

Communities in Preference Networks: Refined Axioms and Beyond

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Abstract—Borgs *et al.* [2016] investigated essential requirements for communities in preference networks. They defined six axioms on community functions, i.e., community detection rules. Though having elegant properties, the practicality of this axiom system is compromised by the intractability of checking two critical axioms, so no nontrivial consistent community function was reported in [Borgs *et al.*, 2016]. By adapting the two axioms in a natural way, we propose two new axioms that are efficiently-checkable. We show that most of the desirable properties of the original axiom system are preserved. More importantly, the new axioms provide a general approach to constructing consistent community functions. We further find a natural consistent community function that is also enumerable and samplable, answering an open problem in the literature.

I. INTRODUCTION

Clustering individuals in a social network, called community detection, is a fundamental task in graph mining and has been adequately studied. Community detection has different forms, depending on whether overlapping communities are allowed [Palla *et al.*, 2005; Baumes *et al.*, 2005; Zhang *et al.*, 2007; Ahn *et al.*, 2009], whether hierarchical structures are taken into account [Sibson, 1973], and in which form of the data is provided, etc. People proposed a number of algorithms to find communities, based on different principles such as spectral clustering [Hoffman, 1973], density-based methods [Ester *et al.*, 1996], modularity-based method [Newman, 2006]. No matter which form and which algorithm we choose, a community is usually considered as a group of closely related individuals. However, “a group of closely related individuals” is a rather rough concept, and there does not exist a widely accepted definition of communities.

In this paper, we try to axiomatize the concept of communities. We allow overlapping communities and assume the data is given in a preference network.

Kleinberg [2003] developed an axiomatic framework for non-overlapping clusterings and defined clustering functions

whose inputs are a set V of n individuals and pairwise distances between these n individuals and whose output is a partition of these vertices (i.e., overlapping communities are not allowed). He found that three desired clustering axioms (scale-invariance, richness and consistency) cannot be satisfied at the same time.

But this negative result did not prevent researchers from investigating axiomatic aspects of clusterings. Meil [2005] proposed axioms to compare clusterings. Ben-David and Ackerman [2008] introduced axioms on clustering quality measures. Zadeh and Ben-David [2009] presented a unique theorem for clustering and showed that the impossibility result in [Kleinberg, 2003] can be avoided by relaxing one of clustering axioms. Ackerman *et al.* [2010] introduced several more properties of clustering functions to taxonomize a set of popular clustering algorithmic paradigms and strengthened the impossibility result on these properties. Gollapudi and Sharma [2009] devised an evaluation method to characterize the axioms.

Following this line but in a more general sense, Borgs *et al.* [2016] recently initiated the axiomization of overlapping communities over preference networks. A preference network is a finite set of individuals each having a preference indicating his/her affinity with others. As an expressive model, it covers graphs, a structure that is commonly used in the field of community detection. Borgs *et al.* [2016] proposed a system of six axioms for community functions and showed various properties of this system.

Among the properties, the lattice structure and the intersection lemma are most striking, because they serve as a guideline to construct desirable community functions. Specifically, the lattice structure means that the axioms satisfying all the six axioms form a lattice under the natural “ \cup ” and “ \cap ” operations. Borgs *et al.* also identified the bottom (the minimum, called \mathcal{C}_{cliq}) and the top (the maximum, called \mathcal{C}_{comp}) of this lattice. The intersection lemma claims that a community function satisfies all the axioms if and only if it is the intersection of \mathcal{C}_{comp} and a community function satisfying *four* axioms which are simple and easy to conform. The significance of the intersection lemma is that it paves a way to construct a

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community function satisfying all the axioms, while the lattice structure enables to derive more such community functions on this basis.

However, Borgs *et al.* [2016] did not *construct* any non-trivial community function satisfying all the axioms. This is mainly due to the fact that it is coNP-complete to check two of the axioms, namely Group Stability (*GS*) and Self-Approval (*SA*). Furthermore, *GS* and *SA* play critical roles in this axiomization for two reasons. First, only they cover the stability of communities, while an unstable community does not make sense, especially in the context of preference networks. As an analogue, consider the stable matching problem where affinity among the players is determined by their preferences and the stability of matchings is the focus (see e.g., Gusfield and Irving [1989]). Second, these two axioms set the top of the lattice. The coNP-completeness of the two most important axioms surely compromises the practicality of the axiom system.

On this ground, we make an effort to improve the axiom system in [Borgs *et al.*, 2016] and construct desired community functions. Our contributions are fourfold.

- 1) We show that axiom *SA* can degenerate so as to be checkable in polynomial time, leaving the axiom system equivalent.
- 2) We propose an efficiently checkable alternative axiom *SGS* to the original *GS*. The modification is natural, and preserves all the good properties of the original axiom system.
- 3) We present a general approach to construct nontrivial community functions that conform to the new axiom systems.
- 4) We find a natural, consistent, constructive, enumerable, and samplable community function, answering an open problem in [Borgs *et al.*, 2016]. Note that both this community function and the just-mentioned general approach remain valid in the original axiom system, but the construction essentially depends on the new axiom *SGS*.

The structure of this paper is as follows. In Section II, we review some necessary notions and notations, and show that *SA* can be replaced by a simpler axiom. We define a strengthening alternative to the crucial axiom *GS* in Section III, named *SGS*. In Section IV, we prove that *SGS* is efficiently checkable, and various desirable properties of the original axiom system are prereserved. In Section V, we present a systematic approach to finding consistent and constructive community functions, and find an ideal community function. Finally, we conclude this paper in Section VI with future works and open problems.

II. PRELIMINARIES

In this paper, we consider community detection as a task that translates a preference network into a set of communities. First of all, we review the definition of preference networks and the axiom system introduced by Borgs *et al.* [2016], and show some properties of the axiom system.

A. Preference networks

A preference network is a finite set of individuals each of which ranks all the individuals (including herself) in (strict) order of preference.

Formally, consider a non-empty set $V = \{1, 2, \dots, n\}$ and $L(V)$, the set of all total orders on V .

A total order $\pi \in L(V)$ is equivalently defined as a bijection $\pi : V \rightarrow V$, denoted by $[v_1 v_2 \dots v_n]$ where $v_i = \pi^{-1}(i)$ for $i = 1, 2, \dots, n$. For any $u, v \in V$ and $\pi \in L(V)$, $\pi(u)$ is referred to as the rank of u in π , and we say that π prefers u to v if $\pi(u) < \pi(v)$, denoted by $u \succ_{\pi} v$.

A *preference profile* Π on V is a list of total orders $\{\pi_u\}_{u \in V} \in L(V)^V$ on V . Given $u, v, w \in V$, we say that u prefers v to w if so does π_u , denoted by $v \succ_{\pi_u} w$. Given a preference profile Π on V , the pair $N = (V, \Pi)$ is called a *preference network*.

A *community function* is a function \mathcal{C} that maps each preference network $N = (V, \Pi)$ to a collection $\mathcal{C}(N) \subseteq 2^V$. Each $S \in \mathcal{C}(N)$ is called a community of N defined by \mathcal{C} , and the term “defined by \mathcal{C} ” will be omitted if clear from context.

As an example, consider the community function

$$\mathcal{C}_{cliq}(V, \Pi) = \{S \subseteq V : \forall u, v \in S, \forall w \in V - S, v \succ_{\pi_u} w\}.$$

Every community S defined by \mathcal{C}_{cliq} is called a clique¹. Roughly speaking, a clique is a subset of individuals preferring each other to those not in the subset.

