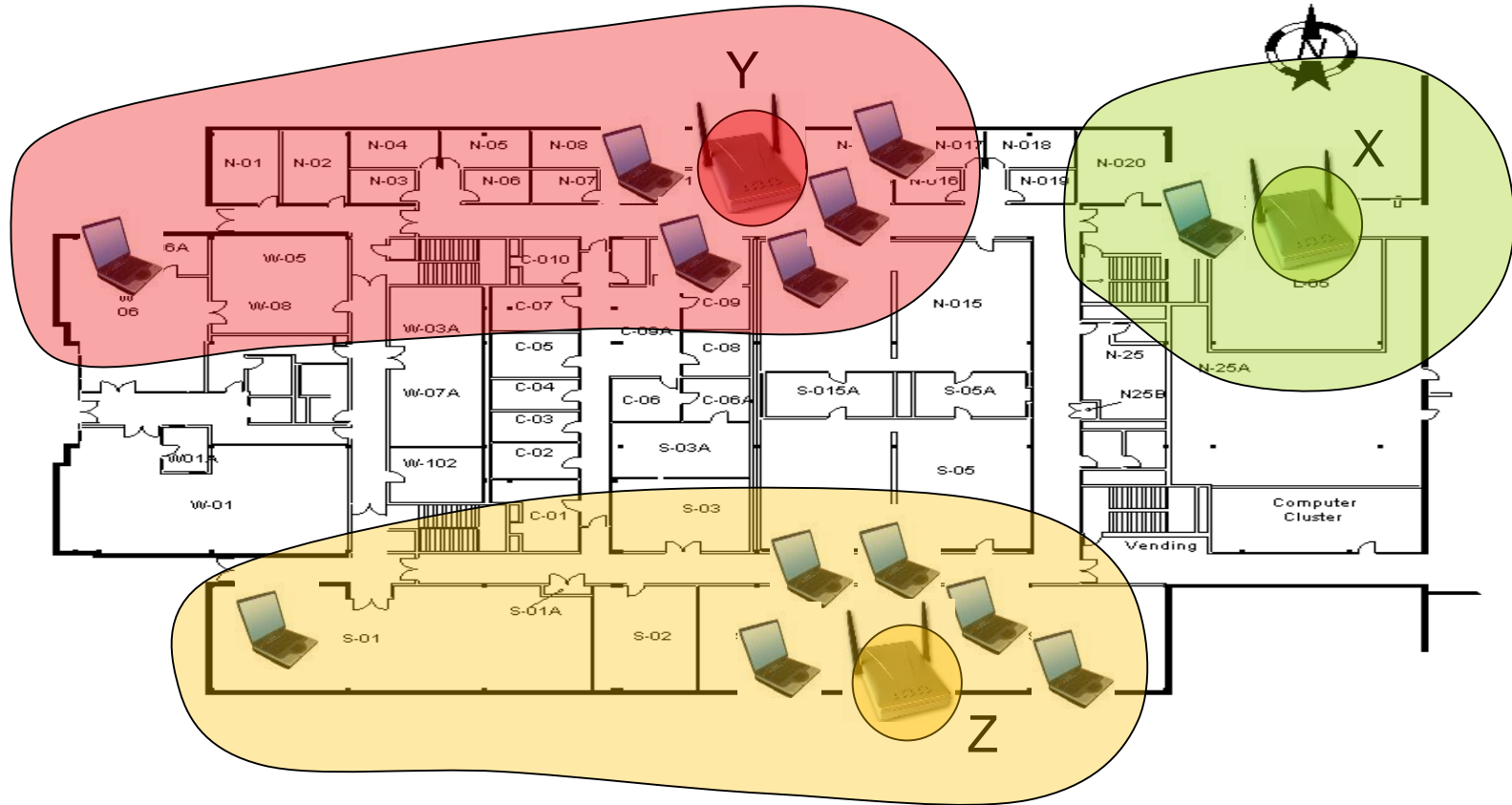
Time for transmitting 20'000 packets:

	Time required	
	standard MAC	"SINR-MAC"
Node u_1	721s	267s
Node u_2	778s	268s
Node u_3	780s	270s

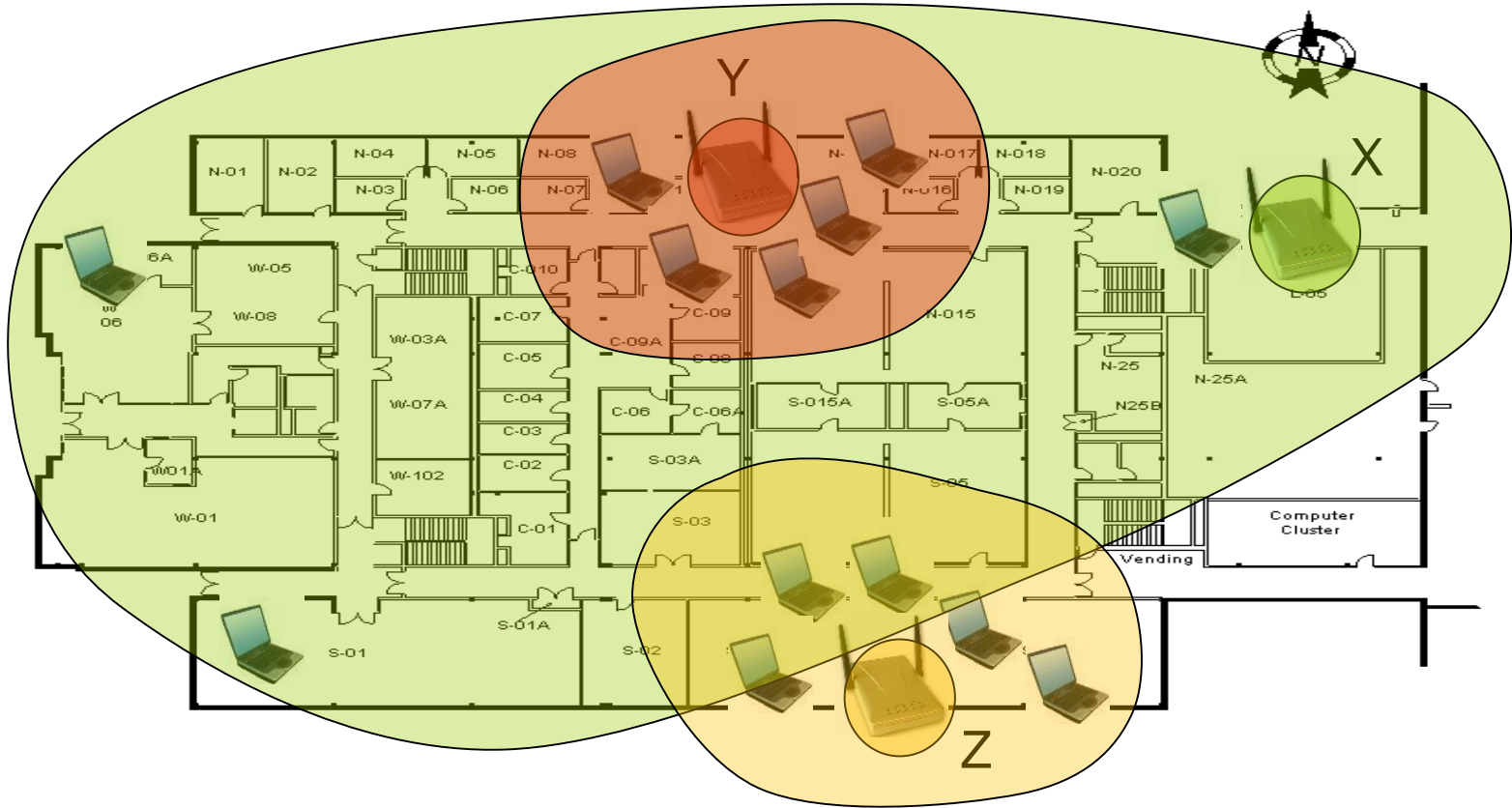
	Messages received	
	standard MAC	"SINR-MAC"
Node u_4	19999	19773
Node u_5	18784	18488
Node u_6	16519	19498

Speed-up is almost a factor 3

Possible Application – Hotspots in WLAN



Possible Application – Hotspots in WLAN



The Capacity of a Network

(How many concurrent wireless transmissions can you have)

Convergecast Capacity in Wireless Sensor Networks

[Moscibroda, W, 2006]

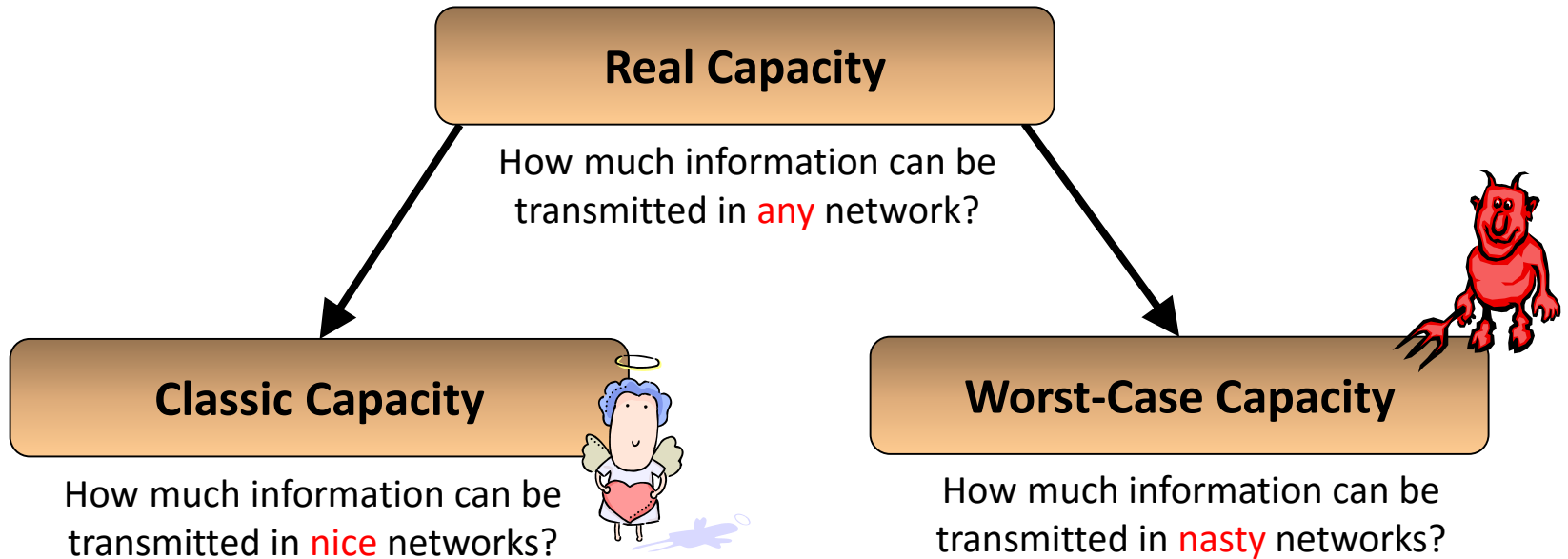
Worst-Case Capacity

[Giridhar, Kumar, 2005]

Classic Capacity

Topology Model/Power	Max. rate in arbitrary, worst-case deployment	Max. rate in random, uniform deployment
Protocol Model	$\Theta(1/n)$	$\Theta(1/\log n)$
Physical Model (power control)	$\Omega(1/\log^3 n)$	$\Omega(1/\log n)$

Capacity of a Network



Core Capacity Problems

Given a set of **arbitrary** communication links

One-Shot Problem

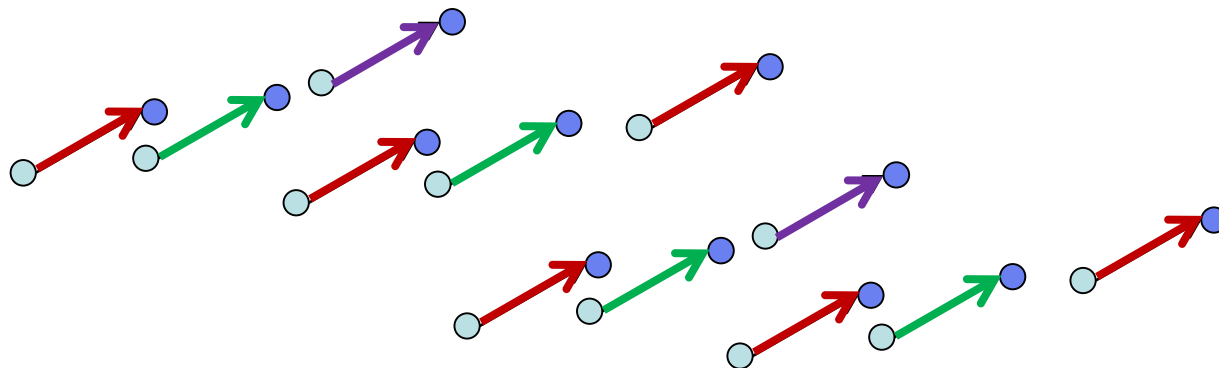
Find the **maximum size** feasible subset of links

$O(1)$ approximations for uniform power [Goussevskiaia, Halldorsson, W, 2009 & 2014] as well as arbitrary power [Kesselheim, 2011]

