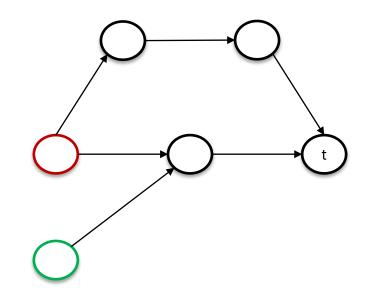
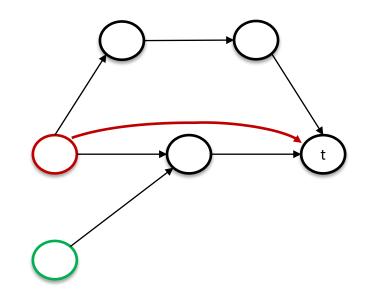
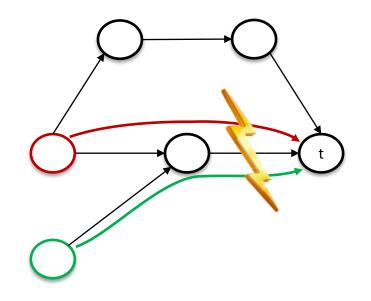
#### Augmenting Anycast Network Flows

Sebastian Brandt, <u>Klaus-Tycho Förster</u>, Roger Wattenhofer January 06, 2016 @ ICDCN 2016 - Singapore

ETH Zurich – Distributed Computing – www.disco.ethz.ch



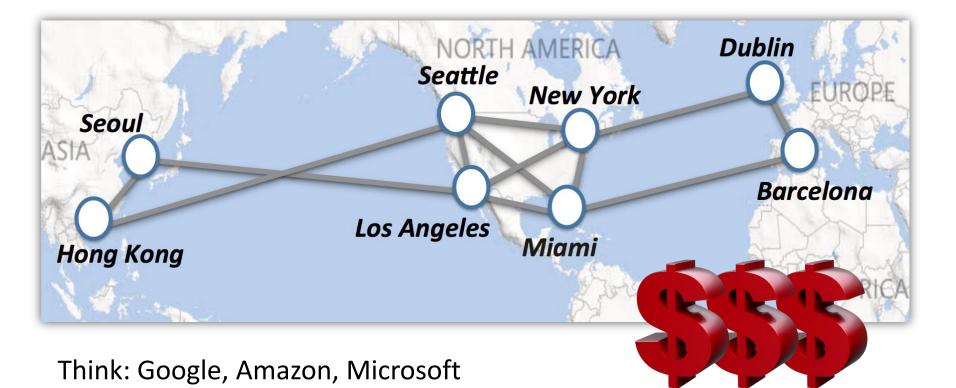




# **Network Updates**

- The Internet: Designed for selfish participants
  - Often inefficient (low utilization of links), but robust
- But what happens if the WAN is controlled by a single entity?
  - Examples: Microsoft & Amazon & Google ...
  - They spend hundreds of millions of dollars per year

# **Network Updates**



# **Network Updates**

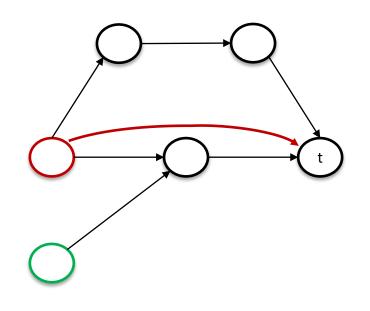
- The Internet: Designed for selfish participants
  - Often inefficient (low utilization of links), but robust
- But what happens if the WAN is controlled by a single entity?
  - Examples: Microsoft & Amazon & Google ...
  - They spend hundreds of millions of dollars per year
- Possible solution: Software Defined Networking (SDN)