Example 1: We consider two families in a village and each family has three members. The members in a family have the same preference and they prefer their own family members to other villagers. In this case, every family is a community. Formally, suppose $N = (V, \Pi)$ is the preference network where $V = \{1, 2, 3, 4, 5, 6\}$, $\pi_1 = \pi_2 = \pi_3 = [123456]$ and $\pi_4 = \pi_5 = \pi_6 = [456123]$. Sets $\{1, 2, 3\}$ and $\{4, 5, 6\}$ are in $\mathcal{C}_{cliq}(N)$.

Actually preference frameworks have been used in various occasions, e.g., voting [Gale and Shapley, 1962; Gusfield and Irving, 1989; Roth, 1984a], network routing [Rekhter and Li, 1994; Caesar and Rexford, 2005], and coalition formation in collaborative games [Brams *et al.*, 2003; Roth, 1984b]. Balcan *et al.* [2013] further elaborated on how the concept of preference networks properly models social networks and why it essentially captures their underlying community structures.

B. Existing axioms

Borgs *et al.* [2016] proposed six axioms for community functions. They are included here for the paper to be self-containing. Let’s arbitrarily fix a non-empty finite set V and a community function \mathcal{C} .

Axiom 1 (Anonymity(A)): Let $S, S' \subseteq V$, and Π, Π' be two preference profiles on V . Assume that $S' = \sigma(S)$ and $\Pi' = \sigma(\Pi)$ for some permutation $\sigma : V \rightarrow V$. Then $S \in \mathcal{C}(N)$ if and only if $S' \in \mathcal{C}(N')$, where $N = (V, \Pi)$ and $N' = (V, \Pi')$.

¹Note that a clique in a preference network does not mean it is a complete subgraph.

This axiom means that whether S is a community is determined by the preference profile rather than by labels of the individuals.

Axiom 2 (Monotonicity(Mon)): Let $S \subseteq V$, and Π, Π' be two preference profiles on V . If $u \succ_{\pi_s} v \Rightarrow u \succ_{\pi'_s} v$ for $\forall s, u \in S$ and $\forall v \in V$, then $S \in \mathcal{C}(N) \Rightarrow S \in \mathcal{C}(N')$, where $N = (V, \Pi)$ and $N' = (V, \Pi')$.

Intuitively, monotonicity means that if the change of the preference profile does not decrease the ranking of any member of a community and preserves the relative ranking among the members, then the community remains. This is reasonable since the affinity among the members of the community is improved after the preference profile is changed.

Axiom 3 (Embedding(Emb)): Given two preference networks $N = (V, \Pi)$ and $N' = (V', \Pi')$ such that $V' \subseteq V$ and $\pi'_u(v) = \pi_u(v)$ for all $u, v \in V'$, then $\mathcal{C}(N') = \mathcal{C}(N) \cap 2^{V'}$.

Intuitively, though V' is embedded into V , its preferences are not influenced. Hence it is reasonable that communities inside V' are formed independently of individuals outside V' .

Axiom 4 (World Community(WC)): For any preference network $N = (V, \Pi)$, $V \in \mathcal{C}(N)$.

This axiom is self-explanatory.

Before introducing the remaining two axioms, the concept of *preference* should be extended to *group-preference*² so that equal-sized sets can be compared. Given a preference network $N = (V, \Pi)$ and non-empty disjoint sets $G, G' \subseteq V$ of the same size, we say that $s \in V$ *group-prefers* G' to G , denoted by $G' \succ_{\pi_s} G$, if one can align the elements $g_1, g_2, \dots, g_{|G|}$ of G and $g'_1, g'_2, \dots, g'_{|G|}$ of G' so that $g'_i \succ_{\pi_s} g_i$ for all $i = 1, 2, \dots, |G|$.

For example, in the preference network in Example 1, $\{3, 4\} \succ_{\pi_2} \{5, 6\}$ since $3 \succ_{\pi_2} 5$ and $4 \succ_{\pi_2} 6$. But $\{3, 4\} \not\succ_{\pi_5} \{5, 6\}$ does not hold, because $5 \succ_{\pi_5} 3$ and $6 \succ_{\pi_5} 4$.

A set $S \subseteq V$ is called *self-approving* with respect to Π if for any $S' \subseteq V - S$ with $|S'| = |S|$, S' is not group-preferred to S by some $s \in S$. S is called *group-stable* with respect to Π if for any non-empty $G \subsetneq S$ and $G' \subseteq V - S$ with $|G'| = |G|$, G' is not group-preferred to G by some $s \in S - G$. Roughly speaking, a set is group-stable if no subset agrees to leave and join another set.

For instance, consider again the preference network in Example 1. The set $S = \{2, 5\}$ is not group-stable, since there exist $G = \{2\}$ and $G' = \{4\}$ such that $5 \in S - G$ group-prefers G' to G . Actually, S is not *self-approving* either, because for any $s \in S$, s group-prefers $\{1, 4\}$ to S .

Axiom 5 (Self-Approval(SA)): For any preference network $N = (V, \Pi)$, if $S \in \mathcal{C}(N)$, then S is self-approving with respect to Π .

Axiom 6 (Group Stability(GS)): Given a preference network $N = (V, \Pi)$, if $S \in \mathcal{C}(N)$, then S is group-stable with respect to Π .

There exist community functions satisfying all the axioms, e.g., \mathcal{C}_{cliq} and

$$\mathcal{C}_{comp}(N) = \{S \subseteq V : S \text{ is self-approving and group-stable}\}.$$

C. Properties

Borgs *et al.* [2016] showed that this axiom system has nice properties. The most impressive ones include the intersection lemma and the lattice structure.

Lemma 1 (Lattice, [Borgs et al., 2016]): Under the operations \cap and \cup , the community functions satisfying A , Mon , Emb , WC , GS , and SA form a lattice whose top is \mathcal{C}_{comp} and bottom is \mathcal{C}_{cliq} .

Due to the lattice structure, given some community functions satisfying all the six axioms, one can construct more such community functions using a mixture of \cap and \cup operations.

Lemma 2 (Intersection Lemma, [Borgs et al., 2016]): For any community function \mathcal{C} satisfying A , Mon , Emb and WC , $\mathcal{C} \cap \mathcal{C}_{comp}$ satisfies all the six axioms.

By this lemma, designing a community function satisfying all the axioms is reduced to find one satisfying the four axioms, which is relatively easier to be satisfied. However, even we have a community function satisfying the four axioms, it remains hard to check whether a given set is a community, since it is computationally hard to check the key axioms GS and SA .

Theorem 3 ([Borgs et al., 2016]): It is coNP-complete to decide whether a subset $S \subseteq V$ is self-approval or group-stable with respect to a preference profile.

Actually, we find that SA can degenerate so as to be efficiently checkable, keeping the axiom system equivalent to the original one. This is because GS almost implies SA , which immediately follows from the next theorem.

Theorem 4: Given a preference network $N = (V, \Pi)$ and a community $S \subseteq V$ with $|S| \geq 2$, if S is group-stable with respect to Π , then it is self-approving with respect to Π .

Proof. Suppose that $|S| \geq 2$ and S is not self-approving with respect to Π . Then there exists $S' \subseteq V \setminus S$ such that for all individuals $s \in S$, s group-prefers S' to S . Fix such a S' and arbitrarily choose $s \in S$. Assume that $s' \in S'$ is the least favorite individual in S' by the preference of s . We have that s group-prefers $S' - \{s'\}$ to $S - \{s\}$.