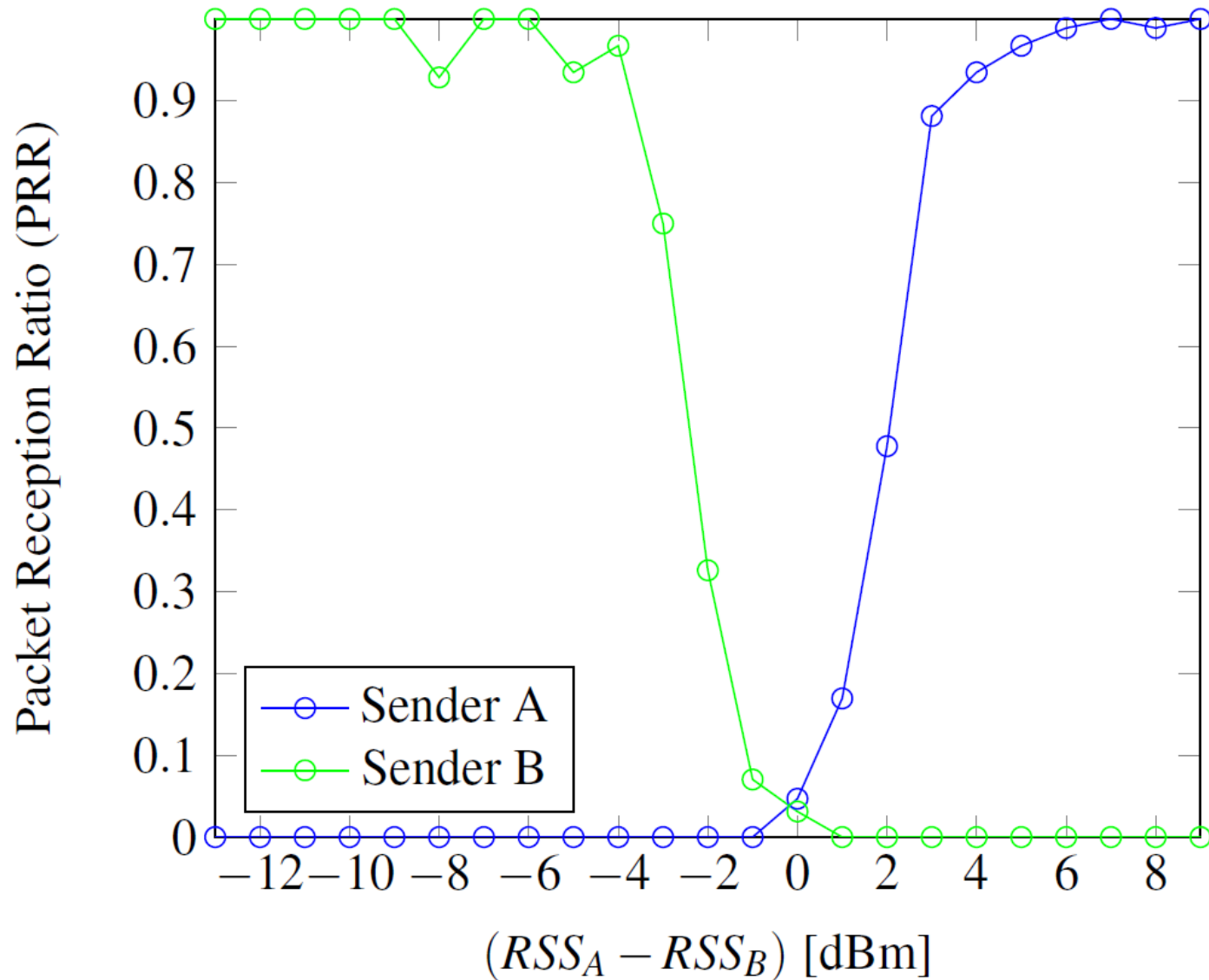
Scheduling Problem

Partition the links into fewest possible slots, to **minimize time**

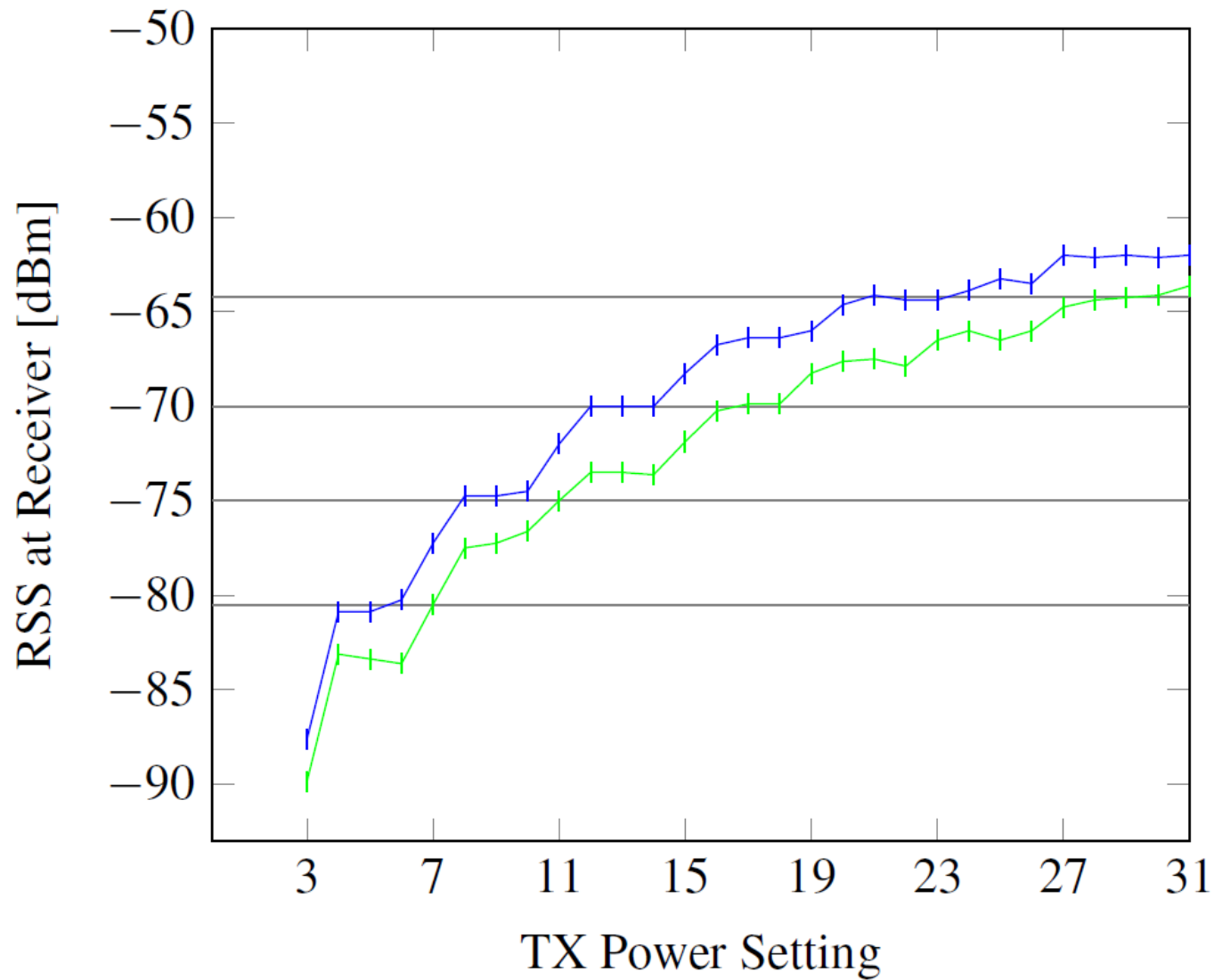
Open problem: Only $O(\log n)$ approximation using the one-shot subroutine



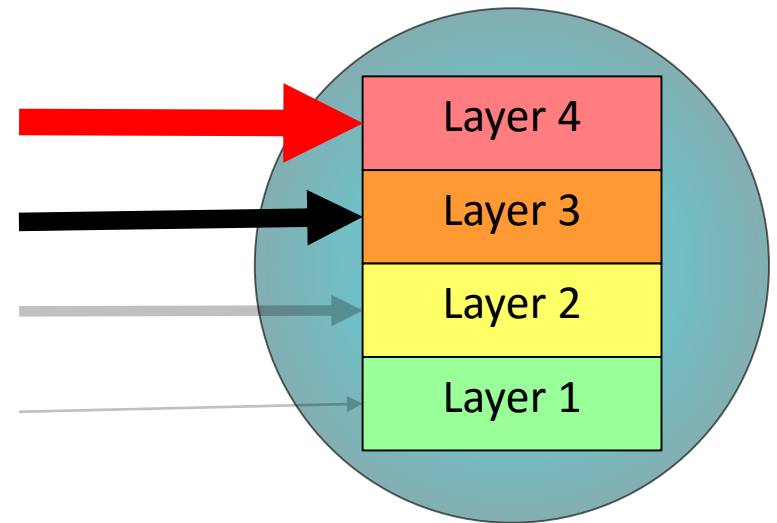
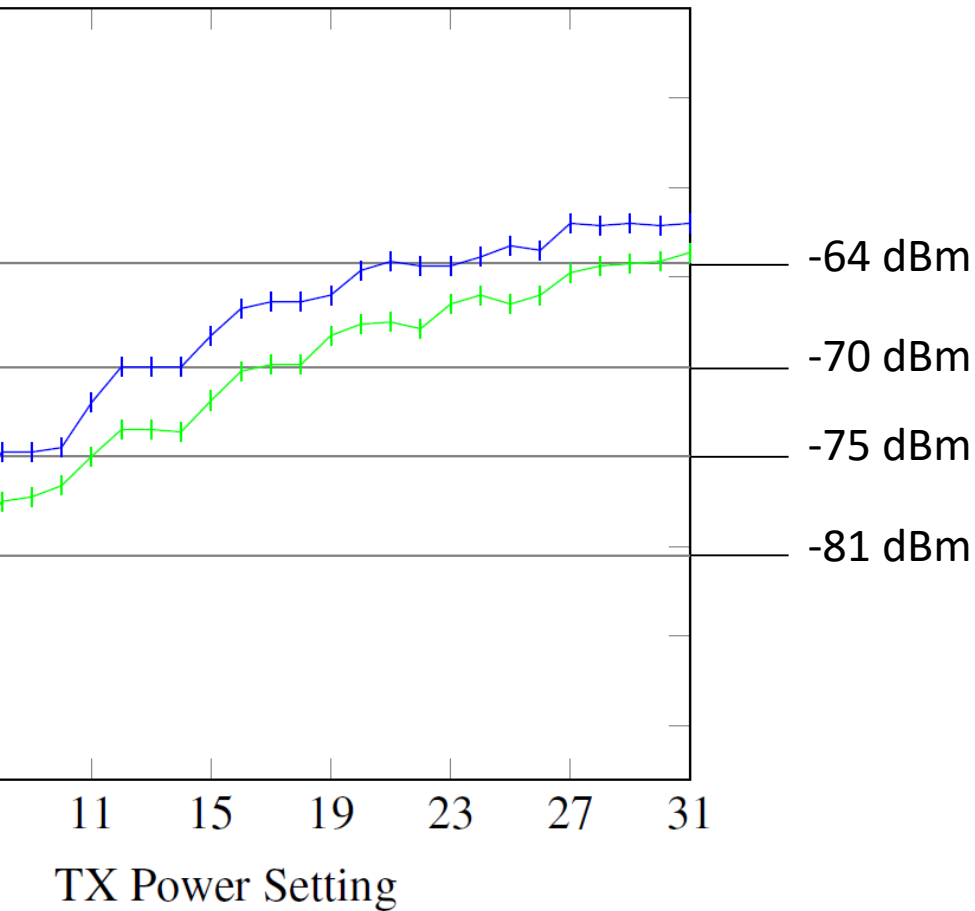
The Capture Effect



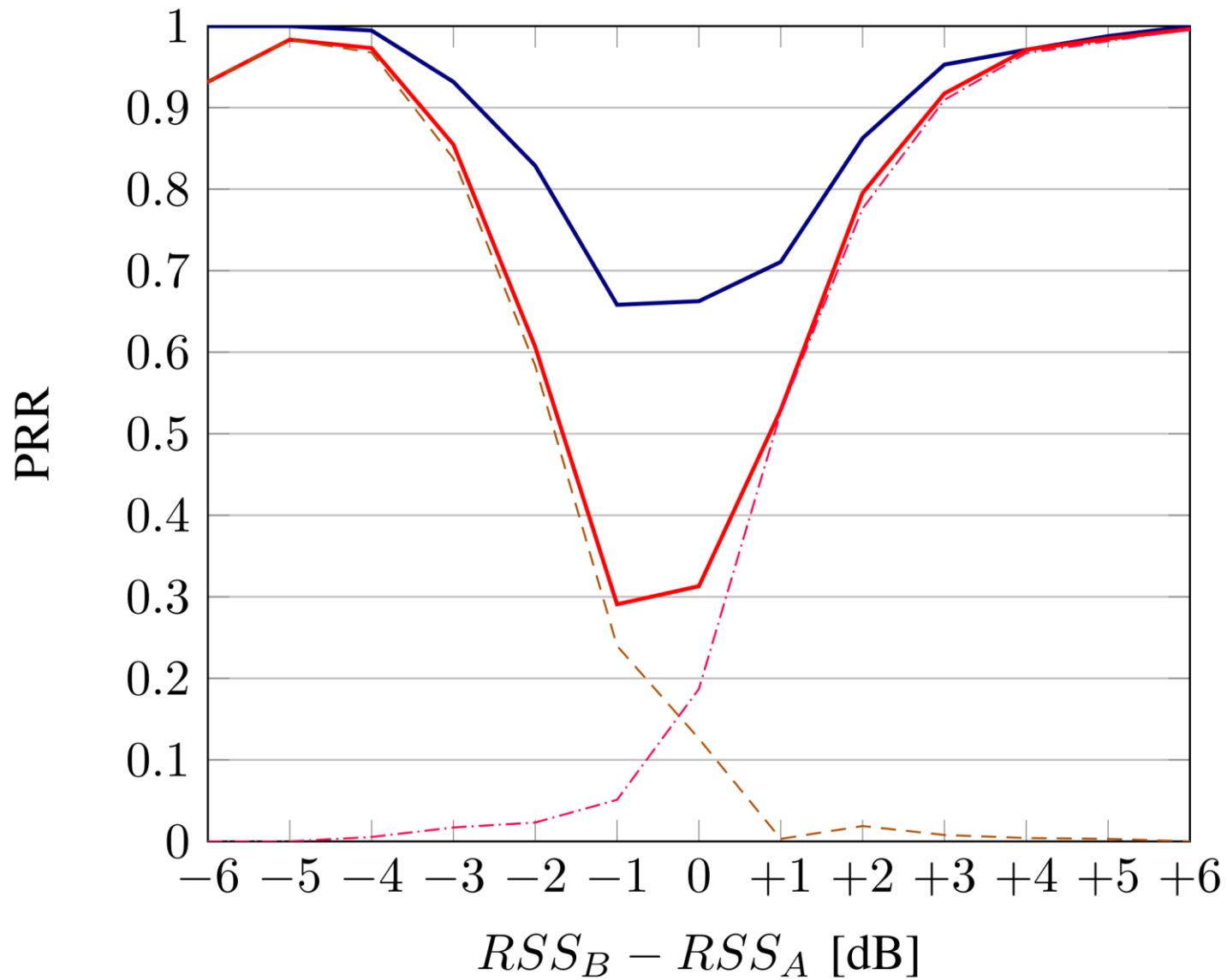
Receiving Different Senders



“Layer” Abstraction



Constructive Interference




Energy Efficiency?