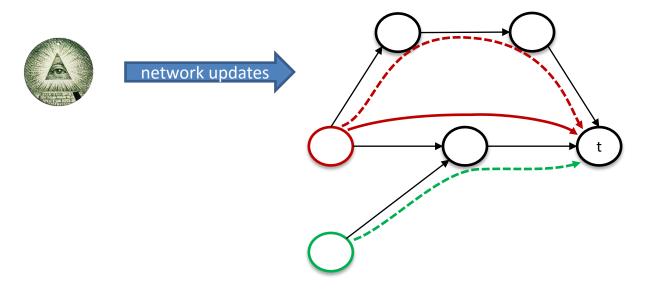


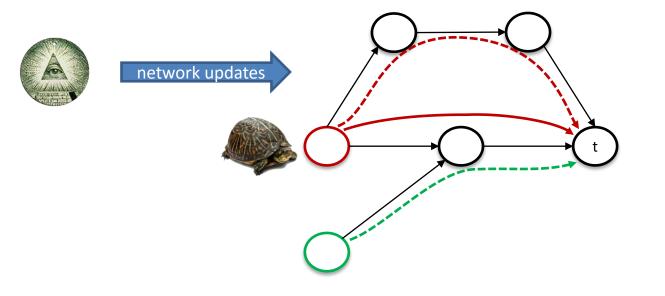
- General Idea: Separate data & control plane in a network
- Centralized controller updates networks rules for optimization
  - Controller (control plane) updates the switches/routers (data plane)