Since $|S| \geq 2$, $S - \{s\}$ is a non-empty set. It follows that S is not group-stable with respect to Π . \square

Therefore, we present a degenerate version of SA so that it only deals with the case where $|S| = 1$. The degenerate version is

SA' : For any preference network $N = (V, \Pi)$, if $\{s\} \in \mathcal{C}(N)$, then $\pi_s(s) = 1$.

By Theorem 4, any community function conforms with $\{GS, SA\}$ if and only if it conforms with $\{GS, SA'\}$. Considering that SA' is efficiently checkable, we will use SA' instead of SA in the remainder of the paper.

III. STRENGTHENING AXIOM GS

After replacing SA with SA' , axiom GS becomes the only difficult-to-check axiom in the axiom system. It is this

²It is referred to as lexicographic preference in [Borgs et al., 2016].

intractability that causes difficulty in constructing community functions. In order to solve this problem, we modify GS in a natural way, resulting in an efficient-to-check axiom called SGS . Besides, we show that a community function satisfies GS if it satisfies SGS .

A. Weak preferences

The hardness of checking group stability of a subset S is partially rooted at enumerating equal-sized subsets of S and $V - S$. A natural idea for tackling this problem is to compare subsets of S directly with $V - S$, rather than with its subsets. The first technical obstacle is to compare subsets of different sizes, so we further extend the concept of preference.

Definition 1 (Weak Preferences): Suppose $S, S' \subseteq V$, $\pi \in L(V)$ and $k = \min(|S|, |S'|)$. The preference profile π *weakly prefers* S to S' , denoted by $S \succ_{\pi} S'$, if π group-prefers the set of the top k elements of S to that of S' .

Consider Example 1 again. We can show that individual 4 weakly prefers $\{4, 1\}$ to $\{2, 3, 5\}$. Specifically, since $\{2, 3, 5\}$ has three individuals and $\{4, 1\}$ only has two, select the top two individuals $\{2, 5\}$ from $\{2, 3, 5\}$ (according to the preference of 4). The result follows because 4 group-prefers $\{4, 1\}$ to $\{2, 5\}$.

It is reasonable to borrow the notation “ \succ ” from group-preference, since weak preference is equivalent to group-preference if $|S| = |S'|$. When $|S| > |S'|$, $S \succ_{\pi} S'$ means that there is $T \subsetneq S$ with $|T| = |S'|$ such that π group-prefers T to S' . When $|S| < |S'|$, $S \succ_{\pi} S'$ means that π group-prefers S to any $T \subsetneq S'$ with $|S| = |T|$.

The following properties of weak preferences will be frequently used in the rest of this paper. They immediately follow from the definition of weak preference, so the proofs are omitted.

Lemma 5: Given $\pi \in L(V)$, for any $S \subseteq V$ and $S' \subseteq U \subseteq V$ with $|S'| \geq |S|$, if π weakly prefers S' to S , then π weakly prefers U to S .

Lemma 6: Given $\pi \in L(V)$, for any $S, S' \subseteq V$, if π weakly prefers S' to S , then for all $T \subseteq S$, π weakly prefers S' to T .

In Example 1, suppose $S = \{2, 4\}$ and $S' = \{1, 3\}$, then individual 4 weakly prefers S' to S . Suppose $U = \{1, 3, 5\}$ and $T = \{2\}$, then $S' \subseteq U$ and $T \subseteq S$, which means that individual 4 weakly prefers U to S by Lemma 5 and weakly prefers S' to T by Lemma 6.

B. Alternative to GS

Now we are ready to define the alternative SGS (namely strong group stability) to GS .

Consider a preference network $N = (V, \Pi)$. A set $S \subseteq V$ is called *strongly group-stable* with respect to Π if for any non-empty $G \subsetneq S$, there exists $s \in S - G$ which does not weakly prefer $V - S$ to G . Strong group-stability differs from group-stability mainly in that weak preference rather than group-preference is used.

Let's arbitrarily fix a community function \mathcal{C} .

Axiom 7 (SGS): For any preference network $N = (V, \Pi)$ and any subset $S \subseteq V$, if $S \in \mathcal{C}(N)$, then S is strongly group-stable with respect to Π .

Roughly speaking, GS rules out a community S if there exists a subset $T \subsetneq S$ such that all members of $S - T$ agree to replace T with a common, equal-sized group outside of S , while SGS does so under a weaker condition that essentially captures two facts. First, T can be replaced by a group of a different size. Second, members of $S - T$ might not have a common replacement of T (in case that $|V - S| > |T|$), but they agree on kicking T out. In this case, it is reasonable to rule out such communities; see the following example.

In Example 1, individual 4 weakly prefers $\{5, 6\}$ to $\{1, 2, 3\}$, which means that $\{1, 2, 3, 4\}$ is not strongly group-stable and cannot be a community according to SGS . This coincides with commonsense, since 4 tends to leave $\{1, 2, 3\}$ and join $\{5, 6\}$. However, $\{1, 2, 3, 4\}$ is group-stable and is a candidate of community in the original axiom system.

Theorem 7: For any community function \mathcal{C} , if it satisfies SGS , then it also satisfies GS .

Proof. If some community function \mathcal{C} satisfies SGS but does not satisfy GS , there must exist a preference network $N = (V, \Pi)$ and a community $S \in \mathcal{C}(N)$ which is strongly group-stable but not group-stable with respect to Π . By the definition of group stability, there are two non-empty subsets $T \subsetneq S$ and $T' \subseteq V - S$ such that $|T'| = |T|$ and u group-prefers T' to T for all individuals $u \in S - T$. By Lemma 5, u weakly prefers $V - S$ to T , meaning that S is not strongly group-stable. A contradiction is reached. \square

IV. COMPLEXITY AND LATTICE

This section shows that the lattice structure, the intersection lemma, and richness of the axiom system in Borgs *et al.* [2016] still hold in our axiom system. More importantly, we show that it takes polynomial time to check whether a given community satisfies SGS .

A. Intersection lemma and lattice structure

Let $\mathcal{A} = \{A, Mon, WC, Emb, GS, SA'\}$ and $\mathcal{SA} = \{A, Mon, WC, Emb, SGS, SA'\}$. For a collection $\mathcal{X} \in \{\mathcal{A}, \mathcal{SA}\}$ of axioms, a community function is said to be \mathcal{X} -consistent if it conforms with all axioms in \mathcal{X} .

We go on with showing that many good properties of \mathcal{A} are preserved by \mathcal{SA} . An example is the intersection lemma, one of the most important properties of \mathcal{A} .

Before introducing the intersection lemma, we define the intersection between community functions. Given two community functions \mathcal{C}_1 and \mathcal{C}_2 , we define $\mathcal{C}_1 \cap \mathcal{C}_2$ as the community function that $(\mathcal{C}_1 \cap \mathcal{C}_2)(N) = \mathcal{C}_1(N) \cap \mathcal{C}_2(N)$. Similarly, one can define $(\mathcal{C}_1 \cup \mathcal{C}_2)(N) = \mathcal{C}_1(N) \cup \mathcal{C}_2(N)$.

Then we define three community functions that will be used: \mathcal{C}_{SGS} , $\mathcal{C}_{SA'}$, and \mathcal{C}_{scomp} , which are the maximum community functions satisfying SGS , SA' , and both of them, respectively. Formally, given a preference network $N = (V, \Pi)$,

$$\mathcal{C}_{SGS}(N) = \{S \subseteq V : S \text{ is strongly group-stable}\},$$

$$\mathcal{C}_{SA'}(N) = \{S \subseteq V : \text{if } S = \{s\} \text{ then } \pi_s(s) = 1\}, \text{ and}$$

$$\mathcal{C}_{scomp} = \mathcal{C}_{SGS} \cap \mathcal{C}_{SA'}.$$

Lemma 8 (Intersection Lemma): If a community function \mathcal{C} satisfying any axiom $X \in \{A, Mon, WC, Emb\}$, then $\mathcal{C} \cap \mathcal{C}_{scomp}$ satisfies X , SGS and SA' .