Clock Synchronization!

Clock Synchronization Example: Dozer

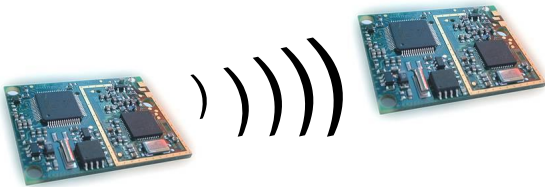
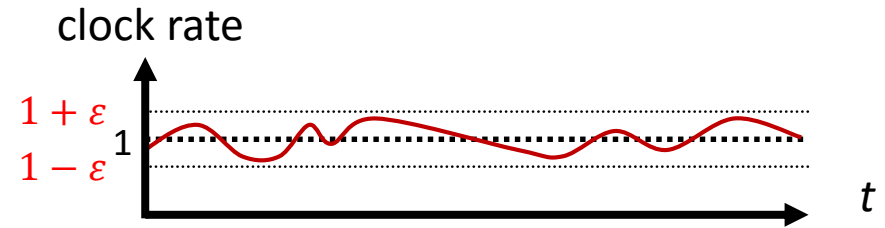
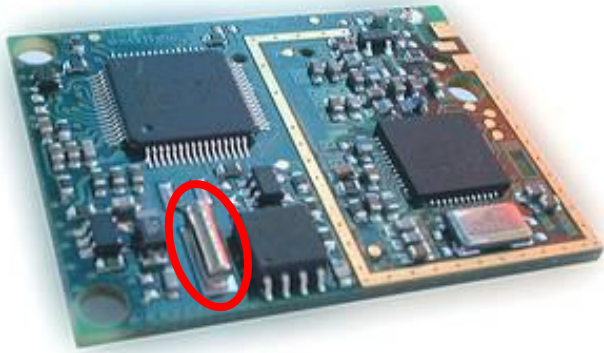
- Multi-hop sensor network with duty cycling
- 10 years of network life-time, mean energy consumption: 0.066mW
- High availability, reliability (99.999%)
- Many different applications use Dozer: TinyNode, PermaSense, etc.

[Burri, von Rickenbach, W, 2007]

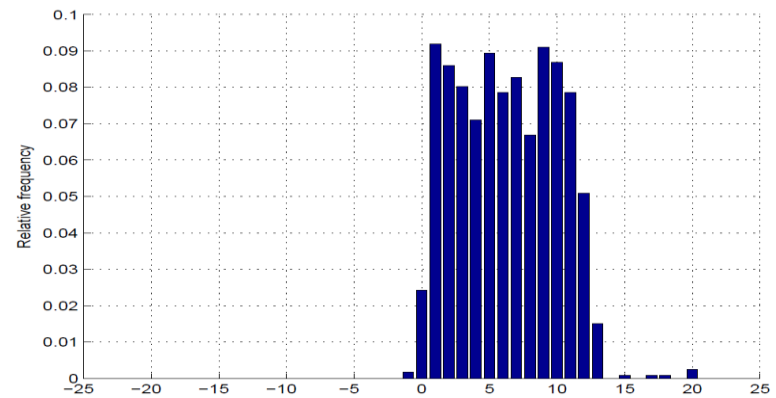


Wireless vehicle detection systems
for outdoor parking lots

Problem: Physical Reality



message delay

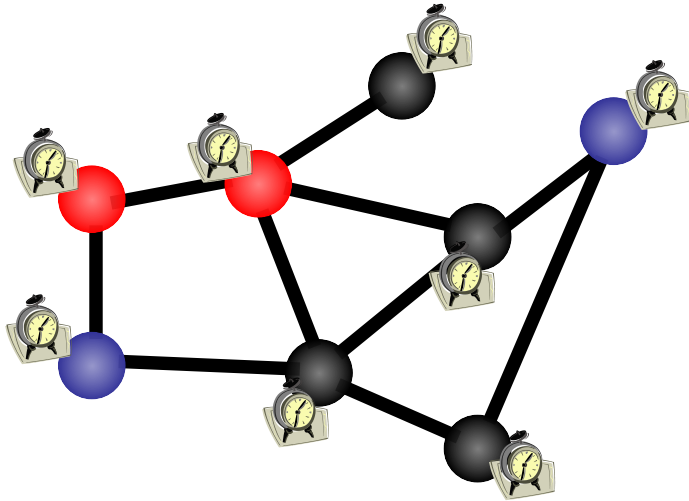


Clock Synchronization in Theory?

Given a communication network

1. Each node equipped with hardware clock with **drift**
2. Message delays with **jitter**

worst-case (but constant)



Goal: Synchronize Clocks (“Logical Clocks”)

- Both **global** and **local** synchronization!

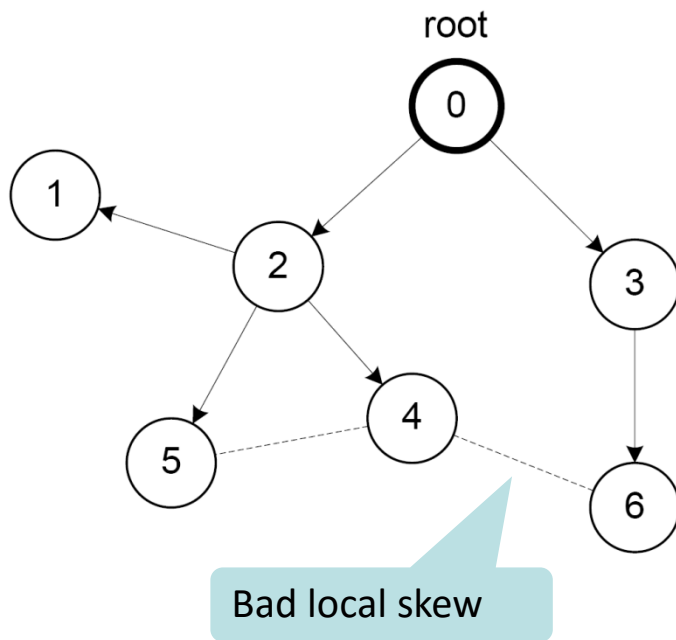
Time Must Behave!

- Time (logical clocks) should **not** be allowed to **stand still** or **jump**

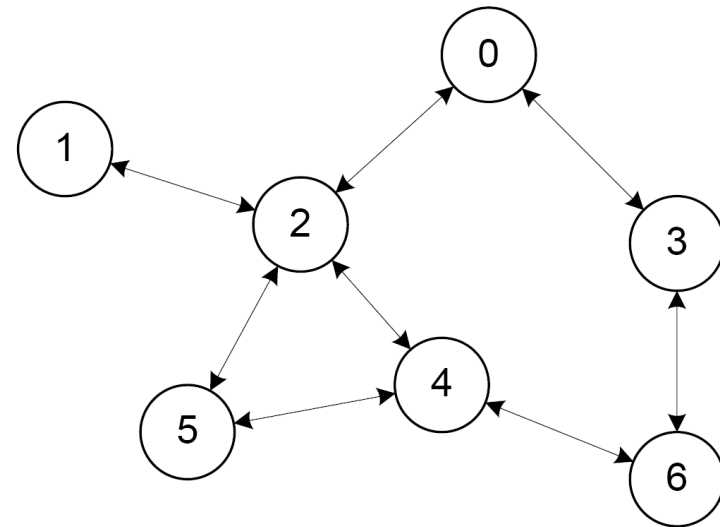


Local Skew

Tree-based Algorithms
e.g. FTSP

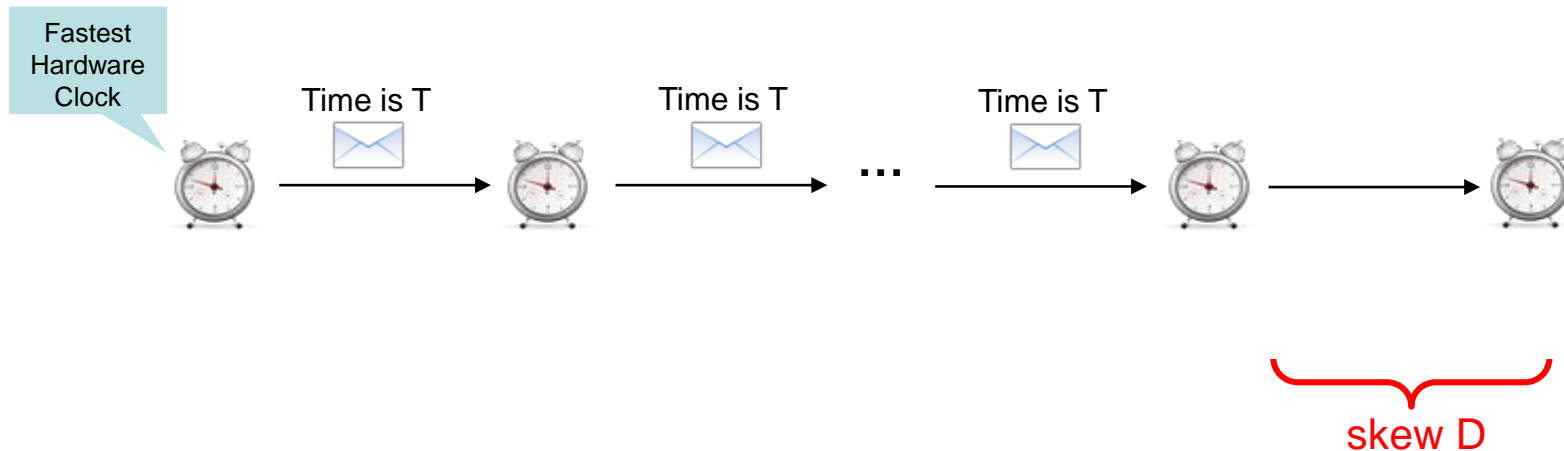


Neighborhood Algorithms
e.g. GTSP

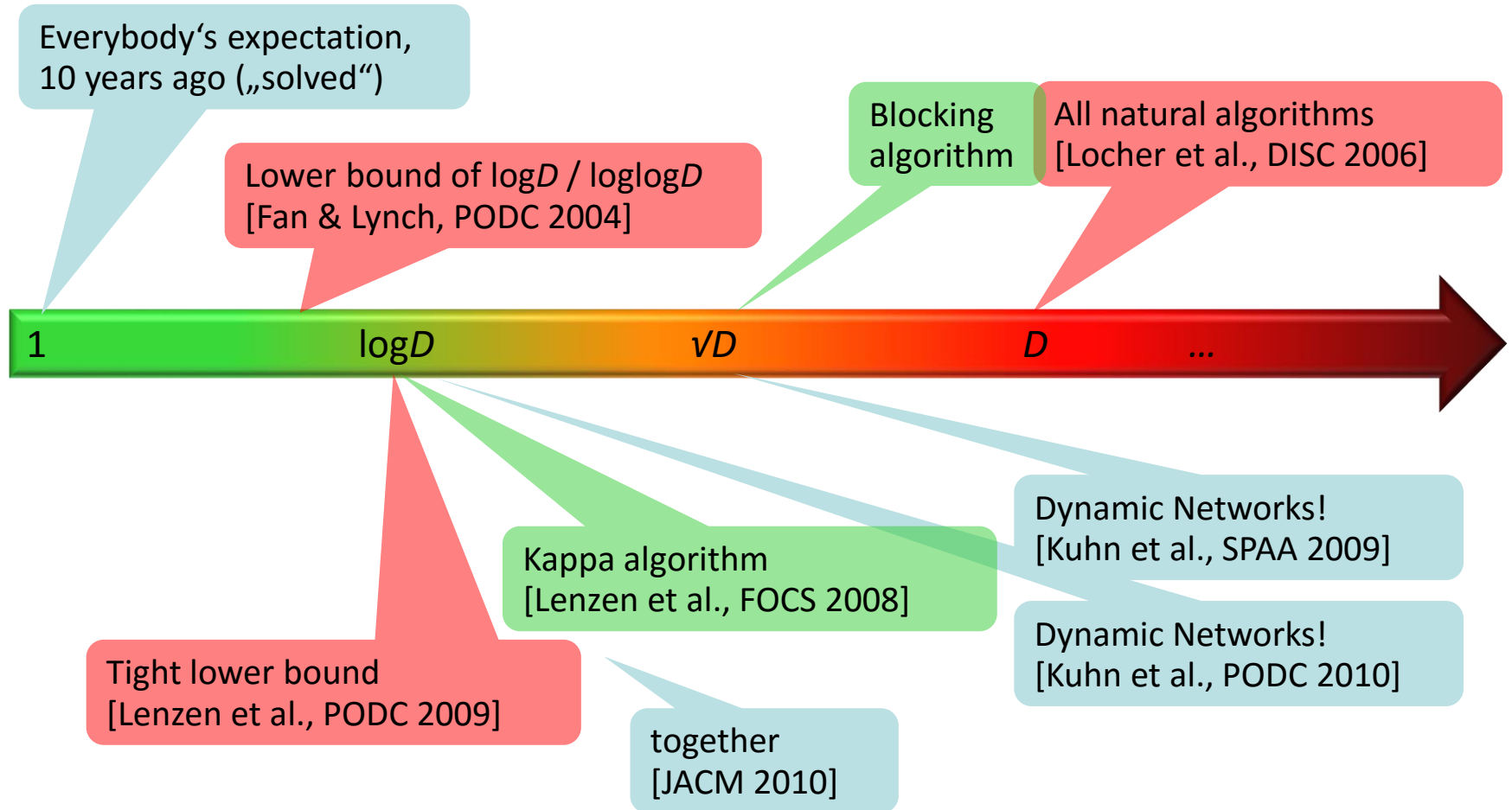


Synchronization Algorithms: An Example (“ A^{\max} ”)