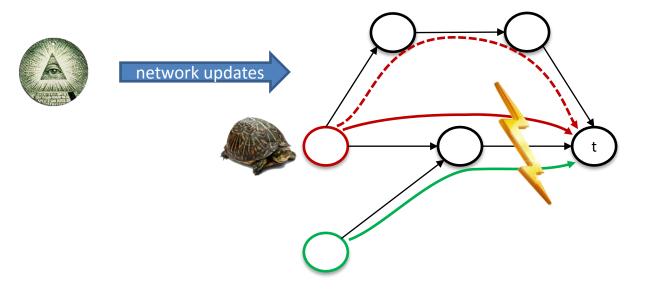


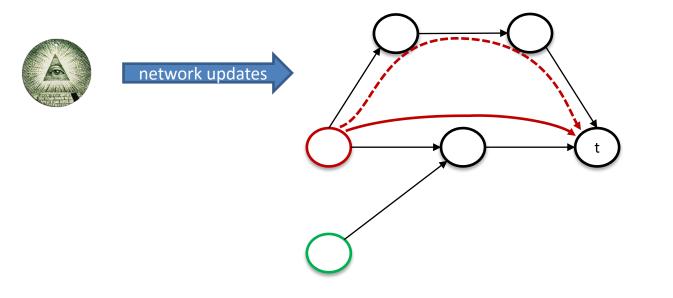


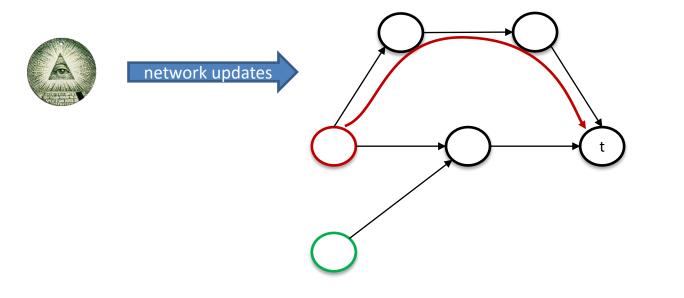


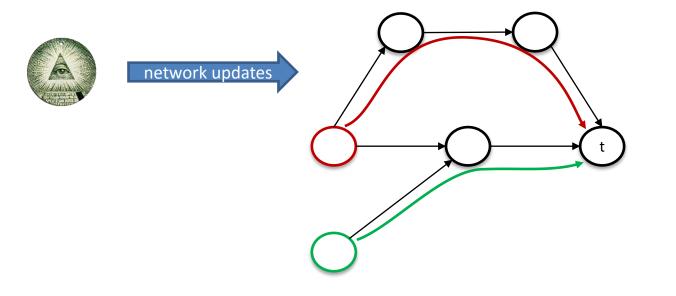








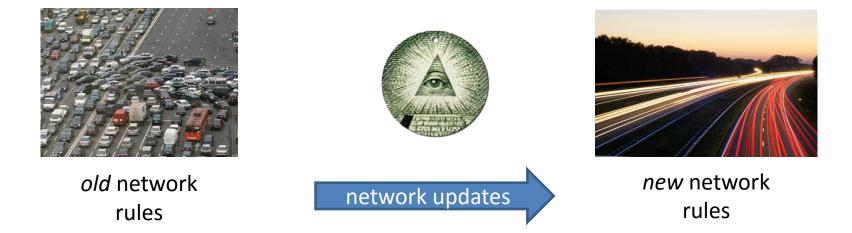




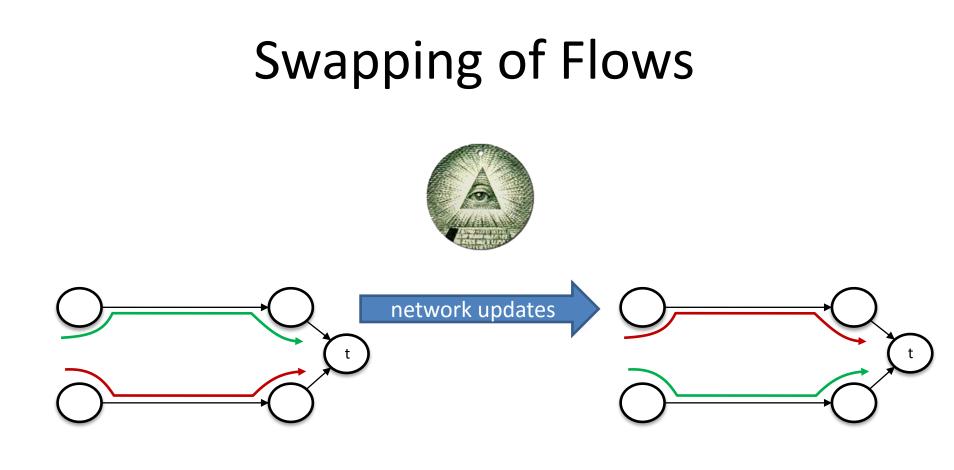
# Structure of the Talk

- Motivation & Software Defined Networking
- Related Work & Splittable Flows
- Our Approach
- Extension beyond Anycast Flows

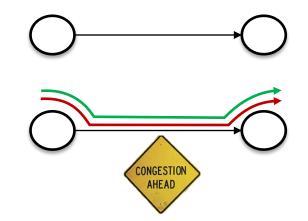
#### Network Updates of Flows without Congestion



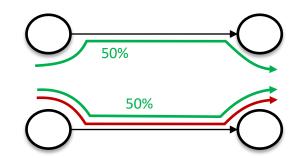
- State of the art: (Partial) moves of flows using linear programming (LPs), e.g.,
  - SWAN [Hong et al., SIGCOMM 2013], *zUPDATE* [Liu et al., SIGCOMM 2013]
  - Dionysus [Jin et al., SIGCOMM 2014]
- Open problems:
  - When are network updates without congestion *possible*?
  - How can we do them *fast*?
- *This paper*: Addresses the case of one (logical) destination for splittable flows



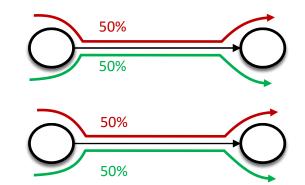
# Just Switch? Congestion!



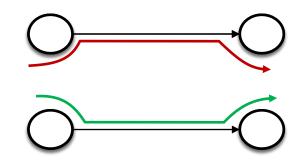
# Migrate only parts of the flow



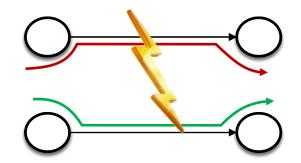
# Can even do both flows at once



### Done in two steps



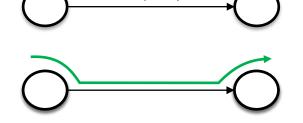
# But not always possible!



# How about other paths?



Number of steps can be unbounded



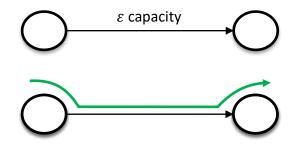
 $\varepsilon$  capacity

# How about other paths?



Number of steps can be unbounded

Binary search with LPs ineffective



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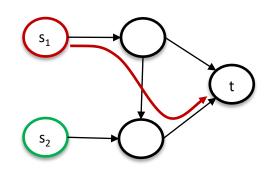
# Our Approach