Proof. Let $\tilde{\mathcal{C}} = \mathcal{C} \cap \mathcal{C}_{scomp}$. If $X \in \{A, WC, Emb\}$, $\tilde{\mathcal{C}}$ satisfies X because \mathcal{C}_{scomp} is obviously $\{A, WC, Emb\}$ -consistent. Hence, we just consider the case where $X = Mon$.

Assume that preference networks $N = (V, \Pi)$, $N' = (V, \Pi')$, and community $S \in \tilde{\mathcal{C}}(N)$ satisfy the condition of Mon , namely, for all $s, u \in S$ and $v \in V$, if $u \succ_{\pi_s} v$ then $u \succ_{\pi'_s} v$. The rest of the proof consists of two steps.

Step 1: Since \mathcal{C} satisfies Mon and $S \in \mathcal{C}(N)$, we immediately have $S \in \mathcal{C}(N')$.

Step 2: We prove that $S \in \mathcal{C}_{scomp}(N')$ as following.

If S is a singleton $\{s\}$, then $S \in \mathcal{C}_{scomp}(N)$ implies that $\pi_s(s) = 1$, namely $s \succ_{\pi_s} u$ for any $u \neq s$. Hence, $s \succ_{\pi'_s} u$ for any $u \neq s$ and we have $S \in \mathcal{C}_{scomp}(N')$.

When $|S| > 1$, $S \in \mathcal{C}_{scomp}(N')$ if and only if $S \in \mathcal{C}_{SGS}(N')$, so we only have to prove that $S \in \mathcal{C}_{SGS}(N')$. For contradiction, suppose $S \notin \mathcal{C}_{SGS}(N')$. Then there exists a non-empty set $T \subsetneq S$ such that $V - S \succ_{\pi'_s} T$ for all $s \in S - T$. Thus, for all $s \in S - T$ and $1 \leq i \leq k$ where $k = \min\{|T|, |V - S|\}$, the top- i -th individual $u_i \in V - S$ and the top- i -th individual $v_i \in T$ satisfy $u_i \succ_{\pi'_s} v_i$. Recalling the assumption about N, N', S , we have $u_i \succ_{\pi_s} v_i$ for all $1 \leq i \leq k$ and $s \in S - T$, meaning that $V - S \succ_{\pi_s} T$ for all $s \in S - T$. This indicates that $S \notin \mathcal{C}_{SGS}(N)$. A contradiction is reached. Therefore, $S \in \mathcal{C}_{scomp}(N')$.

Altogether, we have shown that if N, N' and $S \in \tilde{\mathcal{C}}(N)$ satisfy the condition of Mon , then $S \in \tilde{\mathcal{C}}(N')$. This means that $\tilde{\mathcal{C}}$ satisfies Mon , and the lemma follows immediately. \square

In some sense, the critical role of the intersection lemma lies in that it provides a normal form of consistent community functions: the intersection of a simple community function with \mathcal{C}_{scomp} . Hence finding a SA -consistent community function is reduced to finding a $\{A, Mon, WC, Emb\}$ -consistent one. This reduction will greatly help us construct desirable community functions, as shown in Section V-B. The intersection lemma also leads to the lattice structure of SA -consistent community functions, which is also true for \mathcal{A} and is one of the most striking results in [Borgs et al., 2016].

Theorem 9: Let \mathbb{C} be the class of all SA -consistent community functions. The algebraic structure $\mathcal{T} = \{\mathbb{C}, \cup, \cap\}$ forms a bounded lattice whose top and bottom are \mathcal{C}_{scomp} and \mathcal{C}_{cliq} , respectively.

The proof of Theorem 9 is straightforward by the intersection lemma (Lemma 8) and the proof of Theorem 3.1 in Borgs et al. [2016], so the detail is omitted.

The significance of the lattice structure lies in that given some SA -consistent community functions, one can produce more such community functions by a mixture of \cap and \cup operations.

According to the lattice structure, the richness of SA -consistent community functions is to some extent determined by the difference between \mathcal{C}_{scomp} and \mathcal{C}_{cliq} . Actually, the difference is significant, with evidence from the number of communities. On the one hand, for any preference network

N , the size of $\mathcal{C}_{cliq}(N)$ is linear, by Lemma 19. On the other hand, an example inspired by [Borgs et al., 2016] indicates that the size $\mathcal{C}_{scomp}(N)$ can be exponential for some preference network N .

Lemma 10: For any positive integer n , there exists a preference network $N = (V, \Pi)$ with $|V| = n$ such that $|\mathcal{C}_{scomp}(N)| \geq 2^{\lfloor n/2 \rfloor}$.

Proof. Let $m = \lfloor n/2 \rfloor$. Let $H = \{h_1, h_2, \dots, h_m\}$ be a set of m heroes, $S = \{s_1, s_2, \dots, s_m\}$ be a set of m sidekicks, and $V = H \cup S$. For all $1 \leq i \leq m$, define preferences $\pi_{h_i} = \pi_{s_i} = [h_i, h_1, h_2, \dots, h_m, s_i, s_1, \dots, s_{i-1}, s_{i+1}, \dots, s_m]$. Let Π be the preference profile consisting of all π_{h_i} 's and π_{s_i} 's. Consider the preference network $N = (V, \Pi)$. Arbitrarily choose a subset $U \subseteq S$. We will prove that $H \cup U \in \mathcal{C}_{SGS}(N)$.

For contradiction, suppose $H \cup U \notin \mathcal{C}_{SGS}(N)$. Then there exists a non-empty set $T \subsetneq H \cup U$ such that $V - (H \cup U) \succ_{\pi_w} T$ for all $w \in (H \cup U) - T$. Since any individual in H is always preferred to any individual in S , it must hold that $T \subset U \subseteq S$, implying that $H \subseteq (H \cup U) - T$. Therefore, $V - (H \cup U) \succ_{\pi_h} T$ for any $h \in H$. Now choose $1 \leq i \leq m$ such that $s_i \in U$. We have $V - (H \cup U) \succ_{\pi_{h_i}} T$, contradictory to the fact that $s_i \succ_{\pi_{h_i}} v$ for any $v \in V - (H \cup U) \subseteq S$. Hence $H \cup U \in \mathcal{C}_{SGS}(N)$.

Since $|H \cup U| > 1$, by the definition of \mathcal{C}_{scomp} and $\mathcal{C}_{SA'}$, $H \cup U \in \mathcal{C}_{SGS}(N)$ implies that $H \cup U \in \mathcal{C}_{scomp}(N)$. Because there are 2^m different U 's, $|\mathcal{C}_{scomp}(N)| \geq 2^m$. \square

Since the size of $\mathcal{C}_{cliq}(N)$ is at most linear (proved in Lemma 19), Lemma 10 suggests that the collection of SA -consistent community functions may be rich and the lattice may be far from collapsing. The above lemma also implies that the though SGS is stronger than GS , the restriction is not too much.

B. SGS is efficient to check

Now we present the most important property of SGS , that it can be efficiently checked. This is shown constructively, through an algorithm inspired by the following observation.