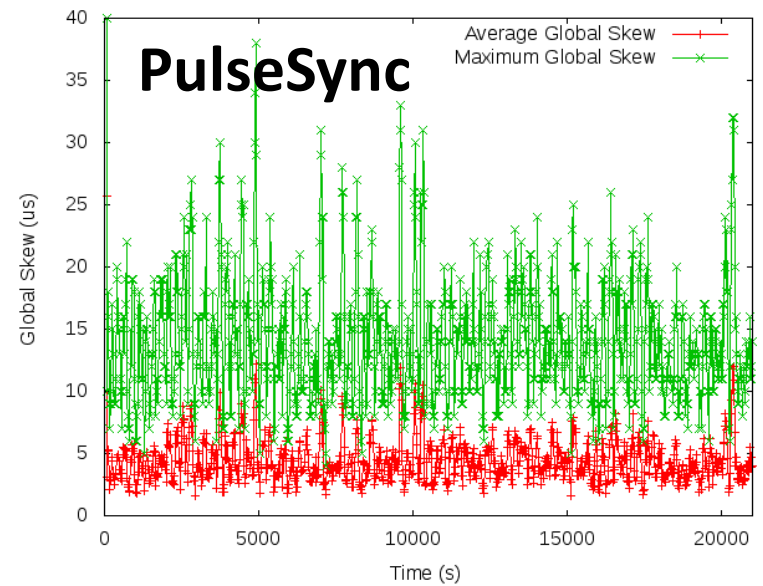
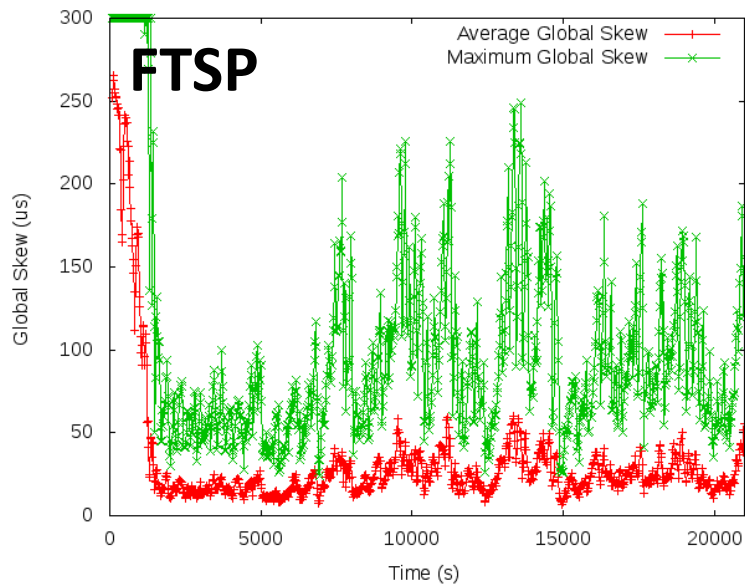
- Question: How to update the logical clock based on the messages from the neighbors?
- Idea: Minimizing the skew to the **fastest** neighbor
 - Set clock to **maximum** clock value you know, forward new values immediately
- First all messages are slow (1), then suddenly all messages are fast (0)!



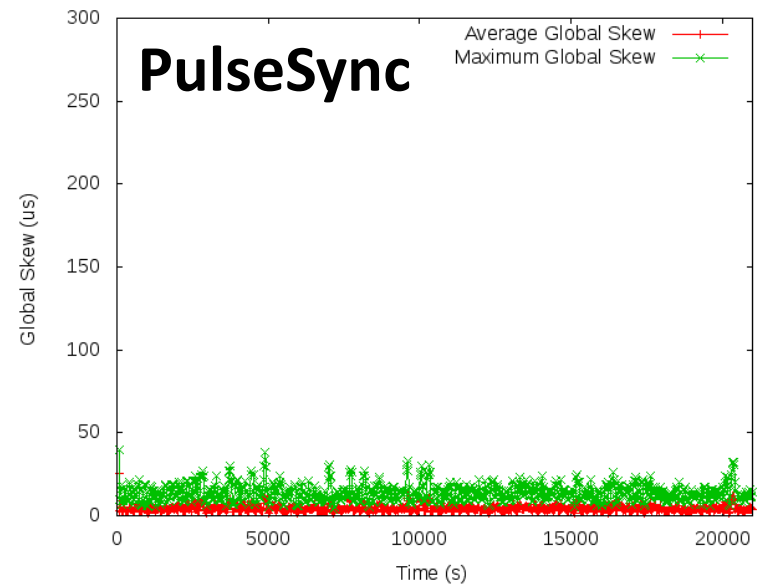
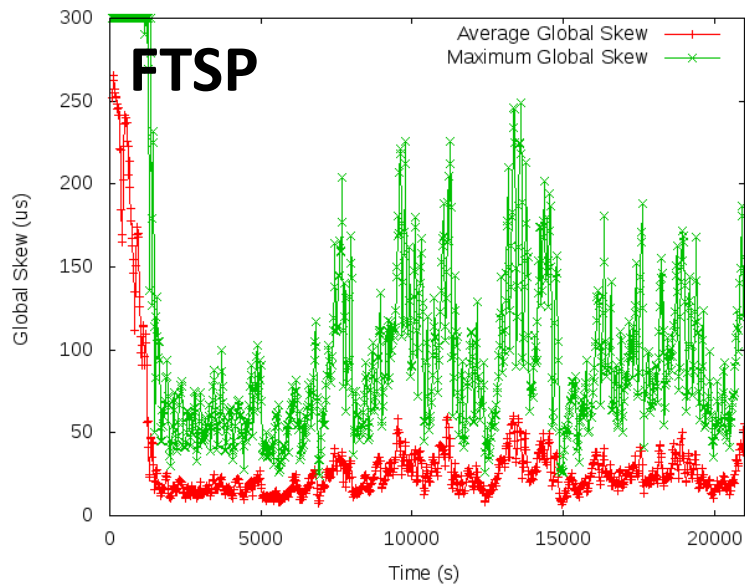
Local Skew: Overview of Results



Experimental Results for Global Skew

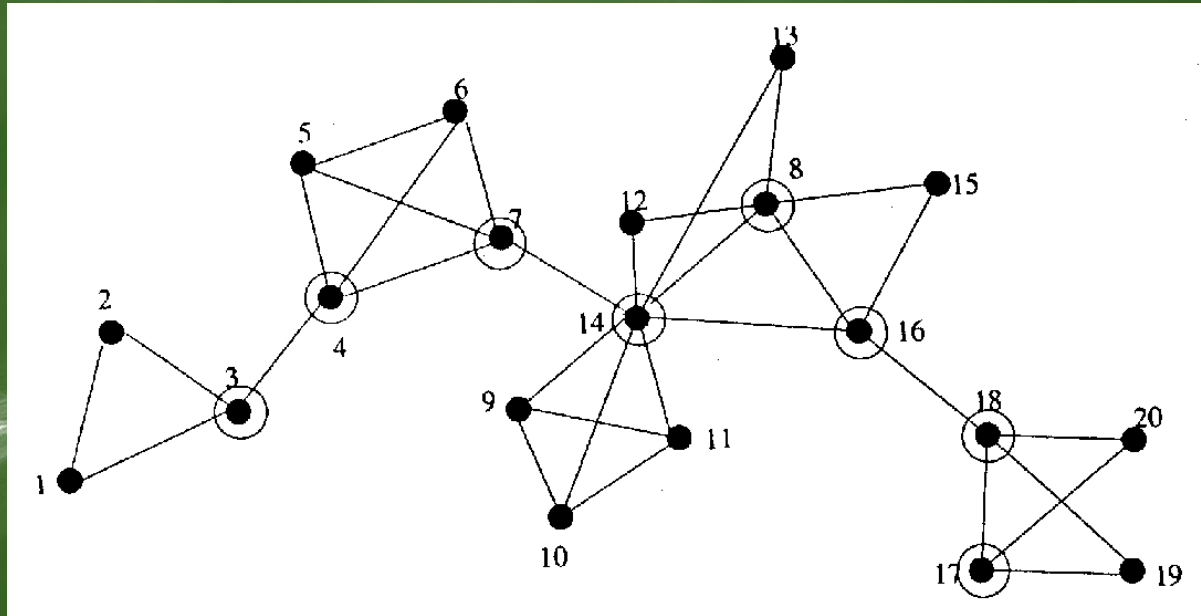


Experimental Results for Global Skew



Network Dynamics?

Distributed Control!



Complexity Theory

Can a Computer Solve
Problem P in Time t ?

Distributed



Complexity Theory

Network

Can a ~~Computer~~ Solve
Problem P in Time t ?

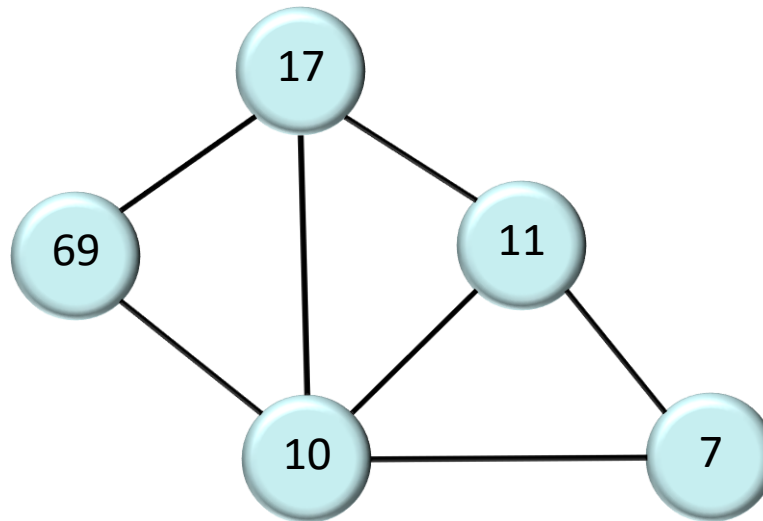
Network
~~Distributed~~

↓ Complexity Theory

Network
Can a ~~Computer~~ Solve
Problem P in Time t ?

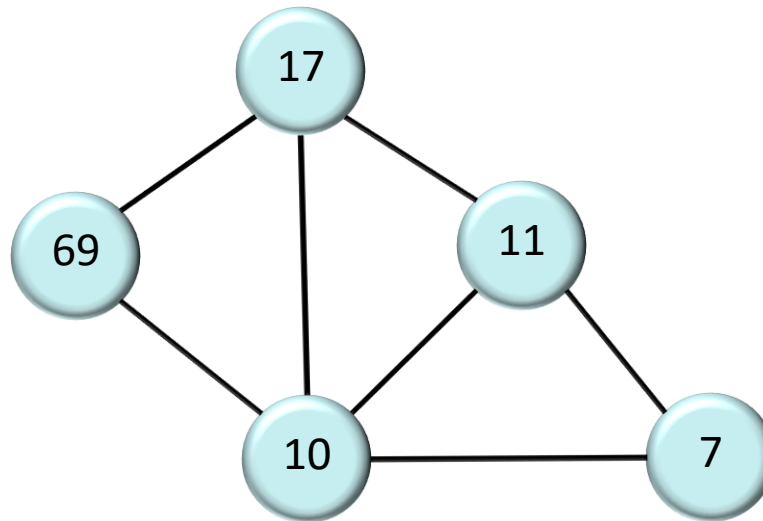
Distributed (Message-Passing) Algorithms

- Nodes are agents with unique ID's that can communicate with neighbors by **sending messages**. In each **synchronous round**, every node can send a (different) message to each neighbor.



Distributed (Message-Passing) Algorithms

- Nodes are agents with unique ID's that can communicate with neighbors by **sending messages**. In each **synchronous round**, every node can send a (different) message to each neighbor.



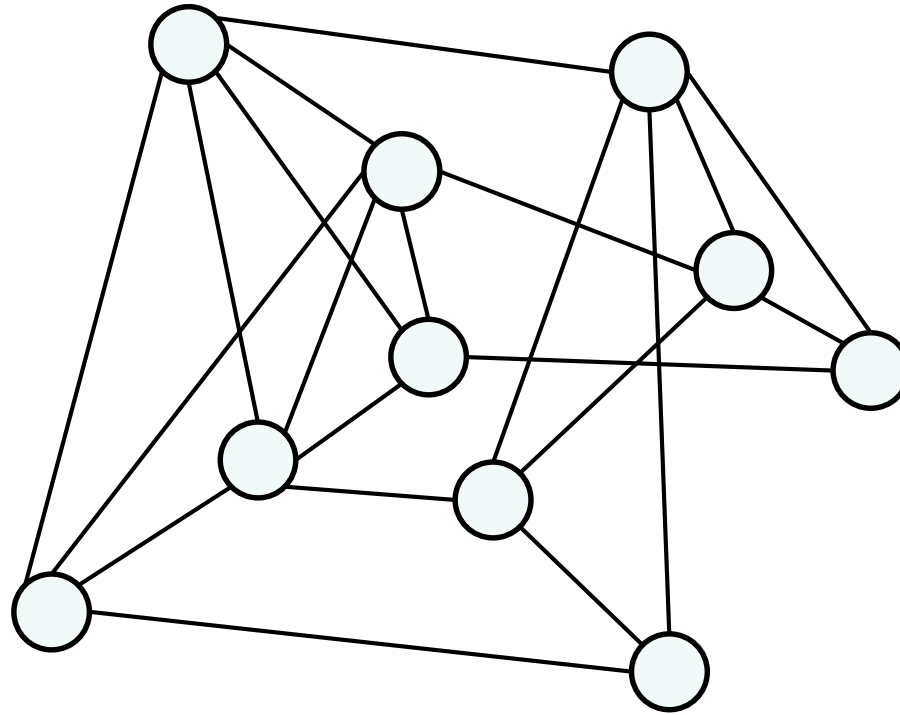
each round:
every node:
1. send msgs
2. rcv msgs
3. compute

- Distributed (Time) Complexity**: How many rounds does problem take?