• Compute our **own** new rules

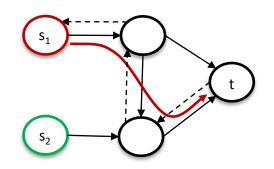
Based off the new demands

- Deviate from linear programming binary search
  Go combinatorial with augmenting flows
  - "Push back" flows for migration

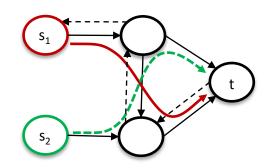
## **Augmenting Flows**



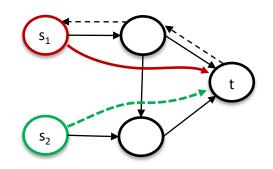
# **Consider Residual Network**



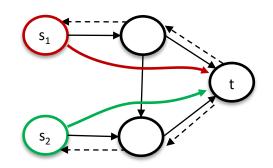
# Find a Way in the Residual Network



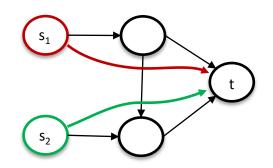
### Push back the old Flow



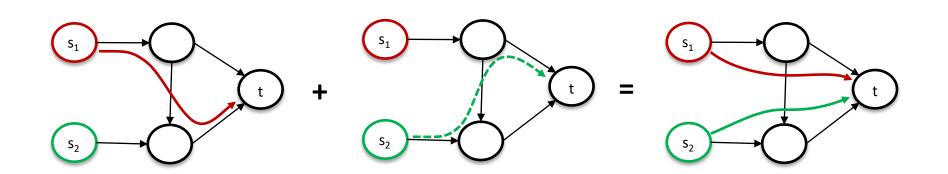
#### Insert the new Flow



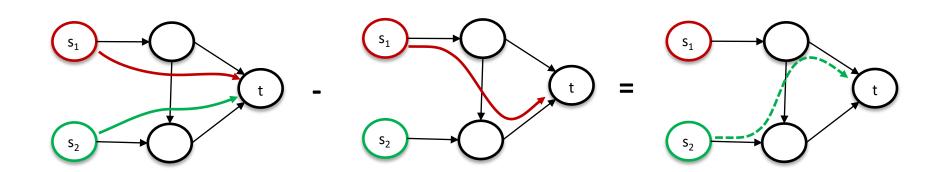
# Migrated without Congestion



# Similar to "Addition"



# Also works as "Subtraction"



# High-level Mechanism Idea

For all commodities (iteratively):

- Increase demand and calculate new flow with LP
  - *offline* calculation
- Apply augmenting flow from the difference
  - linear # re-routing *in the network*

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OLD	CURRENT	NEW
1	1	5
1	1	2
2	2	6
3	3	3
2	2	8

