## How Many Ants Does It Take to Find the Food?



## Ants Nearby Treasure Search

- Introduced by Feinerman, Korman, Lotker and Sereni [PODC 2012].
- $n$ mobile agents, controlled by Turing machines, search for a treasure.
- Communication not allowed.


## Model



- Infinite integer grid.
- Each ant initially located in the origin.


## Model



- Adversarially hidden treasure/food.
- (Manhattan) distance to treasure is $D$.


## Ants Nearby Treasure Search

- How many rounds until the treasure is found?
- We study the number of ants needed to find the treasure at all.


Model


Model


## Model



- One Turing Machine is enough. No communication needed.



## Model

- Ants are controlled by (randomized) finite state machines.
- Communicate by sensing the states of nearby ants.
- Run-time studied by Emek, Langner, Uitto and Wattenhofer [ICALP2014].



## Model

- Synchrony vs. Asynchrony

- A deterministic protocol?


## Model

- Individual algorithm for each ant.

- An algorithm works correctly if the ants find the treasure in expected finite time.



## Deterministic + Asynchronous



## Triangle Search



## Triangle Search



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## Synchronization?

- Can we perform better if the ants have a common sense of time?



## Rectangle Search



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## Randomization

- How about random coin tosses?



## Geometric Search



## Geometric Search



NE

## Geometric Search



NE 1

## Geometric Search



NE 11

## Geometric Search



NE 111

## Geometric Search



NE 1110

## Geometric Search



NE 11101

## Geometric Search



NE 111011

## Geometric Search



NE 1110110

## Geometric Search



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## Geometric Search



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## Geometric Search



## Run-Time

- For every search $i$, we have a probability of at least $A_{i}=\frac{1}{4} \cdot 2^{-(D+1)}$ to find the treasure.
- Let $B_{i}$ be the event that the treasure is not found during any search $j<i$.


## Run-Time

- Let $T$ be the total time required.
- $E[T] \leq \sum_{i=1}^{\infty} P\left(A_{i+1} \cdot B_{i}\right)(O(i)+O(D))$.
- $P\left(A_{i+1} \cdot B_{i}\right) \leq 2^{-(D+3)} \cdot\left(1-2^{-(D+3)}\right)^{i}$.
- $E[T] \leq 2^{-(D+3)} \sum_{i=1}^{\infty}\left(1-2^{-(D+3)}\right)^{i}(O(i)+O(D))=$ $O\left(2^{D}\right)$.


## Lower Bounds?



- Can we do better? In the deterministic and synchronous case, the answer is no.
- Let us start with showing that one ant is not enough.


## One Ant



- A finite state machine repeats its behavior.

One Ant

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One Ant

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One Ant

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## One Ant



A band of constant width

## One Ant

- One ant can only discover a band of constant width.
- How about two ants?



## Two Ants

- Let $t$ be the time of the last meeting.

- Both agents (alone) discover a band after $t$.


## Two Ants

- Lemma: The ants meet infinitely often in some pair of states $\left(q, q^{\prime}\right)$.
- Observation: the time between two such meetings is bounded by a constant.


## Two Ants

$\left(q, q^{\prime}\right)$


## Two Ants



## Two Ants



## Two Ants



## Two Ants



- Two deterministic ants can only discover a band of constant width.
- Two deterministic ants cannot find the food.


## Conclusion

## FA



- Three asynchronous ants?
- Two randomized ants?


## Conclusion

## PDA

| Problem | sync |  | async |  |
| :--- | :---: | :---: | :---: | :---: |
|  | det | rand | det | rand |
| One agent | $\times$ | $\checkmark$ | $\times$ | $\checkmark$ |
| Two agents | $\checkmark$ |  | $\checkmark$ |  |
| Three agents |  |  |  |  |
| Four agents |  |  |  |  |

## Questions?



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