An Example

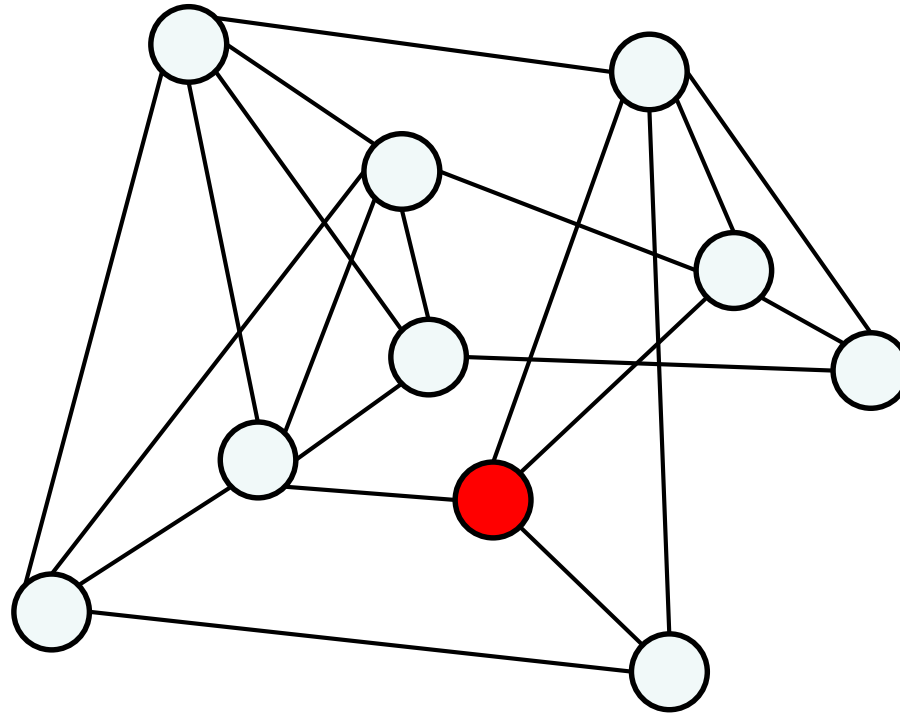
each round:
every node:
1. send msgs
2. rcv msgs
3. compute

How Many Nodes in Network?



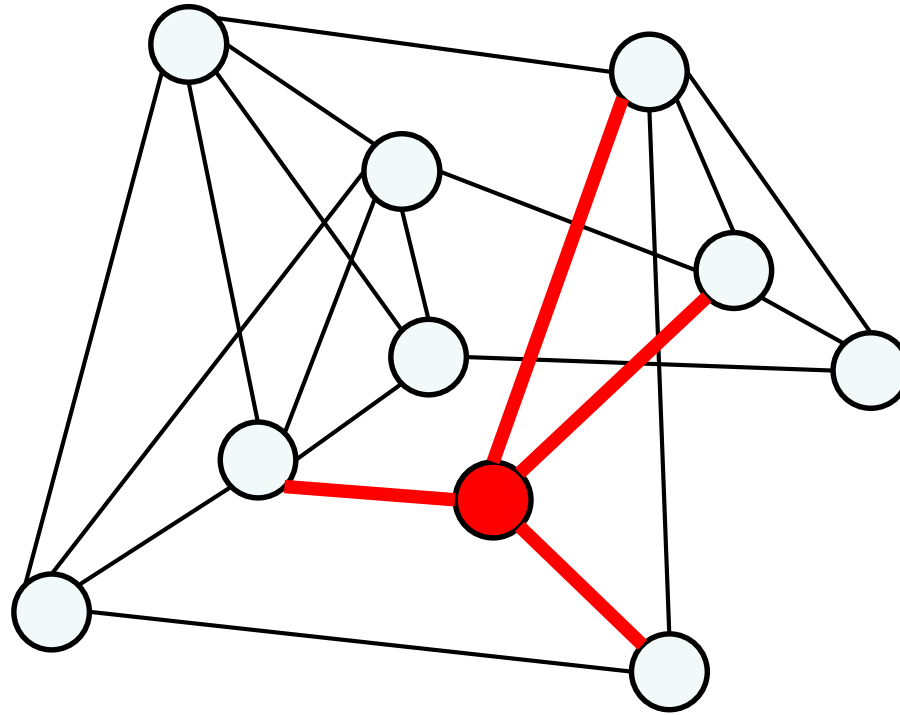
each round:
every node:
1. send msgs
2. rcv msgs
3. compute

How Many Nodes in Network?



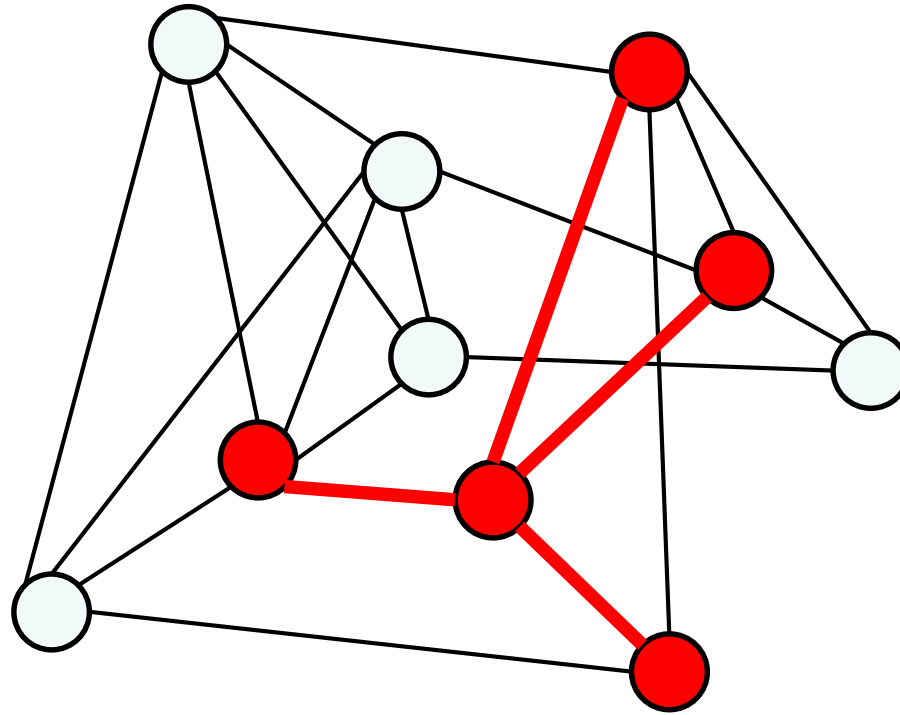
each round:
every node:
1. send msgs
2. rcv msgs
3. compute

How Many Nodes in Network?

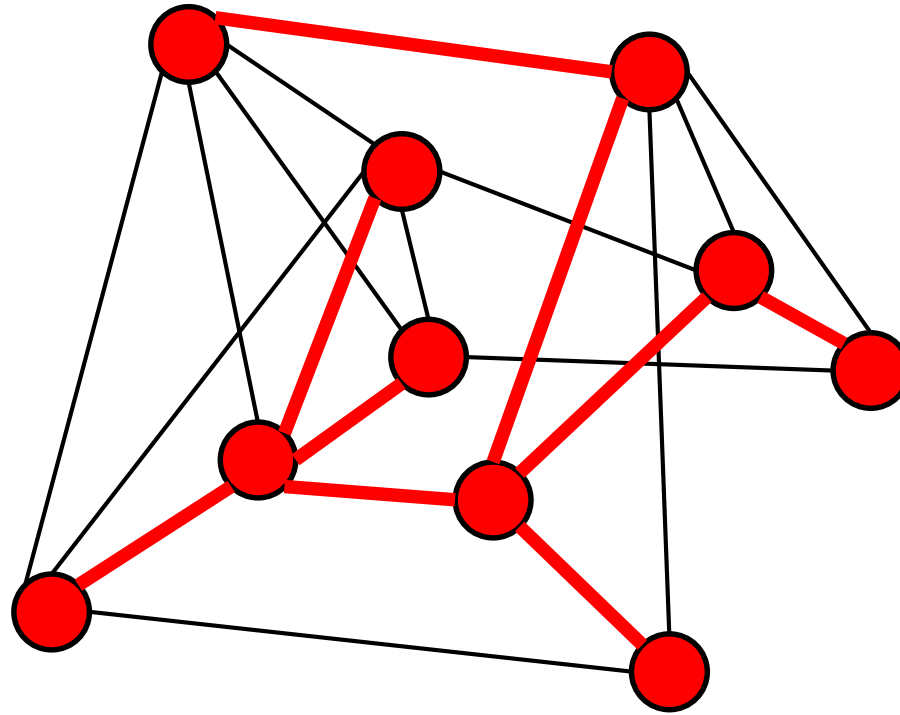


each round:
every node:
1. send msgs
2. rcv msgs
3. compute

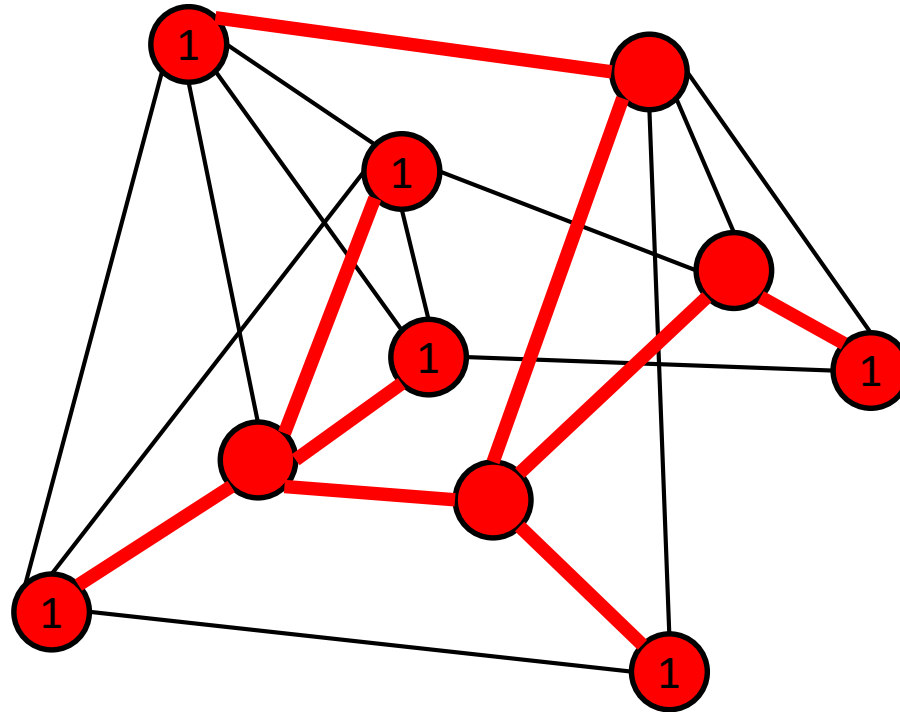
How Many Nodes in Network?



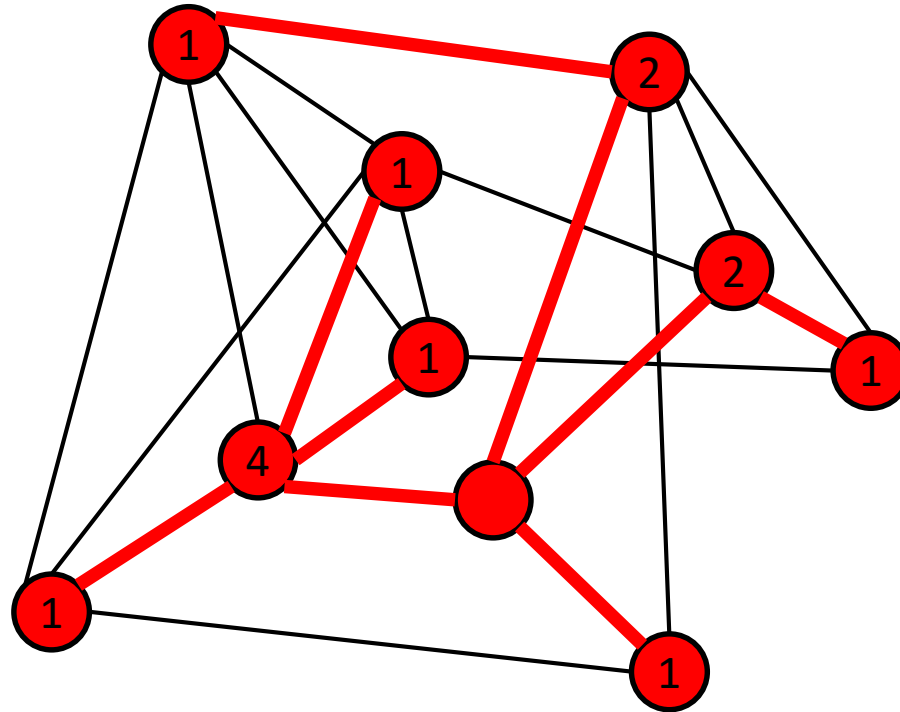
How Many Nodes in Network?