Consider a preference network $N = (V, \Pi)$ and a subset $S \subsetneq V$. Suppose that $S \notin \mathcal{C}_{SGS}(N)$. There must exist a non-empty set $T \subsetneq S$ that all individuals in $S - T$ weakly prefer $V - S$ to T . If we extend any set $U \subseteq T$ by adding any $u \in S$ that does not weakly prefer $V - S$ to U , Lemma 6 indicates that no individual in $S - T$ can be added since all individuals in $S - T$ weakly prefer $V - S$ to U . So, if we happen to start with $U = \{u\}$ for some $u \in T$ and extend U step-by-step, the process must stop before $U = S$.

Based on this observation, we design an algorithm which for each individual $u \in S$, initializes U to be $\{u\}$, and iteratively extends U by adding an individual in S that does not weakly prefer $V - S$ to U . If and only if we get some $U \subsetneq S$ that cannot be extended any more, decide $S \notin \mathcal{C}_{SGS}(N)$. The details are specified in Algorithm 1.

We illustrate how Algorithm 1 works using the preference network in Example 1. Let S be $\{1, 2, 3, 4\}$. At the beginning, u is 1 at Line 2 and $U = \{1\}$. Individual 2 is added to U at Line 7 since individual 2 does not weakly prefer $V - S = \{5\}$

Algorithm 1 Decide membership of $\mathcal{C}_{SGS}(N)$

Input A preference network (V, Π) , a subset $S \subsetneq V$ **Output** A boolean value indicating whether $S \in \mathcal{C}_{SGS}(N)$

```
1: procedure DECIDE( $V, \Pi, S$ )
2:   for  $u \in S$  do
3:      $U \leftarrow \{u\}$ 
4:     while  $U \neq S$  do
5:       if  $\exists s \in S - U$  such that  $V - S \not\prec_{\pi_s} U$  then
6:         Arbitrarily choose such an  $s$ 
7:          $U \leftarrow U \cup \{s\}$ 
8:       else
9:         return False
10:  return True
```

to U . Then $U = \{1, 2\}$, and individual 3 is added to U since individual 3 does not weakly prefer $V - S = \{5\}$ to U . Now $U = \{1, 2, 3\}$ and $S - U = \{4\}$. Because individual 4 weakly prefers $\{5\}$ to U , the algorithm returns *False* at Line 9.

Now we prove the correctness of Algorithm 1 and show that its running time is $O(|S|^4)$.

Lemma 11: Given a preference network $N = (V, \Pi)$ and $S \subseteq V$, it is in time $O(|S|^4)$ to decide whether $S \in \mathcal{C}_{SGS}(N)$.

Proof. Correctness:

Suppose $S \notin \mathcal{C}_{SGS}(N)$. There must exist a non-empty set $T \subsetneq S$ that $V - S \succ_{\pi_s} T$ for all $s \in S - T$. Note that the *For* loop enumerates all $u \in S$. If the algorithm terminates before any $u \in T$ is chosen at line 3, it surely outputs *False*. Otherwise, we eventually have $u \in T$ at line 3. By Lemma 6, $U \subseteq T$ always holds. Hence, the *While* loop will reach a state where $U \neq S$ and the condition at Line 5 is false. Then, the algorithm also returns *False*.

On the other hand, if the algorithm returns *False*, all individuals in $S - U$ weakly prefer $V - S$ to U , which means that $S \notin \mathcal{C}_{SGS}(N)$.

Complexity:

There are two nested loops in Algorithm 1 and each loop runs at most $|S|$ times. The complexity of checking the condition at line 5 is $O(|S|^2)$. So, the complexity of Algorithm 1 is $O(|S|^4)$. \square

We are ready to show that \mathcal{C}_{scomp} is efficiently checkable, in contrast to the coNP-completeness of checking the membership of \mathcal{C}_{comp} .

Theorem 12: Given a preference network $N = (V, \Pi)$ and $S \subseteq V$, it takes $O(|S|^4)$ time to decide whether $S \in \mathcal{C}_{scomp}(N)$ or not.

Proof. If $|S| \geq 2$, it is equivalent to decide whether $S \in \mathcal{C}_{SGS}(N)$, which can be done in $O(|S|^4)$ by Algorithm 1.

By Lemma 11, it takes $O(|S|^4)$ time to decide whether $S \in \mathcal{C}_{SGS}$, which means that it is in $O(|S|^4)$ to decide whether $S \in \mathcal{C}_{scomp}$.

Otherwise, S is a singleton $\{s\}$. In this case, $S \in \mathcal{C}_{scomp}(N)$ if and only if $\pi_s(s) = 1$, which can be decided in constant time. \square

V. CONSTRUCTING COMMUNITY FUNCTIONS

Borgs *et al.* [2016] proposed an open problem: Is there an ideal community function in their axiom system \mathcal{A} ? In this section, we first solve the problem in the our axiom system \mathcal{SA} by constructing such a community function. Since \mathcal{SA} is stronger than \mathcal{A} , this community function is also \mathcal{A} -consistent, answering the original open problem.

A. Consistent, constructive community functions

A community function \mathcal{C} is said to be *constructive* if for any preference network $N = (V, \Pi)$, the membership of $\mathcal{C}(N)$ can be checked in polynomial-time in $|V|$. Constructive community functions are desirable, but it is not easy to figure out a natural, nontrivial, \mathcal{A} -consistent one (*nontrivial* means different from \mathcal{C}_{cliq}). Borgs *et al.* [2016] mainly considered two families of candidates, namely $\mathcal{C}_{cliq(g)}$ and $\mathcal{C}_{harmonious(\lambda)}$ with non-negative function g and real number $\lambda \in [0, 1]$. They showed that though $\mathcal{C}_{cliq(g)} \cap \mathcal{C}_{comp}$ and $\mathcal{C}_{harmonious(\lambda)} \cap \mathcal{C}_{comp}$ are \mathcal{A} -consistent, they are not constructive in general. The definitions of $\mathcal{C}_{cliq(g)}$ and $\mathcal{C}_{harmonious(\lambda)}$ are presented here in order to make this paper self-contained.

Definition 2 ($\mathcal{C}_{cliq(g)}$): Given a non-negative function $g : \{1, 2, 3, \dots\} \rightarrow \{0, 1, 2, 3, \dots\}$, for any preference network $N = (V, \Pi)$ and $S \subseteq V$, $S \in \mathcal{C}_{cliq(g)}(N)$ if $\forall u, s \in S, \pi_s(u) \in [1 : |S| + g(|S|)]$.

Definition 3 ($\mathcal{C}_{harmonious(\lambda)}$): Given $\lambda \in [0, 1]$, for any preference network $N = (V, \Pi)$ and $S \subseteq V$, $S \in \mathcal{C}_{harmonious(\lambda)}(N)$ if $\forall u \in S, v \in V - S$, at least a λ -fraction of $\{\pi_s : s \in S\}$ prefer u to v .

In the context of \mathcal{SA} , we have the following theorem which helps finding consistent constructive community function.

Theorem 13: If a constructive community function \mathcal{C} is $\{A, Mon, Emb, WC\}$ -consistent, then $\mathcal{C} \cap \mathcal{C}_{scomp}$ is \mathcal{SA} -consistent and constructive. \square

Proof. The theorem follows immediately from Lemma 8 and Theorem 12. \square

This theorem greatly simplifies the task of finding \mathcal{SA} -consistent constructive community functions, since it is relatively easier to find a constructive one that is $\{A, Mon, Emb, WC\}$ -consistent.