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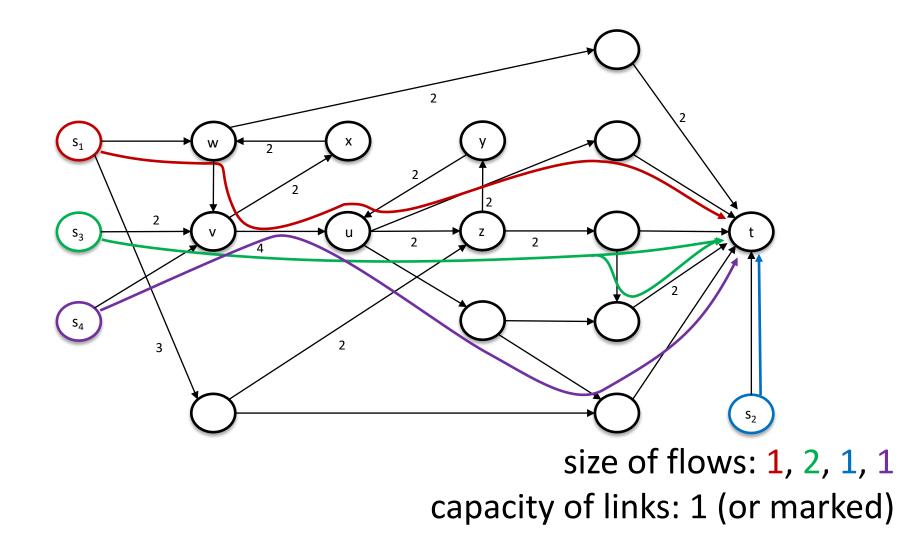
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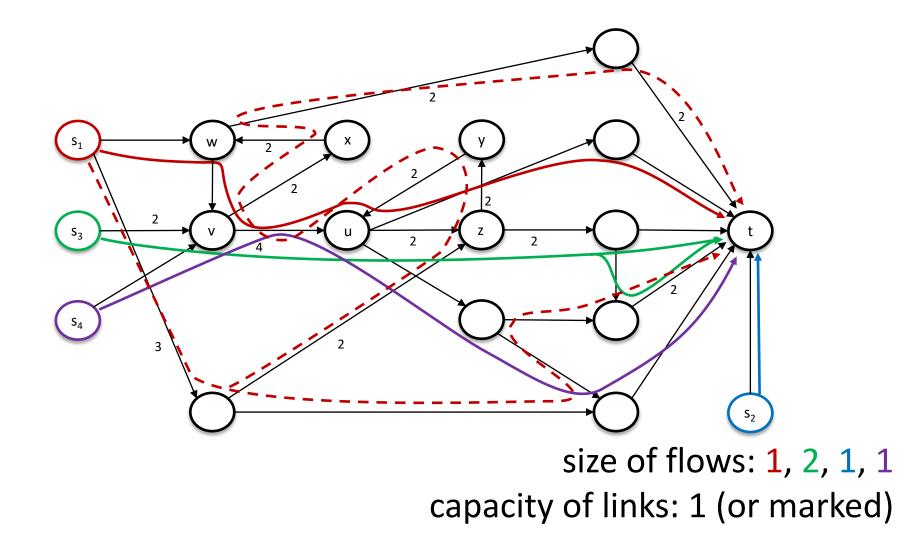
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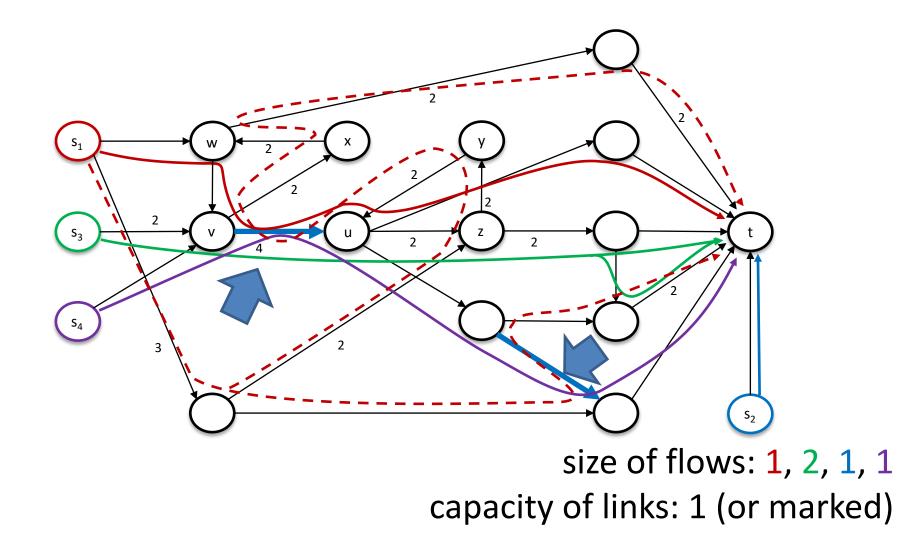