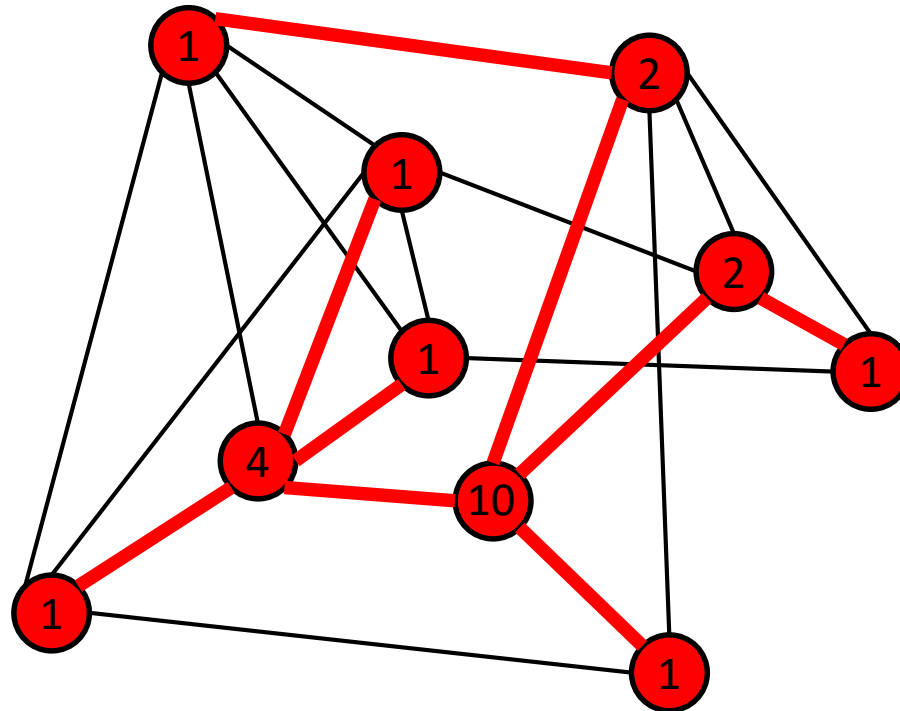
How Many Nodes in Network?



How Many Nodes in Network?

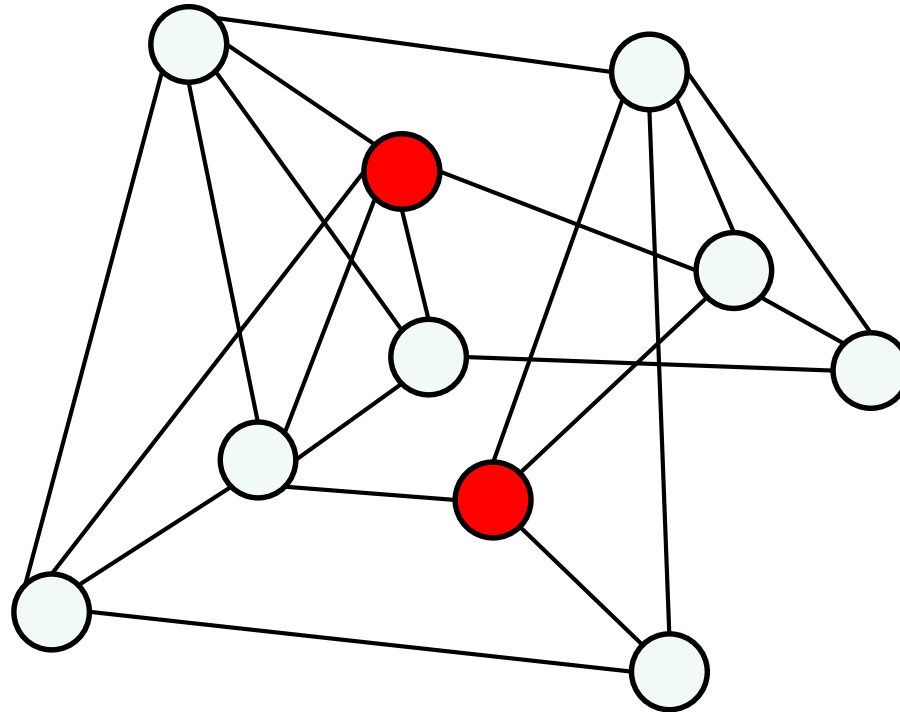


How Many Nodes in Network?



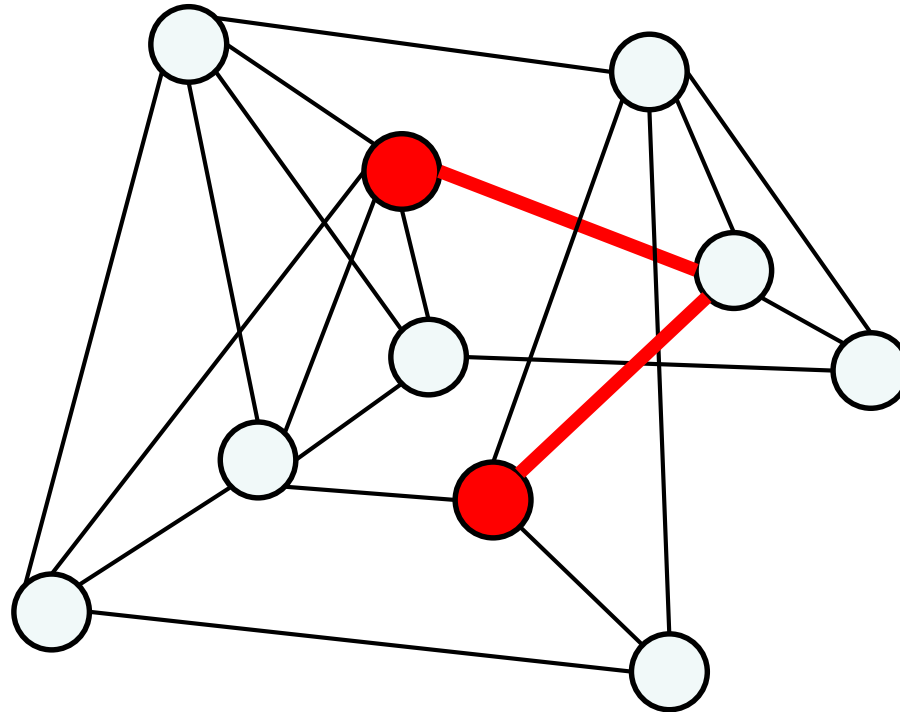
With a simple flooding/echo process, a network can find the number of nodes in **time $O(D)$** , where D is the diameter (size) of the network.

Diameter of Network?



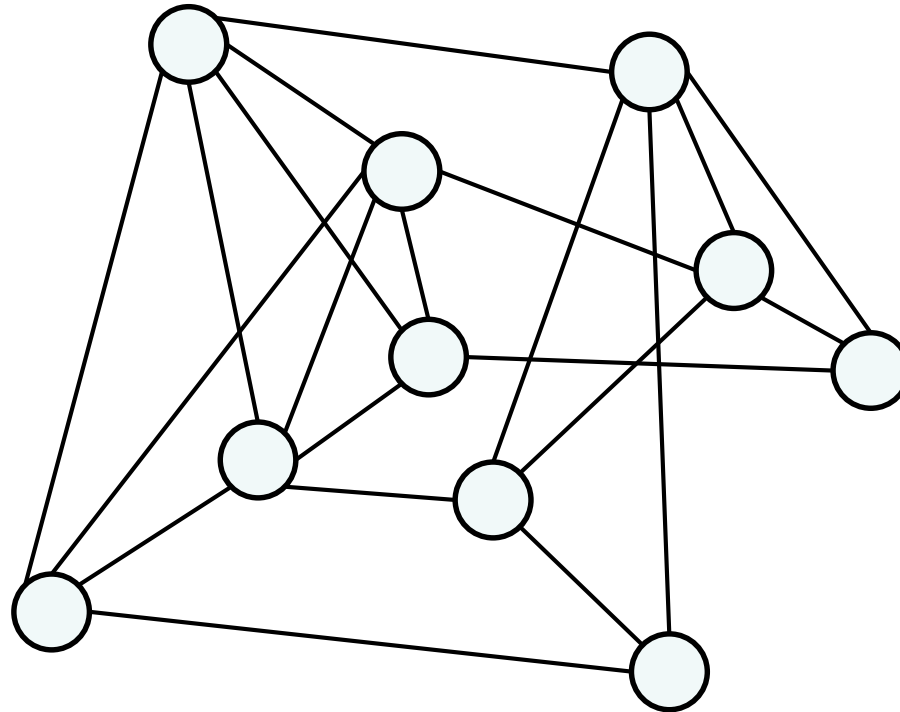
- **Distance** between two nodes = Number of hops of shortest path

Diameter of Network?



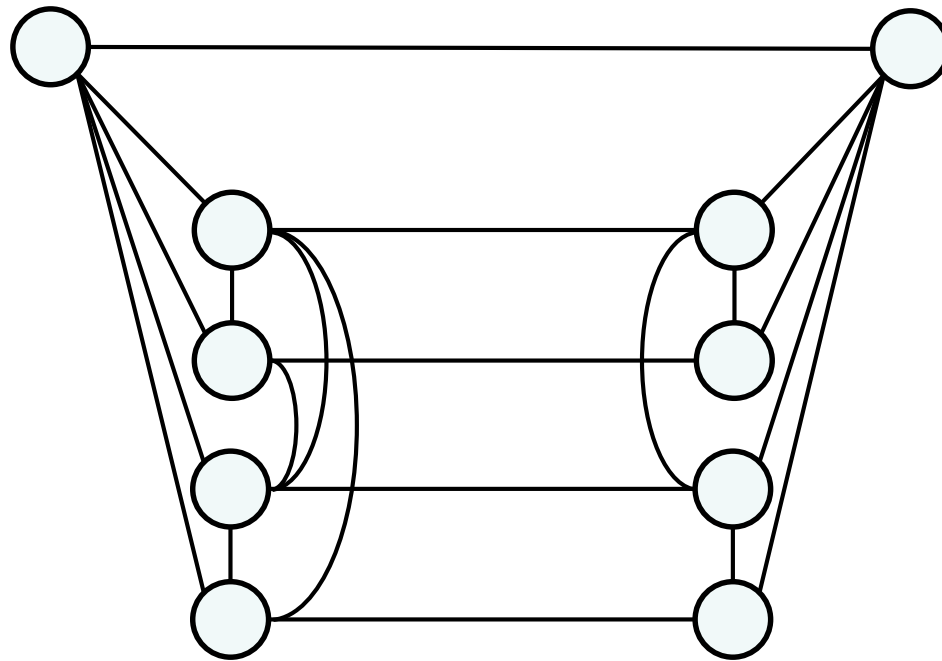
- **Distance** between two nodes = Number of hops of shortest path

Diameter of Network?

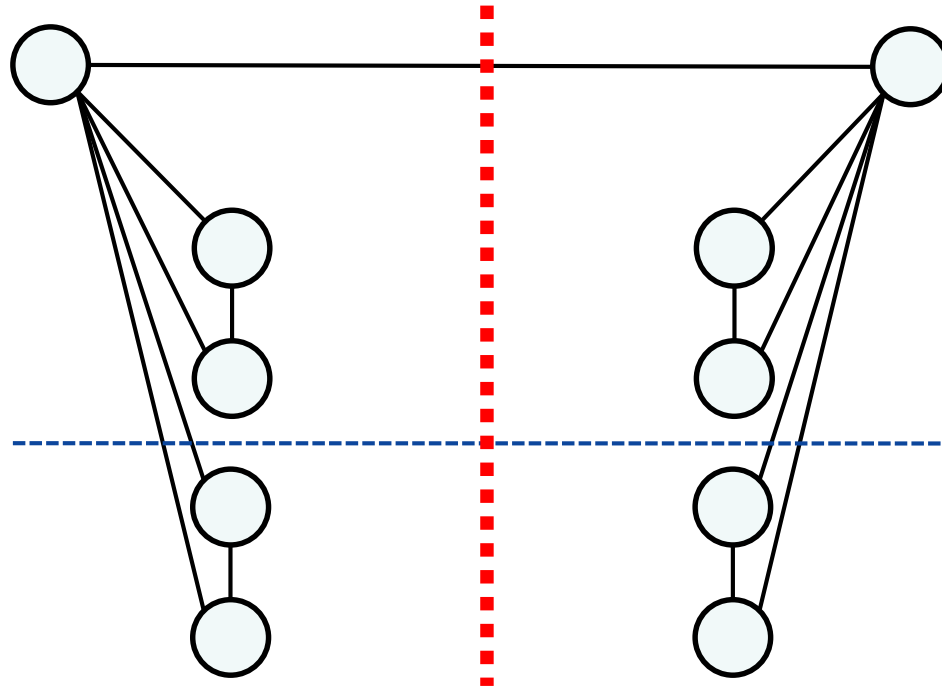


- **Distance** between two nodes = Number of hops of shortest path
- **Diameter** of network = Maximum distance, between any two nodes

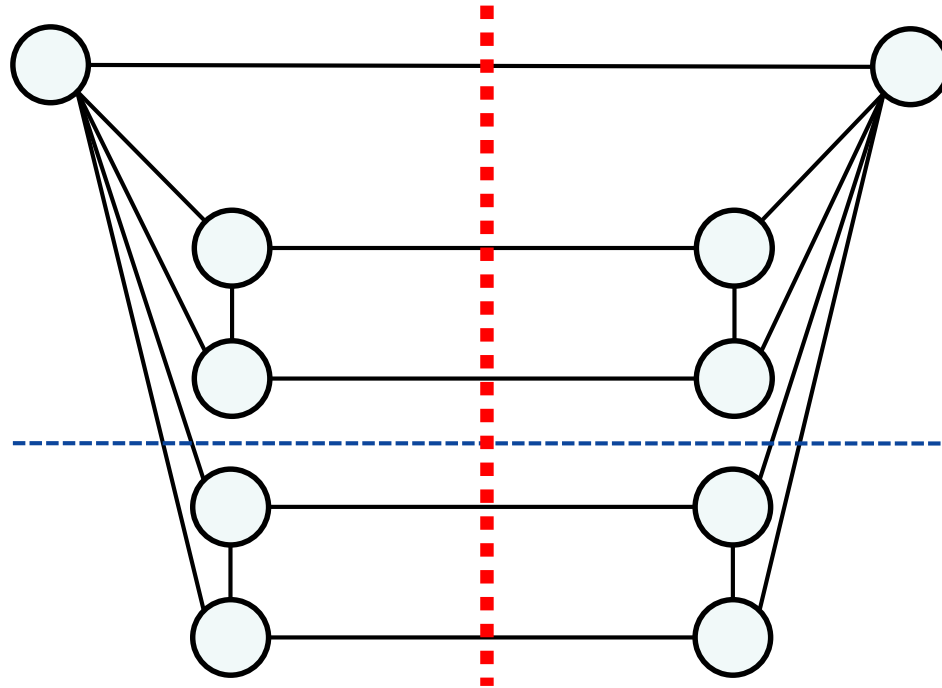
Diameter of Network?