For example, for any non-negative function g and real number $\lambda \in [0, 1]$, both $\mathcal{C}_{cliq(g)}$ and $\mathcal{C}_{harmonious(\lambda)}$ conform with A, Mon, Emb , and WC . Hence, Theorem 13 immediately implies the following lemma.

Lemma 14: The community functions $\mathcal{C}_{cliq(g)} \cap \mathcal{C}_{scomp}$ and $\mathcal{C}_{harmonious(\lambda)} \cap \mathcal{C}_{scomp}$ are \mathcal{SA} -consistent and constructive, for any non-negative function g and any real number $\lambda \in [0, 1]$.

Furthermore, by Theorem 7, the two families of community functions in Lemma 14 are also \mathcal{A} -consistent and constructive.

In the next subsection we go further in this direction and find a community function having more nice properties, solving the open problem posed by Borgs *et al.* [2016].

B. \mathcal{C}_{grow} : an ideal community function

An ideal community function should not only be constructive, but also allows to efficiently enumerate the communities.

Borgs *et al.* [2014] (the full version of [Borgs *et al.*, 2016]) defined two more properties for community functions \mathcal{C} :

- **Samplable:** given a preference network $N = (V, \Pi)$, one can randomly sample any community from $\mathcal{C}(N)$ in time $P(|V|)$, where $P(\cdot)$ is a polynomial function;
- **Enumerable:** given a preference network $N = (V, \Pi)$, one can enumerate $\mathcal{C}(N)$ in time $O(n^k |\mathcal{C}(N)|)$ for some constant k .

Borgs *et al.* [2016] proposed the following open problem.

Problem 1 (ACCSE): Find a natural nontrivial community function that is \mathcal{A} -consistent, constructive, samplable, and enumerable.

We solve the ACCSE problem in this subsection. Actually, the following stronger problem is solved.

Problem 2 (SACCSE): Find a natural community function that allows nontrivially overlapping communities and is \mathcal{SA} -consistent, constructive, samplable, and enumerable.

The basic idea of our solution is to reduce the problem according to Theorem 13. Namely, if we find a community function \mathcal{C} satisfying A , Mon , WC , and Emb , then $\mathcal{C} \cap \mathcal{C}_{scomp}$ is \mathcal{SA} -consistent. Furthermore, if \mathcal{C} is enumerable and $|\mathcal{C}(N)|$ is polynomial in $|V|$ for all preference networks $N = (V, \Pi)$, $\mathcal{C} \cap \mathcal{C}_{scomp}$ will be enumerable and samplable.

But how to find such a community function \mathcal{C} ? Our approach is inspired by the method in [Palla *et al.*, 2005]. Roughly speaking, starting with best communities, i.e., cliques in a preference network, we produce new good communities as many as possible by extending existing communities. Intuitively, this growing process conforms with the formation of communities in real life (for example, consider how friendship forms). The key of the approach is properly defining good communities. For this end, we consider the following community function which is adapted from $\mathcal{C}_{harmonious(\lambda)}$.

Definition 4 (\mathcal{C}_{harmon}): Given a preference network $N = (V, \Pi)$ and a subset $S \subseteq V$, $S \in \mathcal{C}_{harmon}(N)$ if for all $u \in S$ and $v \in V - S$, more than half of the members of $S - \{u\}$ prefer u to v .

The following lemma shows that \mathcal{C}_{harmon} is efficiently checkable.

Lemma 15: Given a preference network $N = (V, \Pi)$ and a set $S \subseteq V$, it takes $O(|V|^3)$ time to determine whether S is in $\mathcal{C}_{harmon}(N)$.

Proof. For each individual u in S and each individual v in $V - S$, it takes $O(|S|)$ time to determine whether a majority of $S - \{u\}$ prefer u to v , which means that it takes $O(|V|^3)$ time to determine whether $S \in \mathcal{C}_{harmon}(N)$ or not. \square

Based on the idea preceding Definition 4, we define \mathcal{C}_{grow} in terms of Algorithm 2, i.e., for any preference network N , $\mathcal{C}_{grow}(N) := \text{CLIQUEGROWING}(N)$.

Algorithm 2 is self-explanatory. We now show that the size of $\mathcal{C}_{grow}(N)$ and the time complexity of Algorithm 2 are both polynomial in $|V|$.

Lemma 16: For any preference network $N = (V, \Pi)$, the size of $\mathcal{C}_{grow}(N)$ is $O(|V|^2)$, and Algorithm 2 terminates within time $O(|V|^6)$.

Algorithm 2 The definition of \mathcal{C}_{grow}

Input A preference network $N = (V, \Pi)$

Output A collection R of communities

```

1: procedure CLIQUEGROWING( $N$ )
2:    $C \leftarrow \mathcal{C}_{cliq}(N)$        $\triangleright C$  keeps the communities to be
   extended
3:    $R \leftarrow \emptyset$          $\triangleright R$  keeps the final communities
4:   while  $C$  is not empty do
5:     Arbitrarily choose an element  $S$  from  $C$ 
6:      $C \leftarrow C \setminus \{S\}$ ,  $R \leftarrow R \cup \{S\}$ 
7:     for  $u$  in  $V - S$  do
8:        $S' \leftarrow S \cup \{u\}$ 
9:       if  $S' \in \mathcal{C}_{harmon}(N)$  then
10:         $C \leftarrow C \cup \{S'\}$ 
11:   return  $R$ 

```

Proof. According to the definition of \mathcal{C}_{harmon} , given $S \subseteq V$, if $S \cup \{u\} \in \mathcal{C}_{harmon}(N)$ for some $u \in V - S$, then for any $v \in V - S$ with $v \neq u$, a majority of individuals in S prefer u to v . As a result, for any $v \in V - S$ with $v \neq u$, $S \cup \{u\} \notin \mathcal{C}_{harmon}(N)$. This means that for any $S \in C$, it can be extended by at most one $u \in V - S$ in the **for** loop. Hence, starting with any $S \in \mathcal{C}_{cliq}$, at most $O(|V|)$ communities can be obtained. Since there are only $O(|V|)$ cliques by Lemma 19, the size of $\mathcal{C}_{grow}(N)$ is $O(|V|^2)$.

Then we discuss the time complexity of Algorithm 2. The **If** statement at line 9 is executed for $|V - S|$ times for each community S . This fact, together with Lemma 15, implies that the time complexity of Algorithm 2 is $O(|V|^6)$. \square

Lemma 16 immediately indicates that \mathcal{C}_{grow} is enumerable, samplable, and constructive. However, \mathcal{C}_{grow} is not \mathcal{SA} -consistent. For example, in the preference network of Example 1, the set $\{1, 2, 3, 4\}$ is a community defined by \mathcal{C}_{grow} , but it is not strongly group-stable since individual 4 weakly prefers $\{5, 6\}$ to $\{1, 2, 3\}$.

Fortunately, $\mathcal{C}_{grow} \cap \mathcal{C}_{scomp}$ is a community function satisfying all the requirements in Problem 2.

Theorem 17: $\mathcal{C}_{grow} \cap \mathcal{C}_{scomp}$ is a natural community function that allows nontrivially overlapping communities and solves both the SACCSE problem and the ACCSE problem.

Proof. By Lemma 16, we know that \mathcal{C}_{grow} is enumerable, samplable, and constructive, and that the size of $\mathcal{C}_{grow}(N)$ is polynomial for any preference network N . Since \mathcal{C}_{scomp} is constructive, $\mathcal{C}_{grow} \cap \mathcal{C}_{scomp}$ is enumerable, samplable, and constructive.