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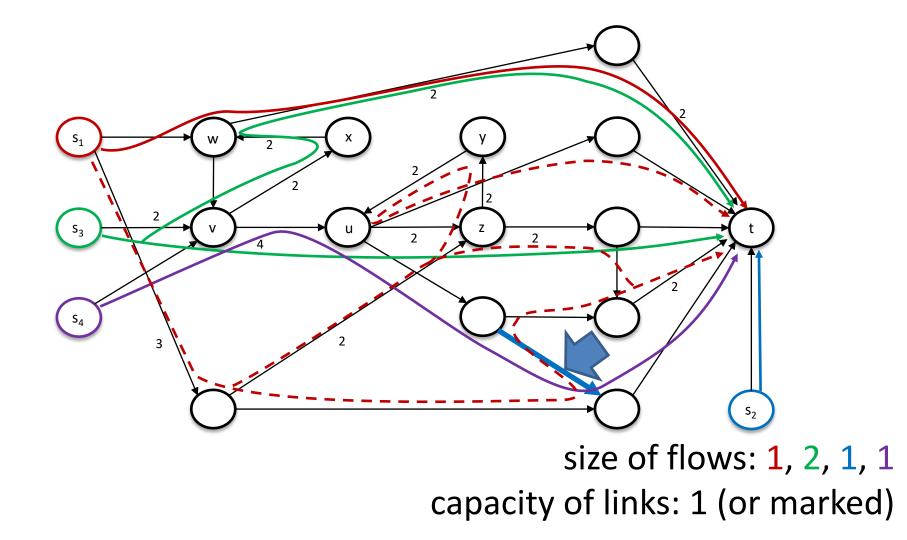
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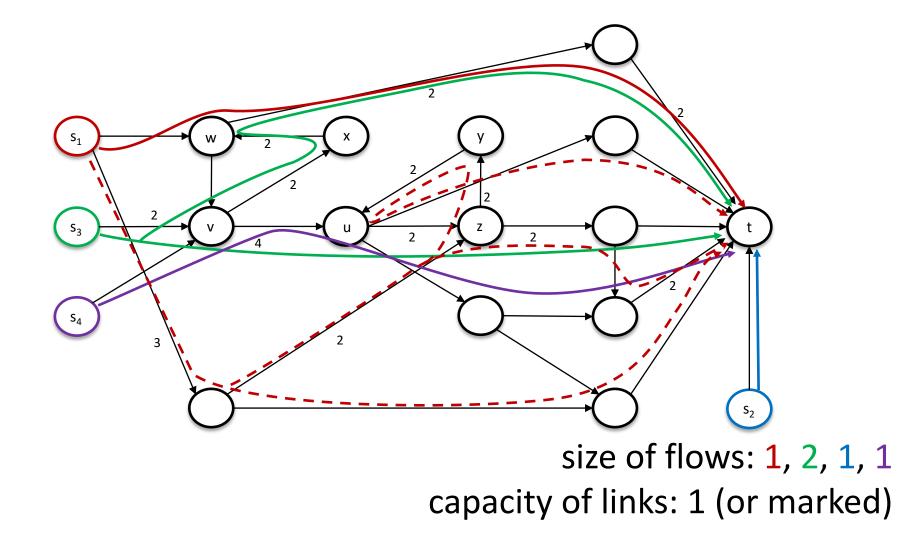
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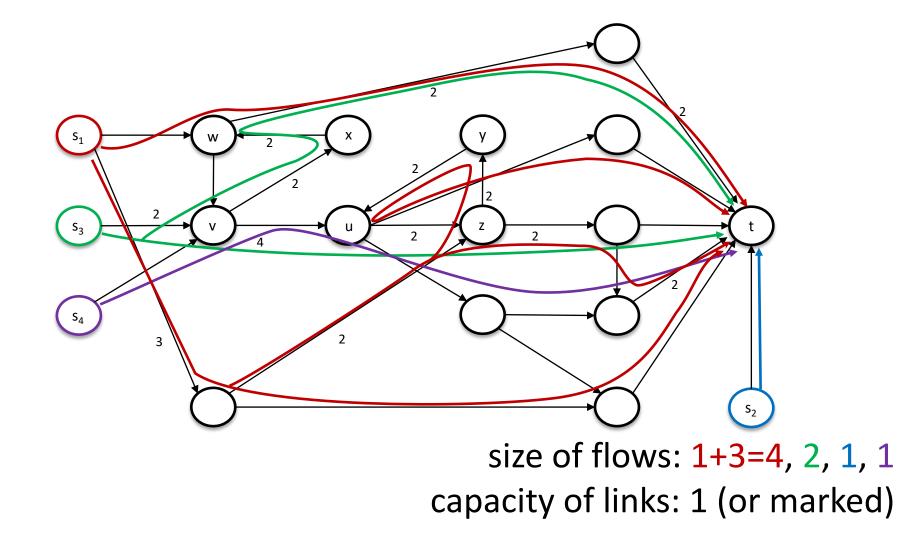