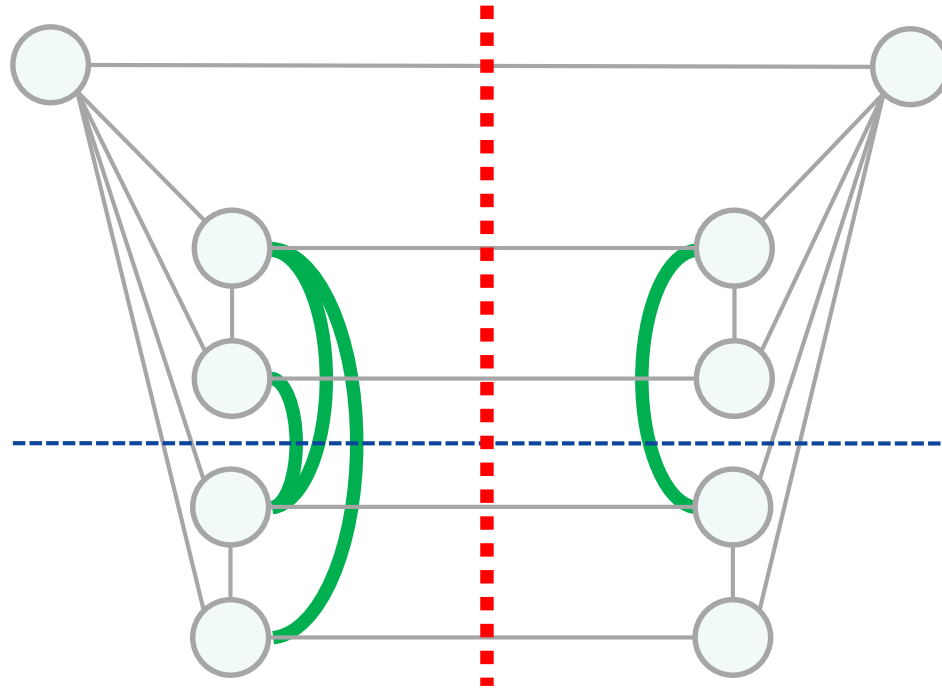
Diameter of Network?



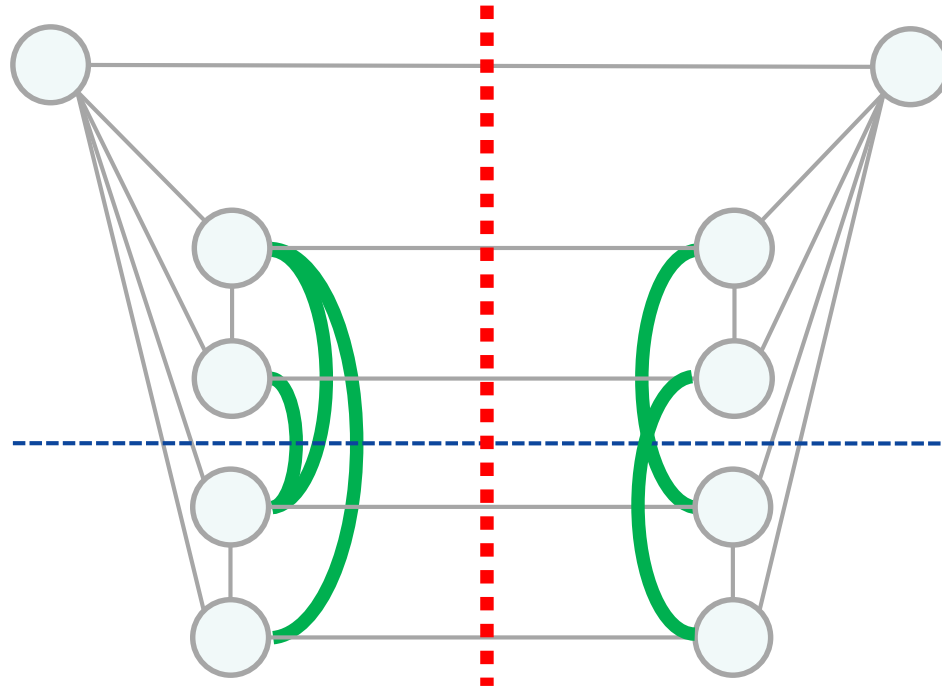
Diameter of Network?



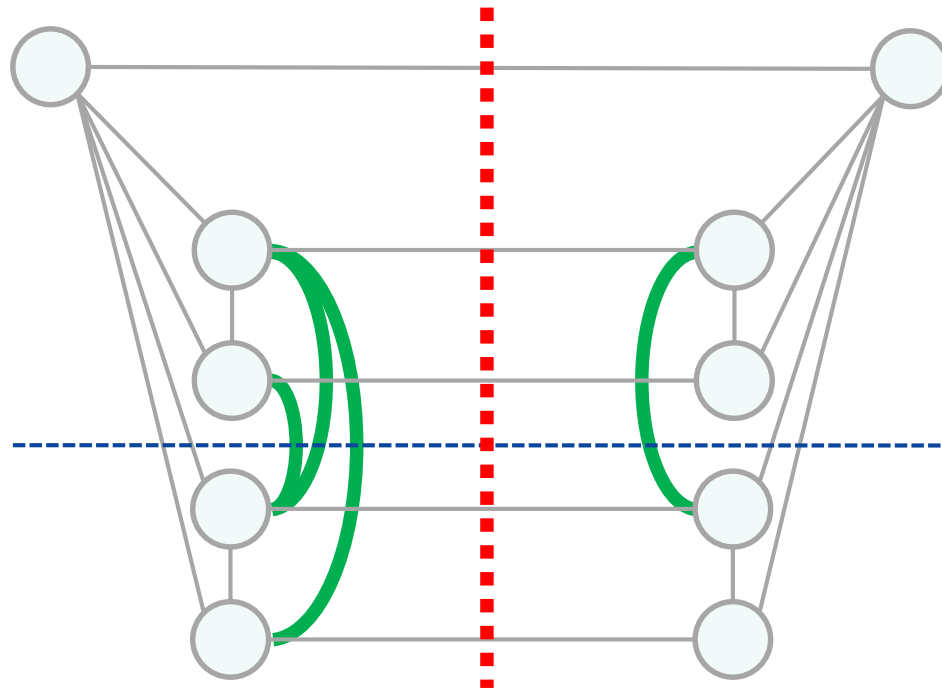
Diameter of Network?



Diameter of Network?

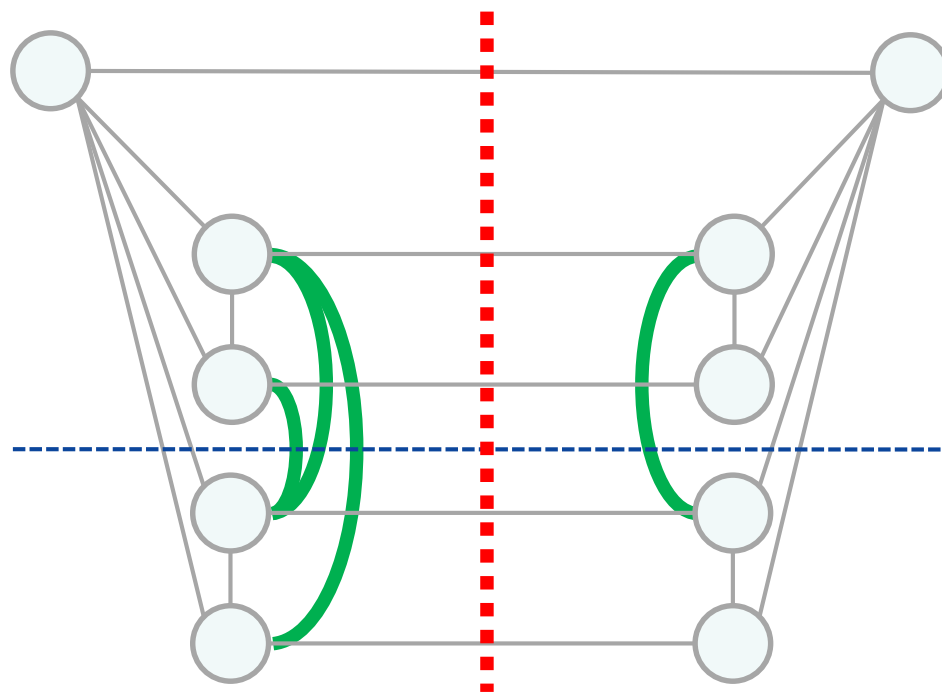


Diameter of Network?



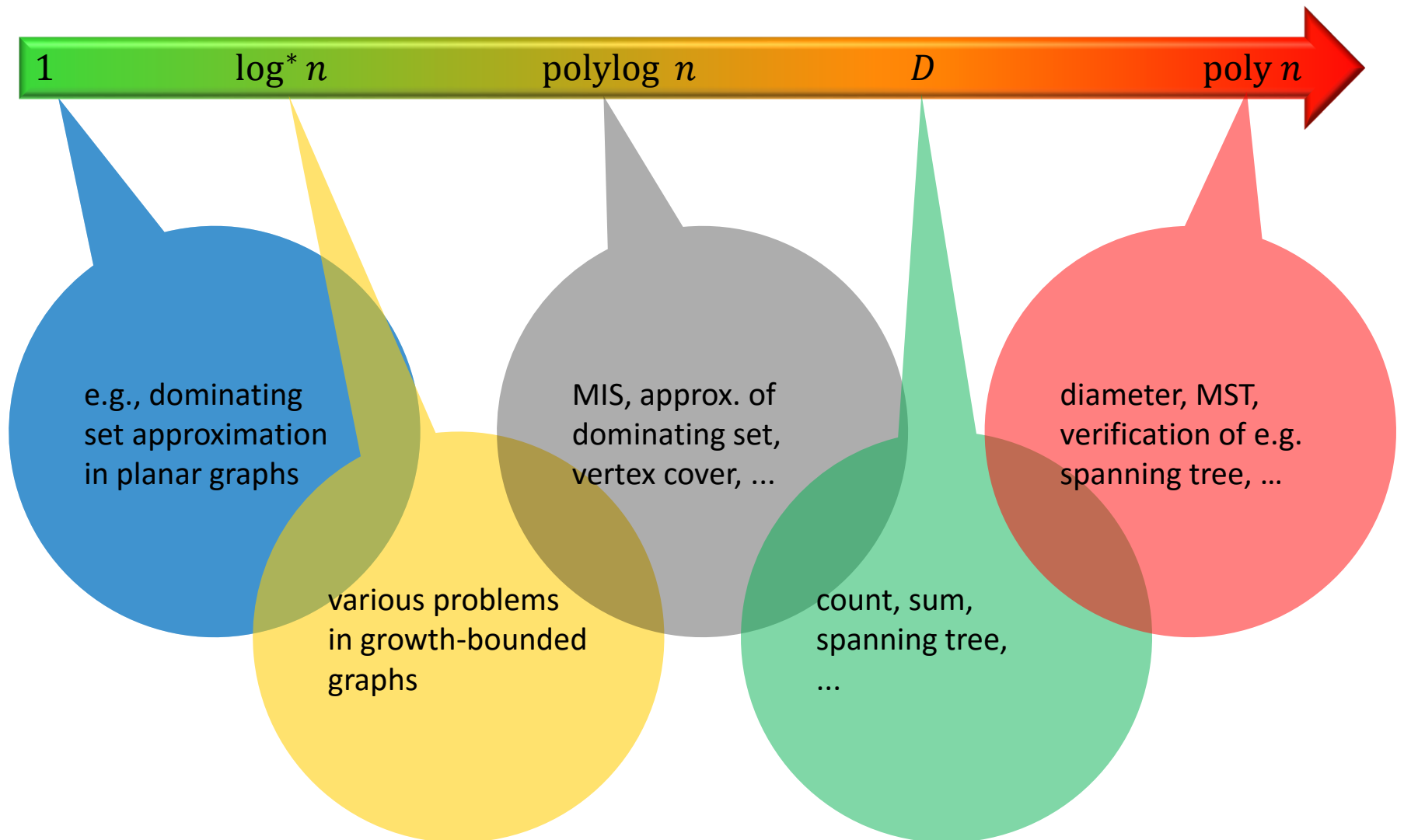
Networks Cannot Compute Their Diameter in Sublinear Time!