It is straightforward to show that \mathcal{C}_{grow} satisfies A , Mon , WC , and Emb . According to Lemma 8, $\mathcal{C}_{grow} \cap \mathcal{C}_{scomp}$ is \mathcal{SA} -consistent.

By Lemmas 19 and 20, $\mathcal{C}_{grow} \cap \mathcal{C}_{scomp} \neq \mathcal{C}_{cliq}$, namely, $\mathcal{C}_{grow} \cap \mathcal{C}_{scomp}$ is nontrivial.

As a result, $\mathcal{C}_{grow} \cap \mathcal{C}_{scomp}$ solves the SACCSE problem. By Theorem 7, it also solves the ACCSE problem. \square

Now we prove Lemmas 19 and 20 just used. First, recall a property of \mathcal{C}_{cliq} .

Lemma 18 ([Borgs et al., 2016]): Given a preference network $N = (V, \Pi)$ and $S_1, S_2 \in \mathcal{C}_{cliq}(N)$, then either $S_1 \subseteq S_2$, $S_2 \subseteq S_1$, or $S_1 \cap S_2 = \emptyset$.

Lemma 19: For any preference network $N = (V, \Pi)$, the size of $\mathcal{C}_{cliq}(N)$ is $O(|V|)$.

Proof. For any preference network $N = (V, \Pi)$, construct a graph $T(N)$ corresponding to $\mathcal{C}_{cliq}(N)$ as follows. Each vertex of $T(N)$ stands for a clique in $\mathcal{C}_{cliq}(N)$, and an edge between two vertices exists if and only if the corresponding two cliques satisfy the condition that one is a maximal clique inside the other. By Lemma 18, $T(N)$ is a tree. Now construct another tree $T'(N)$ by extending $T(N)$ in this way: for any vertex u of $T(N)$, if the corresponding clique S is such that there is only one maximal clique S' inside S , add a virtual vertex v corresponding to $S - S'$ and add an edge between u, v . Obviously, $T'(N)$ remains a tree and the number of vertices of $T'(N)$ is no smaller than that of $T(N)$.

Now we view $T'(N)$ as a rooted tree whose root is the vertex corresponding to V .

The rooted tree $T'(N)$ has two properties. First, every inner vertex has degree at least 2. Second, the leaves are disjoint subsets of V , implying that there are at most $|V|$ leaves. It is easy to see that such a tree has at most $|V| - 1$ inner vertices. As a result, the number of vertices of $T'(N)$ is at most $2|V| - 1$, meaning that the size of $\mathcal{C}_{cliq}(N)$ is at most $2|V| - 1$. \square

Lemma 20: For any finite set V , there exists a preference profile Π such that the size of $\mathcal{C}_{grow} \cap \mathcal{C}_{scomp}(N)$ is $\Omega(|V| \log |V|)$, where $N = (V, \Pi)$.

This lemma will be proved constructively. Before presenting the proof, we show the basic idea of the construction, and briefly explain how the construction algorithm works.

Intuitively, to maximize the size of $\mathcal{C}_{grow} \cap \mathcal{C}_{scomp}(N)$, we should maximize the size of $\mathcal{C}_{cliq}(N)$. Besides, in order to absorb as many individuals as possible, for any two members inside a community, they should have the same preference order on the individuals outside of the community. For example, let $N = (V, \Pi)$ be a preference network where $V = \{1, 2, 3, 4, 5\}$ and $\pi_1 = [1, 2, 3, 4, 5]$, $\pi_2 = [2, 1, 3, 4, 5]$, $\pi_3 = [3, 4, 1, 2, 5]$, $\pi_4 = [4, 3, 1, 2, 5]$ and $\pi_5 = [1, 2, 3, 4, 5]$, then $\{1, 2, 3\}$ and $\{3, 4, 1\}$ are in $\mathcal{C}_{grow} \cap \mathcal{C}_{scomp}(N)$ but not in $\mathcal{C}_{cliq}(N)$.

Following the above idea, we design Algorithm 3 which outputs the desired preference profile. Let's first introduce the notations.

Throughout the algorithm, σ stands for an arbitrarily fixed total order on V . For any $S \subseteq V$, $\sigma|_S$ is defined to be the total order restricted to S . For example, if $S = \{1, 2, 3, 5\} \subset V = \{1, 2, \dots, 5\}$ and $\sigma = [1, 4, 3, 2, 5]$, then $\sigma|_S$ is $[1, 3, 2, 5]$. For a preference network $N = (V, \Pi)$ and $v \in V$, we use $\Pi|_v$ to stand for the preference of v in N . For any disjoint finite sets W', W'' and preferences $\pi' \in L(W')$, $\pi'' \in L(W'')$, the concatenation of π' with π'' , denoted by $\pi' \bullet \pi''$, is defined to be the preference $\pi \in L(W' \cup W'')$ as follows: $\forall w' \in W'$, $\pi(w') = \pi'(w')$ and $\forall w'' \in W''$, $\pi(w'') = |W'| + \pi''(w'')$. Intuitively, π locally preserves the ordering of π' and π'' on W' and W'' , but globally prefers W' to W'' .

Algorithm 3 Find a preference profile with big $\mathcal{C}_{grow} \cap \mathcal{C}_{scomp}$

Input A set V

Output A preference profile Π on V

```

1: procedure GetProfile( $V$ )
2:   if  $V$  is a singleton  $\{v\}$  then  $\pi_v \leftarrow [v]$ 
3:   else
4:     Partition  $V$  into  $V_1$  &  $V_2$  s.t.  $0 \leq |V_1| - |V_2| \leq 1$ 
5:      $\Pi_1 \leftarrow \text{GETPROFILE}(V_1)$ 
6:      $\Pi_2 \leftarrow \text{GETPROFILE}(V_2)$ 
7:     for  $v \in V_1$  do
8:        $\pi_v \leftarrow \Pi_1|_v \bullet (\sigma|_{V_2})$ 
9:     for  $v \in V_2$  do
10:       $\pi_v \leftarrow \Pi_2|_v \bullet (\sigma|_{V_1})$ 
11:   return  $\Pi = \{\pi_v\}_{v \in V}$ 

```

The procedure *GetProfile*(V) is in a divide-and-conquer style. It divides V into two subsets V_1 and V_2 which are balanced in size. Then it recursively determine the preference profiles Π_1 and Π_2 on V_1 and V_2 , respectively. Finally, the preferences in Π_1 and Π_2 are extended, resulting in the preference profile Π on V . The extension of each preference in Π_i is by concatenating it with $\sigma|_{V_{3-i}}$, for $i \in \{1, 2\}$.

Now we prove that the preference profile Π meets the requirement of Lemma 20.

Proof. Fix V throughout this proof. Let $N = (V, \Pi)$ with Π being the preference profile found by the algorithm. For any set $U \subseteq V$ that appears in the recursion, if $|U| > 1$, the algorithm divides it into U_1 and U_2 . Let $\mathcal{K}(U) = (\mathcal{C}_{grow} \cap \mathcal{C}_{scomp})(N) \cap 2^U$, the set of communities of N inside U . Define $F(U) = |\mathcal{K}(U)|$, and $g(U) = |\mathcal{K}(U) \setminus (2^{U_1} \cup 2^{U_2})|$. Obviously, $F(U) = F(U_1) + F(U_2) + g(U)$ always holds. The show that $g(U)$ is linear in the size of U , we make two claims.

Claim 1: $|(\mathcal{C}_{grow}(N) \cap 2^U) \setminus (2^{U_1} \cup 2^{U_2})| = \Theta(|U|)$.