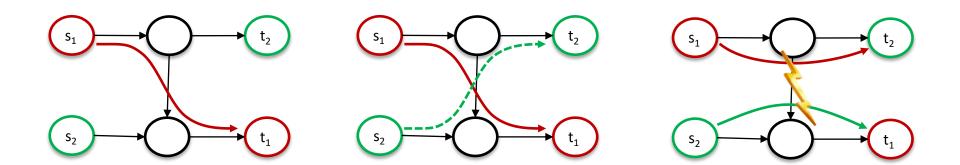




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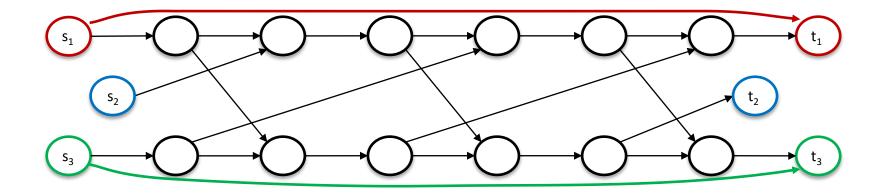
### Flow Augmentation for many Destinations



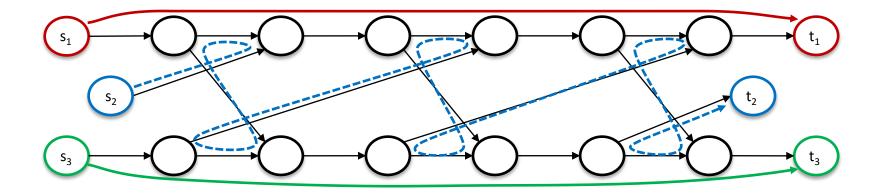
• Flows end up at the wrong destination!

• So let's stick with augmenting flows that don't mix destinations

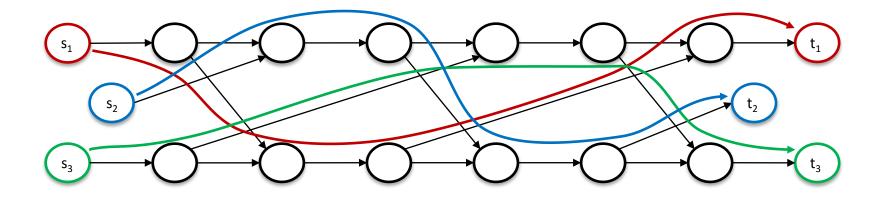
#### Extension beyond one logical destination?



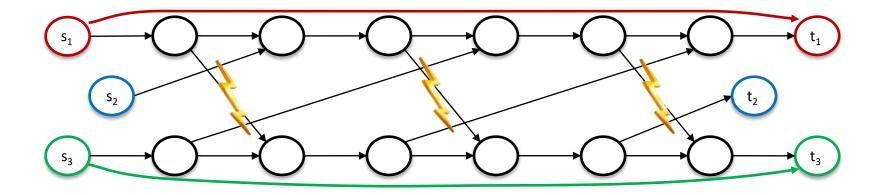
#### Augmenting flows that don't mix up the destinations?



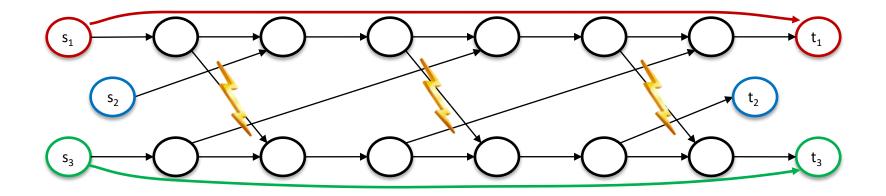
#### Augmenting flows that don't mix up the destinations?



#### But impossible to migrate!



#### But impossible to migrate!



"it is unlikely that similar techniques can be developed for constructing multicommodity flows" [Hu, 1963]

# Summary

- We studied how to migrate flows with one logical destination in SDNs without congestion
  - We can decide fast if demands can be met
  - We can migrate with linear # re-routings in the network using augmenting flows
- Open question:

– How to extend beyond one logical destination?

### Thank you

Sebastian Brandt, <u>Klaus-Tycho Förster</u>, Roger Wattenhofer January 06, 2016 @ ICDCN 2016 - Singapore

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