(even if diameter is just a small constant)

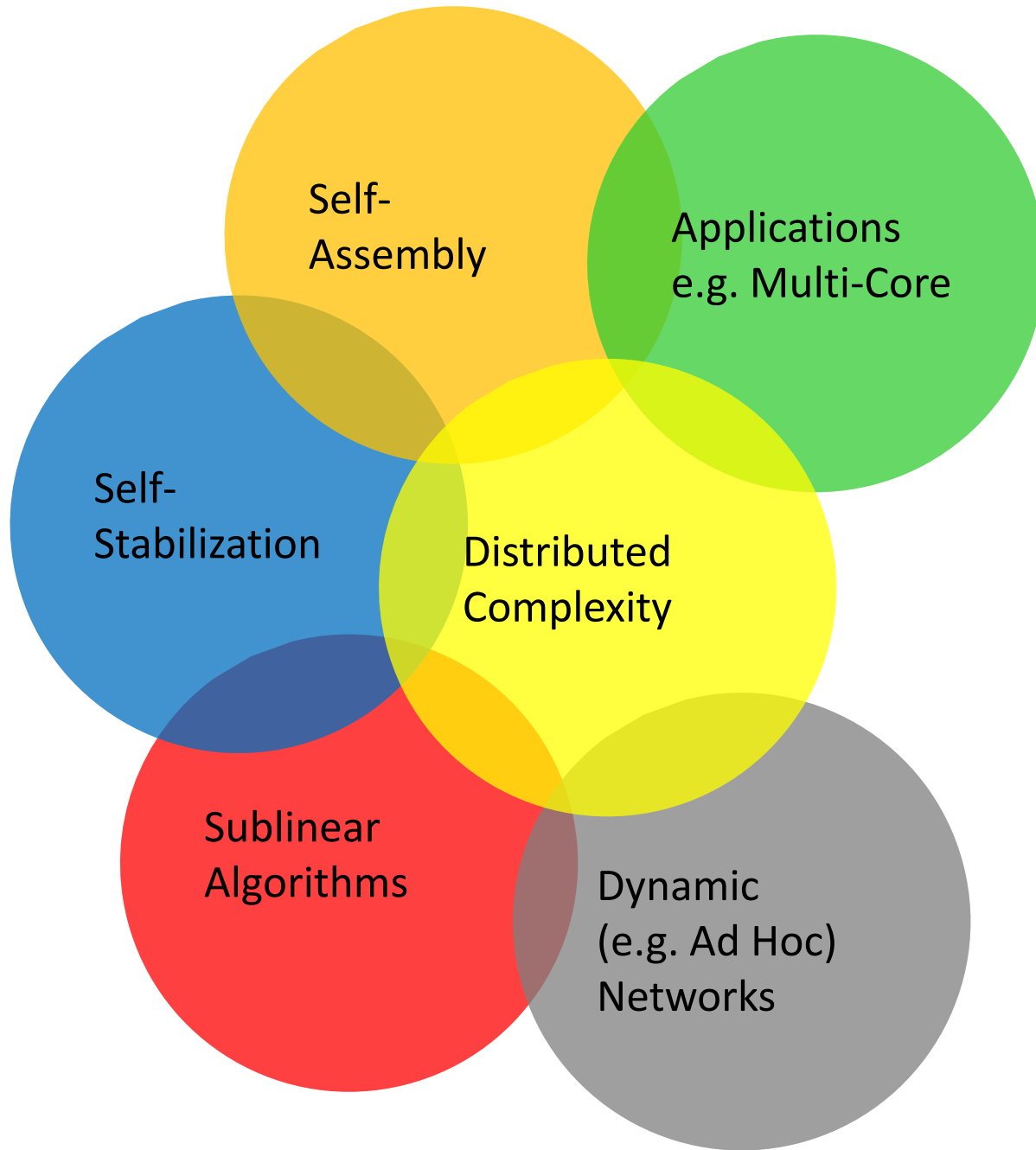


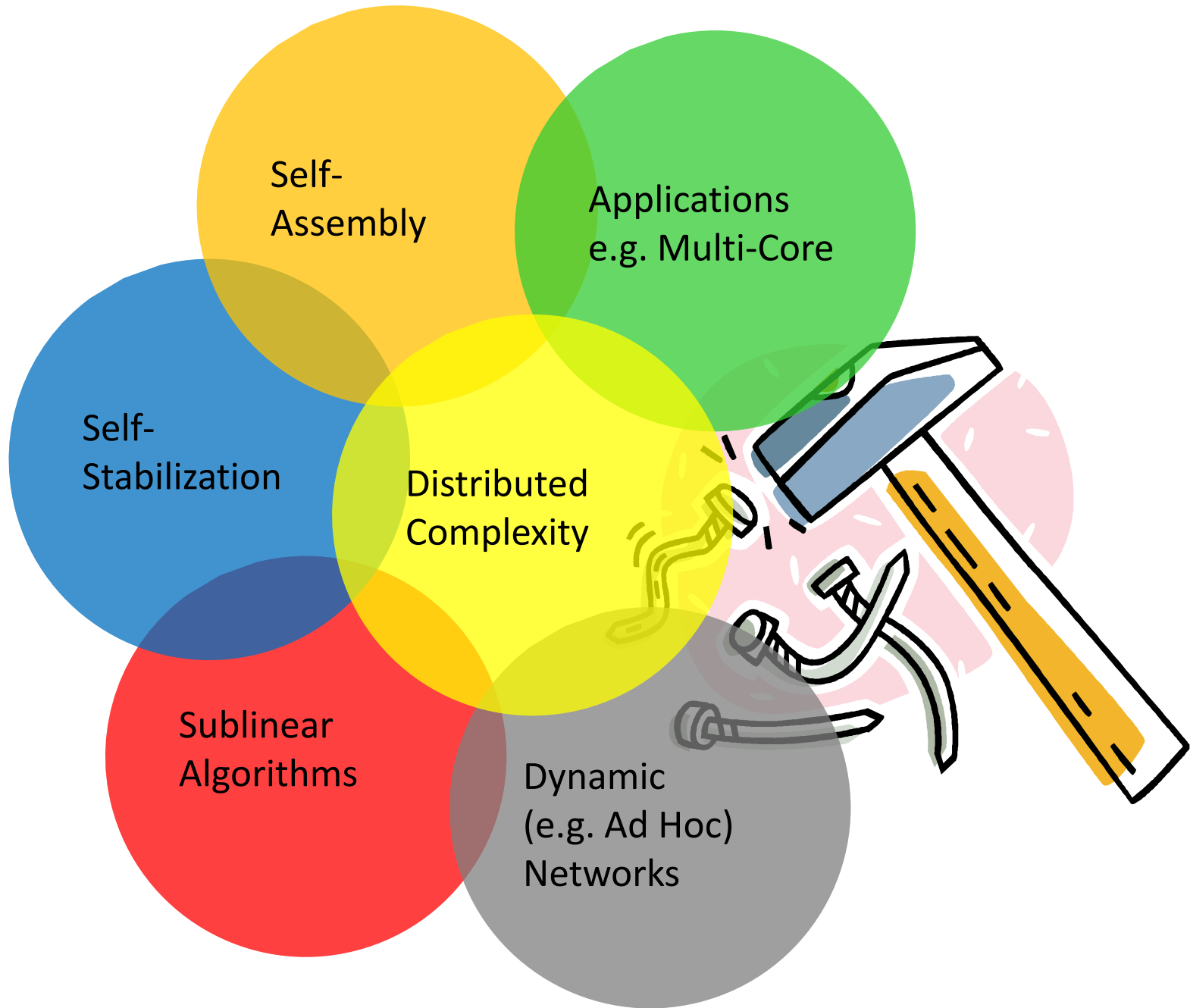
Pair of rows connected neither left nor right? Communication complexity:
Transmit $\Theta(n^2)$ information over $O(n)$ edges $\rightarrow \Omega(n)$ time!

Distributed Complexity Classification



e.g., [Kuhn, Moscibroda, W, 2016]





Self-Assembly

Applications
e.g. Multi-Core

Self-Stabilization

Distributed Complexity

Sublinear Algorithms

Dynamic
(e.g. Ad Hoc)
Networks

$O(1)$ -APX,
 $O(1)$ -time

$w(n)$ -APX
 $\log^* \text{-time}$

Series-parallel
→ planar
trees

planar
proj.
plane

planar
2-fold
cover

(bounded tree-w.)

triangle-free

some forbidden ind. subgr.

no $K_{3,3}$

no $K_{3,5}$

some
forbidden
minor

no K_5

sparse

sparse,
 d_1, d_2, d_3

bounded arb.

bound
indep.
dom. p.
claw-free
line graph
 $f(n)$ -reg.

d -regular

sparse,
 d_1, d_2

bounded
degree

growth-
bounded

$O(1)$ -APX
 $\log^* \text{-time}$

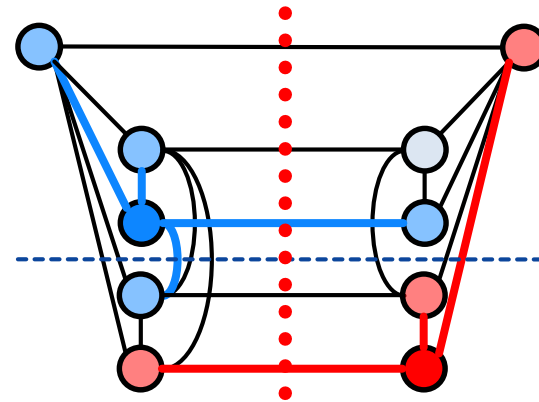
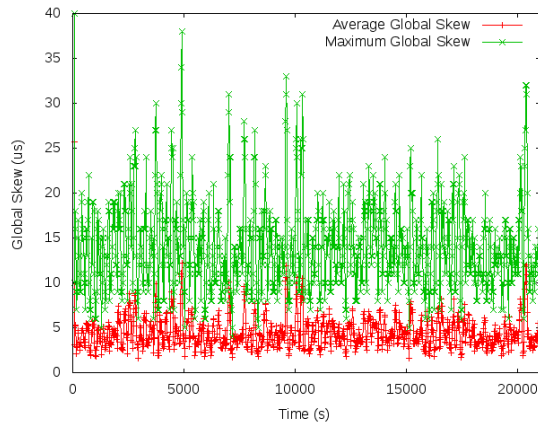
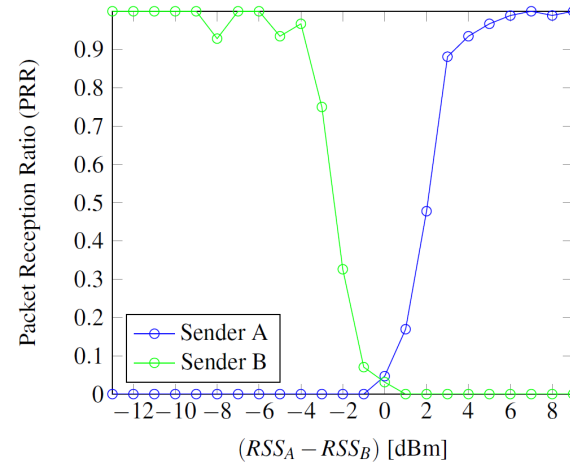
bounded
diam.

gb +
sparse

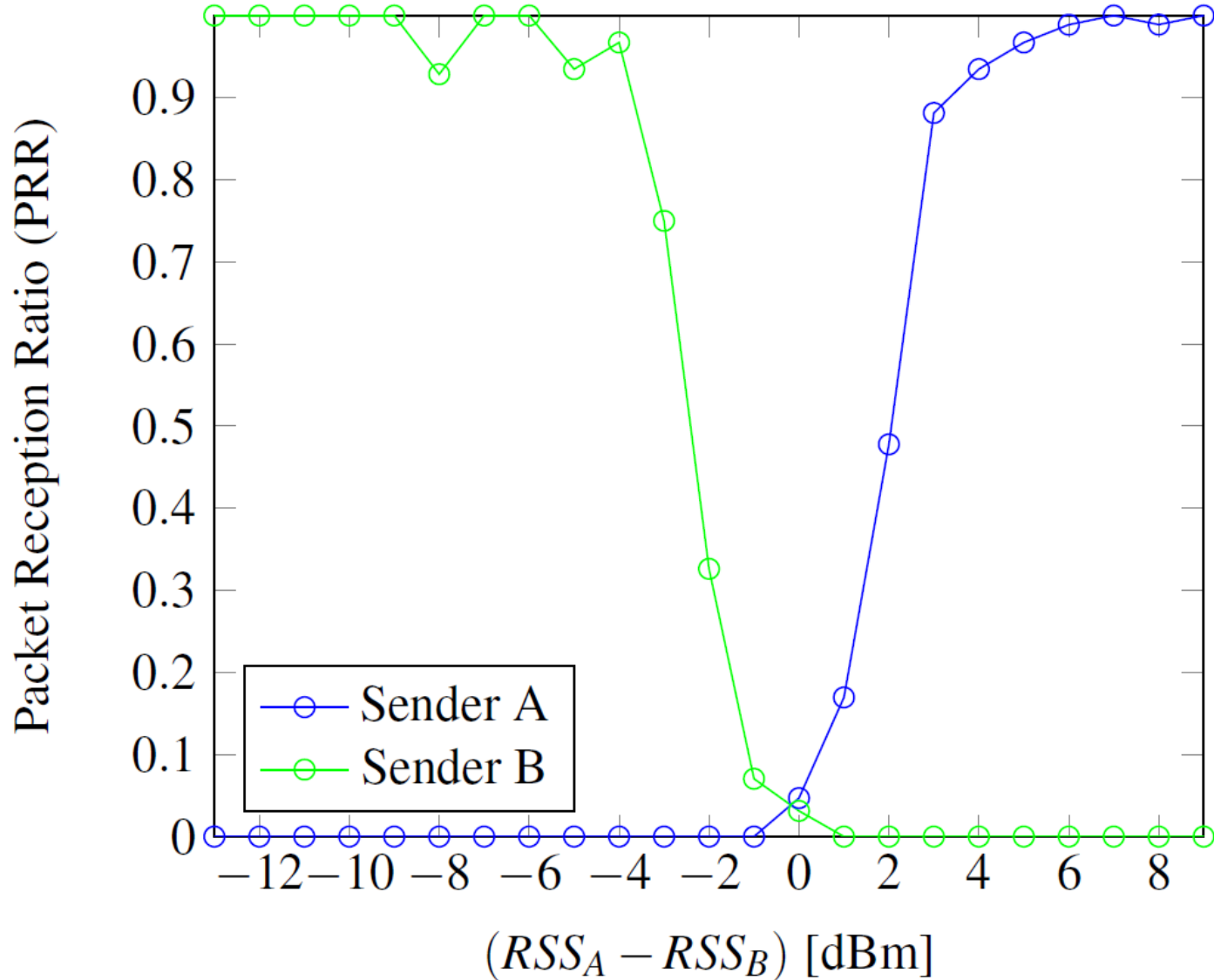
cliques



Summary



The Capture Effect



Theory for sensor networks, what is it good for?!

How many lines of pseudo code
Can you implement on a sensor node?

The best algorithm is often complex
And will not do what one expects.

Theory models made lots of progress
Reality, however, they still don't address.

My advice: invest your research \$\$\$
in ... impossibility results and lower bounds!



Thank You!

Questions & Comments?



Thanks to my co-authors, mostly
Silvio Frischknecht
Magnus Halldorsson
Stephan Holzer
Michael König
Christoph Lenzen
Thomas Moscibroda
Philipp Sommer

www.disco.ethz.ch