## How Many Ants Does It Take to Find the Food?

Jara Uitto

ETH Zurich – Distributed Computing – www.disco.ethz.ch

### Ants Nearby Treasure Search

 Introduced by Feinerman, Korman, Lotker and Sereni [PODC 2012].

• *n* mobile agents, controlled by Turing machines, search for a treasure.

• Communication not allowed.



- Infinite integer grid.
- Each ant initially located in the origin.



- Adversarially hidden treasure/food.
- (Manhattan) distance to treasure is D.

### Ants Nearby Treasure Search

• How many rounds until the treasure is found?

 We study the number of ants needed to find the treasure at all.











• One Turing Machine is enough. No communication needed.



 Ants are controlled by (randomized) finite state machines.



- Communicate by sensing the states of nearby ants.
- Run-time studied by Emek, Langner, Uitto and Wattenhofer [ICALP2014].



• Synchrony vs. Asynchrony

• A deterministic protocol?



 Individual algorithm for each ant.

 An algorithm works correctly if the ants find the treasure in expected finite time.



#### Deterministic + Asynchronous


















































# **Triangle Search**



# Synchronization?

• Can we perform better if the ants have a common sense of time?









































#### Randomization

• How about random coin tosses?









NE





#### NE 11





















#### NE 1110110
### Geometric Search



# **Run-Time**

• For every search *i*, we have a probability of at least  $A_i = \frac{1}{4} \cdot 2^{-(D+1)}$  to find the treasure.

• Let  $B_i$  be the event that the treasure is not found during any search j < i.

# **Run-Time**

- Let *T* be the total time required.
- $E[T] \leq \sum_{i=1}^{\infty} P(A_{i+1} \cdot B_i) (O(i) + O(D)).$

• 
$$P(A_{i+1} \cdot B_i) \le 2^{-(D+3)} \cdot (1 - 2^{-(D+3)})^i$$
.

•  $E[T] \le 2^{-(D+3)} \sum_{i=1}^{\infty} (1 - 2^{-(D+3)})^i (O(i) + O(D)) = O(2^D).$ 

### Lower Bounds?



• Can we do better? In the deterministic and synchronous case, the answer is no.

 Let us start with showing that one ant is not enough.



• A finite state machine repeats its behavior.









A band of constant width

• One ant can only discover a band of constant width.



• Let *t* be the time of the last meeting.



• Both agents (alone) discover a band after t.

• Lemma: The ants meet infinitely often in some pair of states (q, q').

• Observation: the time between two such meetings is bounded by a constant.













- Two deterministic ants can only discover a band of constant width.
- Two deterministic ants cannot find the food.

# Conclusion

	$\mathbf{FA}$				
Problem	Sy	ync	async		
	det	rand	det	rand	
One agent		×		×	
Two agents	$\times$	?	$\times$	?	
Three agents	$\checkmark$	$\checkmark$	?	$\checkmark$	
Four agents			$\checkmark$		

- Three asynchronous ants?
- Two randomized ants?

## Conclusion

	$\mathbf{PDA}$				
Problem	sync		async		
	det	rand	det	rand	
One agent	×	$\checkmark$	×	$\checkmark$	
Two agents	$\checkmark$		$\checkmark$		
Three agents					
Four agents					



Thanks to my co-authors Yuval Emek, Tobias Langner, David Stolz and Roger Wattenhofer