This claim can be proved in two steps.

First, by Algorithm 3, each clique $C \in \mathcal{C}_{cliq}(N) \cap 2^U$ with $C \neq U$ must satisfy either $C \subseteq U_1$ or $C \subseteq U_2$. Recalling the **for** loop of Algorithm 2, we have that if $S \in \mathcal{C}_{grow}(N)$ is obtained by extending $S' \subsetneq U_1$ with some $u \in V - S'$, then $u \in U_1$. So, any community $S \in (\mathcal{C}_{grow}(N) \cap 2^U) \setminus (2^{U_1} \cup 2^{U_2})$ must be extended from (hence include) either U_1 or U_2 .

Second, recall how Algorithm 2 extends U_1 with individual in U_2 . Suppose $\sigma|_{U_2} = [u_1, \dots, u_{|U_2|}]$. Since $|U_1| \geq |U_2|$ and $\pi_v|_{U_2} = \pi_{v'}|_{U_2} = \sigma|_{U_2}$ for any $v, v' \in U_1$, one knows that for any community $S \in 2^U$ with $U_1 \subset S$, it is a community in $\mathcal{C}_{grow}(N)$ if and only if $S = U_1 \cup \{u_1, \dots, u_k\}$ for some $k \leq |U_2|$. Likewise, any community $S \in 2^U$ with $U_2 \subset S$ is a community in $\mathcal{C}_{grow}(N)$ if and only if $S = U_2 \cup \{v_1, \dots, v_k\}$ for some $k \leq |U_1|$, where the $v_i \in U_1$ are such that $\sigma|_{U_1} = [v_1, \dots, v_{|U_1|}]$.

As a result, Claim 1 holds.

Claim 2: If $U \neq V$, there are at most one community in $(\mathcal{C}_{grow}(N) \cap 2^U) \setminus (2^{U_1} \cup 2^{U_2}) \setminus \mathcal{C}_{scomp}(N)$.

This claim can be proved as follows. Arbitrarily choose $S \in (\mathcal{C}_{grow}(N) \cap 2^U) \setminus (2^{U_1} \cup 2^{U_2})$. We begin with the case $U_1 \subseteq S$.

Suppose that $S \notin \mathcal{C}_{scomp}(N)$. Then there is $T \subsetneq S$ such that $V - S \succ_{\pi_u} T$ for any $u \in S - T$. Since $U \neq V$, $|V - U| \geq |U| - 1 \geq |T|$. We again go in two steps.

First, arbitrarily choose $v \in T \cap U_1$ if $T \cap U_1 \neq \emptyset$, otherwise choose $v \in T \cap U_2$. Suppose that $U_1 \not\subseteq T$; arbitrarily choose $u \in U_1 \setminus T$. By Algorithm 3 and the characterization of $(\mathcal{C}_{grow}(N) \cap 2^U) \setminus (2^{U_1} \cup 2^{U_2})$ in the proof of Claim 1, we have $v \succ_{\pi_u} w$ for any $w \in V - S$. Considering that $|V - S| \geq |V - U| \geq |U| - 1 \geq |T|$ and $V - S \succ_{\pi_u} T$, it holds that $w \succ_{\pi_u} v$ for some $w \in V - S$, which is a contradiction. Hence, $U_1 \subseteq T$.

Second, arbitrarily choose $u \in S - T \subseteq U$. By Algorithm 3, for any $v \in U, w \in V - U$, it holds that $v \succ_{\pi_u} w$. Recall that $V - S \succ_{\pi_u} T$ and $|V - S| \geq |T|$, so for any $v \in T$, there is a $w_v \in V - S$ such that $w_v \succ_{\pi_u} v$. Because $V - S = (U_2 - S) \cup (V - U)$, $T = (T - U_1) \cup U_1$, and $|U_2 - S| < |U_2| \leq |U_1|$, we have $w_v \in V - U$ for some $v \in T$, contradictory to the fact that $v \succ_{\pi_u} w$ for any $v \in T \subseteq U, w \in V - U$.

Therefore, for any $S \in (\mathcal{C}_{grow}(N) \cap 2^U) \setminus (2^{U_1} \cup 2^{U_2})$, if $U_1 \subseteq S$, we get $S \in \mathcal{C}_{scomp}(N)$.

Likewise, if $U_2 \subseteq S$, we can show that $S \in \mathcal{C}_{scomp}(N)$ except the unique case $S = \{u\} \cup U_2$ with $u \in U_1$.

As a result, by Claim 1, Claim 2 holds.

Claims 1 and 2 implies that $g(V) = O(|V|)$ and for any set $U \subsetneq V$ that appears in the recursion of Algorithm 3, $g(U) = \Theta(|U|)$. A simple calculation indicates that $F(V) = \Theta(|V| \log |V|)$. \square

Remark 1: Actually, Borgs *et al.* [2016] did not define enumerable or samplable community functions, but we find the definitions in [Borgs *et al.*, 2014], the arXiv version of [Borgs *et al.*, 2016]. Borgs *et al.* [2016] also defined the stability and the open problem required that the community function has stable communities which are samplable and enumerable. We find that even \mathcal{C}_{cliq} is not stable under their definition, so we just ignore it.

Remark 2: Lemma 8 (the intersection lemma) plays an important role in solving the problem $\mathcal{S.ACCSE}$ because \mathcal{C}_{grow} itself does not satisfy SGS . To see this, consider Example 1. Since $\{1, 2, 3\}$ is in $\mathcal{C}_{cliq}(N)$ and $\{1, 2, 3, 4\}$ is in $\mathcal{C}_{harmon}(N)$, $\{1, 2, 3, 4\}$ is in $\mathcal{C}_{grow}(N)$, while it is not in $\mathcal{C}_{SGS}(V, \Pi)$ because π_4 weekly prefers $\{5\}$ to $\{1, 2, 3\}$.

According to Theorem 3, if one uses axioms GS and SA , then it is coNP-complete to determine whether a given community is in $\mathcal{C}_{grow} \cap \mathcal{C}_{GS} \cap \mathcal{C}_{SA}(N)$. This example also shows the importance of getting rid of the computational difficulties of checking the key axioms of GS and SA .

VI. CONCLUSION

In this paper, we focus on axiomization of network community detection. As far as we know, only one paper [Borgs *et al.*, 2016] in this line studied overlapping communities, and it was on a general structure – preference networks. Among the six axioms in [Borgs *et al.*, 2016], two of them (GS and SA) play a critical role, but the hardness in checking them compromises the practicality of the axiom system. We showed that in the context of the axiom system, SA can be

equivalently replaced by a degenerate version SA' . We also naturally modified axiom GS to a stronger version called SGS . We showed that both SGS and SA' can be checked in polynomial time, and most of the properties of the original system are preserved. Furthermore, by the intersection lemma, we found two constructive and SA -consistent community functions. We also found an SA -consistent, constructive, samplable and enumerable community function that allows nontrivially overlapping communities, thus answering to an open problem in [Borgs *et al.*, 2016].

Although the current work is purely theoretical, it would be beneficial to evaluate and improve existing community detection algorithms in the framework of our axioms. This is a direction of our future work.

Another direction is to further improve the axiom system. Before Lemma 10, we discussed the richness, namely, many community functions might be consistent with the axioms. This seems advantageous since good community functions are not likely to be precluded. However, selectivity should also be considered, otherwise users would be burdened with selecting desirable community functions from too many candidates. Can we further improve the system of axiom by making a better trade-off between richness and selectivity?

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